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IMPLEMENTATION OF THE ANALYSIS PART
OF A PASCAL TRANSLATOR

by

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B.S. in C.S., Middle East Technical University, 1981

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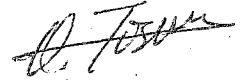
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IMPLEMENTATION OF THE ANALYSIS PART
OF A PASCAL TRANSLATOR

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A B S T R A C T

This thesis describes the implementation of the analysis part of a single pass translator for the programming language STANDARD PASCAL. The translation process includes syntactic analysis and semantic analysis at the declaration level, but excludes code generation. The details of the implementation are described, with a short description of alternative approaches in each section. The source and test runs of the program are given in the appendixes A and B respectively.

Ö Z E T

Bu tez, programlama dili STANDARD PASCAL için geliştirilen, tek geçişli bir çeviricinin gerçekleştirimini anlatır. Çeviri işlemi, sözdizimsel ve kısmen anlamsal çözülemeyi içerir, ancak kod üretimi göz önüne alınmamıştır. Projenin detayları her bölümde değişik yaklaşım yöntemleri belirtilerek açıklanmıştır. Kaynak izlenice ve test geçişleri, sırasıyla Ek A ve Ek B de verilmiştir.

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LIST OF SYMBOLS

<...> The angular brackets are used to distinguish non-terminals from terminals. The element delimited by these brackets is supposed to be a non-terminal.

\Longrightarrow For a grammar G , we say that the string v directly produces the string w , written

$$v \Longrightarrow w$$

if we can write

$$v = xUy \text{ and } w = xuy$$

for some strings x and y , where $U ::= u$ is a rule of G .

$\Longrightarrow +$ For a grammar G , we say that the string v produces the string w , written

$$v \Longrightarrow + w$$

if there exists a sequence of direct derivations

$$v \Longrightarrow u_0 \Longrightarrow u_1 \text{ -----} \Longrightarrow U_n$$

where $n > 0$.

$\Longrightarrow *$ This symbol can be equivalently written as;

$$v = w \text{ or } v \Longrightarrow + w$$

$::=$ Used as an abbreviation for "defined as".

I. INTRODUCTION

This thesis describes a partially completed PASCAL translator (BUPASCAL) implemented on the UNIVAC 1106 of Boğaziçi University. The program is written in standard PASCAL and occupies approximately 60 K words of memory.

BUPASCAL accepts programs written in STANDARD PASCAL as input and generates a listing of its input with the associated error messages. After each input program line is processed, error codes corresponding to the errors fetched (if any) are printed with their positions on the line, just after the line is printed. For example, errors of line 5 of the following statements are printed as;

<u>line#</u>	<u>source lines</u>
4	J := 3
5	K := L + ; M := 11 ;
	1 2 3

ERROR 14 ... POSITION 1
ERROR 2 ... POSITION 2
ERROR 205 ... POSITION 3

6 if A > 10 then C := 3 ;

When the input program is completely processed, explanations of the errors detected are printed as follows;

ERROR EXPLANATIONS

***ERROR 14 : ';' EXPECTED
***ERROR 2 : IDENTIFIER EXPECTED
***ERROR 205 : NULL STRING NOT ALLOWED

Translator is a misleading term for BUPASCAL since it is actually an analyzer and it does not generate code. In addition to error detection and reporting, BUPASCAL collects information about the identifiers of the input program, into a symbol table. This table is optionally dumped as the source listing is produced.

The major components of BUPASCAL are

- (1) Lexical analyzer -- described in section II
- (2) Syntactic analyzer -- described in section III
- (3) Error recovery -- described in section IV
- (4) Symbol table -- described in section V

II. LEXICAL ANALYSIS

A. Definition

The lexical analyzer is a part of a translator which reads the source program one character at a time and constructs the source program symbols (identifiers, keywords, constants and delimiters).

One may justly ask, why the lexical analysis can not be incorporated into the syntax analysis. After all, we can use BNF to describe the syntax of symbols. For example, PASCAL identifiers can be described by

```

<identifier> ::= <letter> { <letter> | <digit>}o
<letter>     ::= A|B|C|... |Z
<digit>     ::= 0|1| ... |9

```

There are several reasons for splitting the analysis of the source program into two phases, lexical analysis and syntactic analysis;

1. Main purpose is to simplify the overall design of the translator. Considering a large portion of compile-time is spent in scanning characters, by separation we can concentrate solely on reducing this time.
2. The syntax of symbols can be described by very simple grammars. If we separate scanning from syntax recognition we can develop efficient parsing techniques which are particularly well suited for these grammars.

3. Generally speaking, syntax analyzer requires much more programming effort than the lexical analyzer does. Our aim must be, therefore to ease life for the syntax analyzer. So, some cumbersome functions such as keeping track of line numbers, producing an output listing, ignoring blanks and comments etc. can be performed by the lexical analyzer.
4. Since the lexical analyzer returns a symbol instead of a character, the syntax analyzer actually gets more information about what to do at each step. Moreover, the steps required for the syntax checking will be decreased.
5. Separation allows us to write one syntactic analyzer and several lexical analyzers (which are simpler and easier to write) one for each source program representation and/or input device. Each lexical analyzer translates the symbols into the same internal form used by the syntactic analyzer.

A lexical analyzer may be programmed as a separate pass which performs a complete lexical analysis of the source program and which gives to the lexical analyzer a table containing the source program in an internal symbol form. Alternatively, it can be a subroutine called by the syntax analyzer, whenever the syntax analyzer needs a new symbol (Figure 2.1).

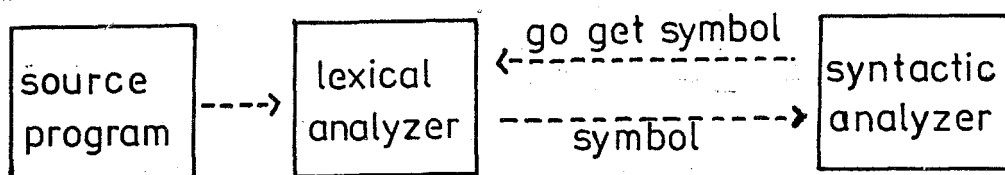


Figure 2.1- Lexical Analyzer as a Subroutine of Syntactic Analyzer.

When called lexical analyzer recognizes the next program symbol and passes it to the syntactic analyzer. This alternative is generally better, because the whole internal source program need not to be constructed and kept in memory.

In this thesis, lexical analyzer is implemented in this manner. It is a PASCAL procedure named SCAN. Implementation of SCAN will be explained in detail within the rest of this section.

B. The Output Of a Lexical Analyzer

A lexical analyzer builds an internal representation of each symbol. In most cases this is a fixed length integer. As opposed to variable length strings, which are the actual symbols, these internal representations are much easier to manipulate and parse.

We include in this internal representation a constant meaning "identifier" and another for "integer" etc. That is, all identifiers have the same internal number to represent them. This is natural, because the term "identifier" is a terminal symbol to the syntax analyzer, and which identifier it happens to be is of no consequence. However, the identifier itself is needed by the symbol table management routines, so it must be stored somewhere. Similarly, it is enough for the syntax analyzer to know an integer is met (whatever its value is), but the semantic analyzer should know its internal value.

The solution is to output two values; the first is the internal representation and the second is the actual symbol itself or a pointer to it. Thus, the output of the lexical analyzer will consist of two fields.

1. CLASS : This is the internal representation of the symbol scanned. This field is utilized by the syntax analyzer.
2. VALUE : This field gives further and more precise information pertaining to the class of the recognized symbol. Only a few classes of symbols make use of this field and it is utilized by the semantic analyzer.

The information placed in the VALUE field is dependent on the CLASS field and can sometimes be empty. VALUE field may be used in two different ways. First use is, in the case of identifiers or literal constants to keep actual symbol or internal values as explained above.

Second, in some cases further information is required to indicate which particular instance of a class, the syntactic primitive is. An example of this is the PASCAL <multiplication operation> which has five possible instances, -,*,/,div,mod, and. Although, the precise nature of the multiplication operation does not affect the syntax analyzer, it has to be known for the semantic analyzer.

C. State Transition Diagrams

Before discussing the implementation of SCAN, I would like to mention about state transition diagrams, which I believe, is the best way, to show how a symbol is to be parsed. Figure 2.2 shows a simple state diagram.

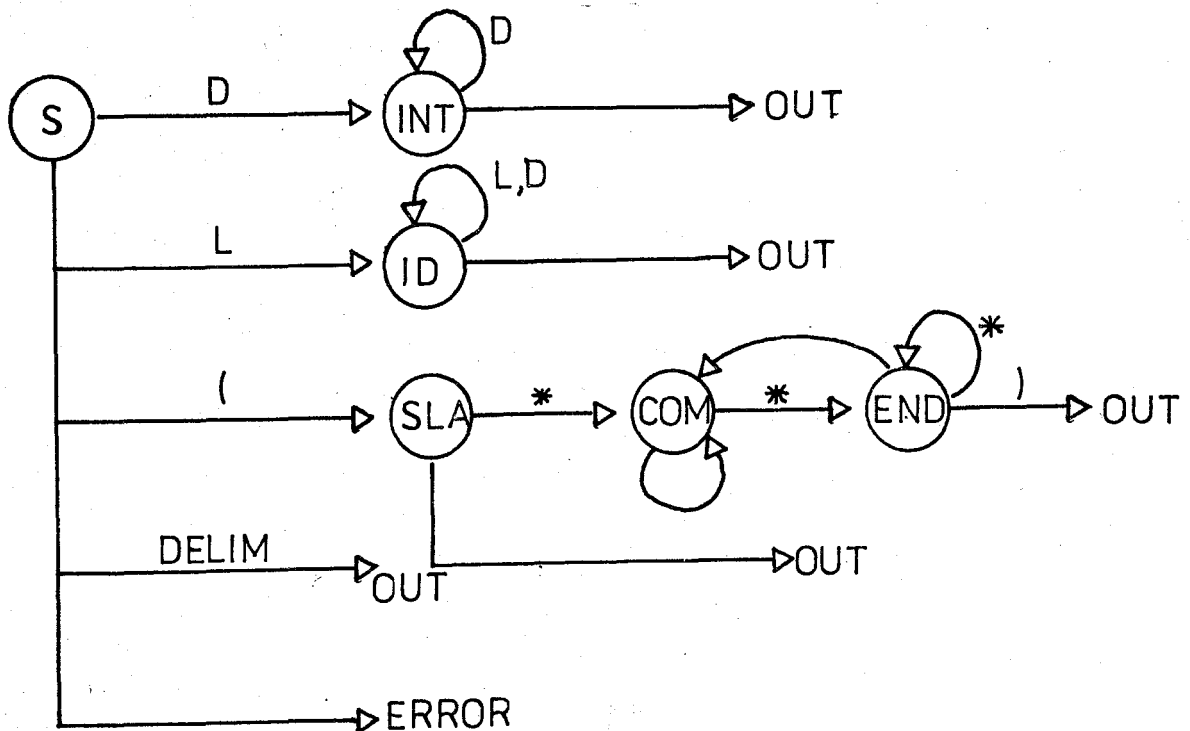


Figure 2.2- State Diagrams Scanning Symbols.

In this figure, the label D is an abbreviation for the labels 0,1,2,...,9. That is, D represents the class of digits. This is done to keep the diagram simple. Similarly, the label L represents the class of letters A,B,C,...,Z, and DELIM represents the class of single character delimiters +,-,*,/ etc. Note that, character "(" is not in this class, since it must be handled in a special manner.

Several of the arcs have no label on them. These are the arcs to be taken, if the character scanned is one that does not explicitly appear on an arc. For example, when in state INT, as long as we scan a digit we stay in state INT. If a non-digit character is scanned we proceed along the arc to OUT.

The arc OUT indicates the class and value of the symbol recognized. That is, by arc OUT we mean that we have detected the end of a symbol and we want to leave the lexical analyzer.

Having provided a transition diagram, there are two approaches, to program the lexical analyzer. First is, to represent transition diagram in core as a set of tables indexed by states and characters and then move within these tables. Second approach, is to develop the lexical analyzer algorithmically. Here we do not have any explicit tables. A transition diagram is viewed as a specialized kind of flowchart. States of a transition diagram correspond to the boxes of a flowchart. So, according to the arrows connecting these states, program structures (loops, if-then-else etc.) are designed until an out arc is reached. This approach is particularly useful when the program is a lexical analyzer, because the action taken is highly dependent on what characters have been seen recently.

The second approach is chosen (except for identifiers for which a set of transition tables is used) in the implementation of SCAN. SCAN uses a case statement based on the first character of the incoming symbol, to determine which arc to take from the initial state S. Following characters

are handled as described above.

When an ERROR arc is reached SCAN takes the corrective action and tries again. Errors detected by SCAN is referred as LEXICAL TIME errors and they will be presented under the heading E.

D. Implementation

Within this section, the PASCAL procedure SCAN, which is the lexical analyzer of BUPASCAL will be explained.

1. Symbols of PASCAL

Symbols of PASCAL can be classified into five groups. In each case the class field represents the type of syntactic primitive, and the value field, if utilized, supplies further information.

i) IDENTIFIERS

Identifiers can be predefined or user defined. Predefined identifiers are

- Subprogram identifiers such as READ, WRITE, EOLN etc.
- Type identifiers such as INTEGER, REAL, BOOLEAN etc.
- Const identifiers such as TRUE, FALSE, MAXINT etc.

In general all identifiers (constants, variables, procedures, functions, types, fields etc) whether user defined or predefined are identified by the same class code. The only difference between predefined and user defined identifiers is that, predefined identifiers are located during initialization into the symbol table. But this is not a concern of SCAN.

Value field is unutilized. Instead a global string variable keeps the actual symbol, for symbol table search.

ii) KEYWORDS

These are the syntactic primitives that serve to identify the different types of PASCAL statements and program components.

They have the same syntax with identifiers, but SCAN separates them from identifiers, by a set of transition tables, thus eliminating a need for a search of a table of keywords. There are thirty keywords and each keyword is represented by a distinct class code. Note that, value field is again unutilized.

iii) LITERAL CONSTANTS

In this category we include unsigned integers, unsigned reals, character constants and string constants. Thus, four internal codes are required to indicate the type of the constant in the class field. The value field points to the location that contains the value of the constant in the dynamic constant table.

iv) OPERATORS

Pascal operators can be classified into three groups, each group requiring a distinct class code.

1. Relational operators

These are "<", ">", "=", "<>", ">=", "<=" and in. The value field indicates the precise operator within the group.

2. Adding operators

These are "+", "-" and or and the value field pinpoints the precise operation.

3. Multiplying operators

Similar to the previous two groups except the operators are "*", "/", div, mod, and.

Note here that the generalizations in groups (1) and (2) are not accurate. The symbol "=" when used in type and const statements is not a relational operator. Similary "+" and "-" can sometimes be used as signs.

This discrepancy is avoided by taking corrective action during semantics check.

v) DELIMITERS

In this category we include all other single character e.g. ".", ";", "(", " :") and double character delimiters (e.g. "..", " :=") of PASCAL.

Each of these delimiters will be represented by a unique class code and the value field will not be utilized.

SCAN is also responsible to detect and ignore seperators which are principally for the benefit of the human reader. Seperators are blanks, end-of-lines or comments. They have no effect on the meaning of the program.

2. Data Structures to Represent PASCAL Symbols

At this point, we can present data structures used in recognition of PASCAL symbols. These structures are declared globally (Table 2.1).

In these declarations, CLASS and VALUE fields are gathered within a record type named TOKENS. The record contains three variants which states different uses of VALUE field.

TABLE 2.1- Data Structures for the Lexical Analyzer.

```

TYPE
  SYMBOL = (* ...SYMBOL CLASSES... *)
  (
    IDENT, INTCONST, REALCONST, CHARCONST,
    STRINGCONST, EOFPGM, MULOP, ADDOP,
    RELOP, LPARENT, RPARENT, LBRACKET,
    RBRACKET, COMMA, SEMICOLON, PERIOD,
    DOTDOT, ARROW, COLON, ASSGNOP,
    ARPAYSY, BEGINSY, CASESY, CONSTSY,
    JOSY, DOWNSY, ELSESY, ENDSY,
    FILESY, FORSY, FORWARDSY, FUNCTIONSY,
    GOTOSY, IFSY, LABELSY, NOTSY,
    OFSY, PACKEDSY, PROCESY, PROGRAMSY,
    RECORDSY, REPEATSY, SETSY, THENSY,
    TOSY, TYPESY, UNTILSY, VARSY,
    WHITESY, WITHSY, OTHERSY ) ;

  OPERATOR =
  (
    ASTR, RDIV, IDIV, ANDOP,
    IMOD, PLUS, MINUS, OROP,
    LTOP, LEOP, GEOP, STOP,
    NEOP, EQOP, INOP ) ;

  LITERALS = INTCONST .. STRINGCONST ;
  OPERATORS = MULOP .. RELOP ;
  DELIMITERS = LPARENT .. ASSGNOP ;
  KEYWORDS = ARPAYSY .. WITHSY ;

  MULTIPLICATION_OPERATORS = ASTR .. IMOD ;
  ADDITIONAL_OPERATORS = PLUS .. OROP ;
  RELATIONAL_OPERATORS = LTOP .. INOP ;

  CSTCLASS = (INTGR, REEL, STRING, KHAR) ;
  STRGRANGE = 0 .. 78 ;
  CHARSTRING = PACKED ARRAY [STRGRANGE] OF CHAR ;

  CSTADDR = @ CONSTANT ;
  CONSTANT = PACKED RECORD
  (
    CASE CLASS : CSTCLASS OF
      INTGR : (IVAL : INTEGER) ;
      REEL : (RVAL : REAL) ;
      KHAR : (ORDCH : INTEGER) ;
      STRING : (SLGTH : STRGRANGE,
               SVAL : CHARSTRING) ;
    END ;
  ) ;

  (* ...FOLLOWING ENUMERATED TYPE INDICATES HOW THE VALUE
  FIELD WILL BE UTILIZED. *)
  TKNCLASS = (NOTNEEDED, OPRTR, DRADRES) ;

  (* ...OUTPUT TYPE OF SCAN IS THE FOLLOWING RECORD WHERE
  VARIANTS OP, CSTADR DENOTES DIFFERENT USES OF VALUE
  FIELD. *)
  TOKENS = RECORD
  (
    CLASS : SYMBOL ;
    CASE TKNCLASS OF
      OPRTR : ( OP : OPERATOR ) ;
      DRADRES : ( CSTADR : CSTADDR ) ;
      NOTNEEDED : ( ) ;
    END ;
  ) ;

  VAR TOKEN : TOKENS ;

```

3. Algorithm of the Lexical Analyzer

SCAN uses the following variables and routines.

1. CH : CH is a global variable which will always hold the current character of the source program being scanned.
2. IDENTIFIER: is a location which will contain the string of characters making up the symbol.
3. NEXTCHAR : is a function to return the next source program character. NEXTCHAR will take care of reading and printing next source line, when it detects end of line, while trying to scan next source character.
4. READLINE : is a procedure which reads the next source program line into an internal buffer. Function NEXTCHAR actually uses this buffer. READLINE is called whenever the buffer is wholly utilized.
5. PRINTLINE : is a procedure which dumps the above mentioned buffer. It also prints the errors of the current line (buffer) if there are any.

Figure 2.3 gives the transition diagram of whole PASCAL symbols.

SCAN uses the following data structure to implement arcs for single character delimiters.

```
var CHRCLASS : array (.CHAR.) of TOKENS;
```

It will be initialized as

```
CHRCLASS ['['] .CLASS := LBRACKET ;
CHRCLASS ['@'] .CLASS := ARROW ;
CHRCLASS [')'] .CLASS := RPARENT ;
CHRCLASS [']'] .CLASS := RBRACKET ;
CHRCLASS [','] .CLASS := COMMA ;
CHRCLASS [':'] .CLASS := COLON ;
CHRCLASS [';'] .CLASS := SEMICOLON ;

WITH CHRCLASS ['+'] DO BEGIN CLASS := ADDOP ; OP := PLUS ; END ;
WITH CHRCLASS ['-'] DO BEGIN CLASS := ADDOP ; OP := MINUS ; END ;
WITH CHRCLASS ['/'] DO BEGIN CLASS := MULOP ; OP := RDIV ; END ;
WITH CHRCLASS ['*'] DO BEGIN CLASS := MULOP ; OP := ASTR ; END ;
WITH CHRCLASS [=] DO BEGIN CLASS := RELOP ; OP := EQOP ; END ;
```

TABLE 2.2- Algorithm of the Lexical Analyzer

```

PROCEDURE SCAN ;
  LABEL 1 ;

  PROCEDURE PRINTLINE ; BEGIN ... END ;
  PROCEDURE READLINE ;
  BEGIN IF EOF(INPUT) THEN
    TOKEN.CLASS := EOFPGM ; (* SO RETURN FROM SCAN *)
  END ;
  FUNCTION NEXTCHAR : CHAR ; BEGIN ... END ;

BEGIN (* SCAN *)
  L : WHILE CH = ' ' DO
    CH := NEXTCHAR ; (* IGNORE BLANKS *)

  CASE CH OF
    'A'..'Z' : (* IDENTIFIERS *)
    '0'..'9' : (* NUMERIC CONSTANTS *)
    '<' : (* FOR THESE DELIMITERS FOLLOW *)
    '>' : (* THE ACTIONS OF CORRESPONDING *)
    '(' : (* ARCS SPECIFIED IN PASCAL *)
    ')' : (* TRANSITION DIAGRAM. *)
    '(' : (* THIS CASE LABEL IS PERFORMED TO SHOW
          HOW A LABEL WILL BE TREATED IN GENERAL *)

    BEGIN
      CH := NEXTCHAR ;
      IF CH = '*' THEN (* COMMENTS *)
        BEGIN
          REPEAT
            UNTIL NEXTCHAR = '*' ;
          UNTIL NEXTCHAR = ')' ;
          CH := NEXTCHAR ; GOTO L
        END
      ELSE IF CH = '.' THEN
        BEGIN
          TOKEN.CLASS := LBRACKET ;
          CH := NEXTCHAR
        END
      ELSE TOKEN.CLASS := LPARENT
    END ; (* CASE LABEL '(' *)

  OTHERWISE
    BEGIN
      TOKEN.CLASS := CHRCLASS [CH] ;
      IF TOKEN.CLASS = OTHERSY THEN
        BEGIN
          ERROR (6) ; (* ILLEGAL SYMBOL *)
          CH := NEXTCHAR ;
          GOTO L
        END
      END
    END
  END ; (* CASE SCAN *)
END ;

```

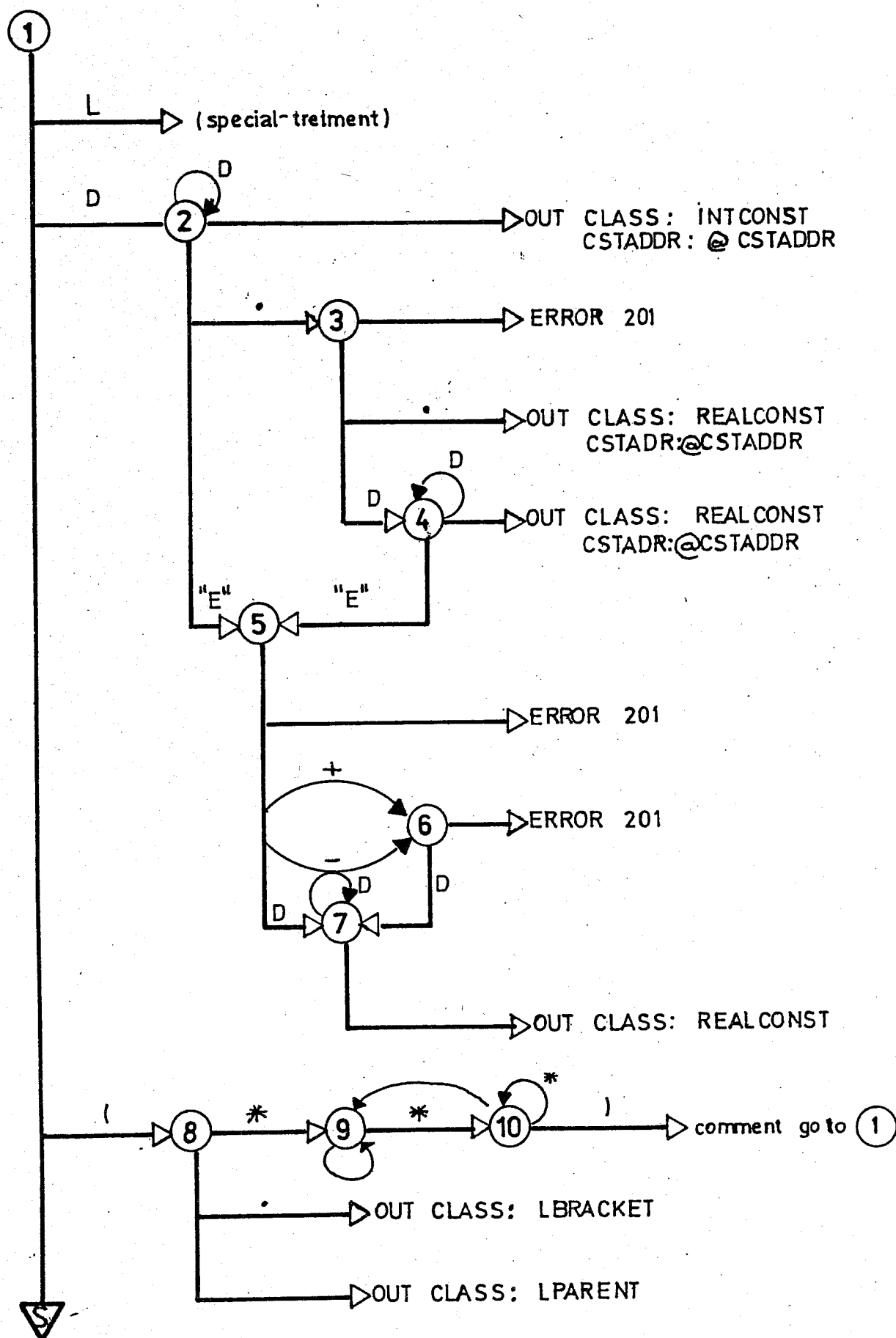


Figure 2.3- State Diagram for PASCAL.

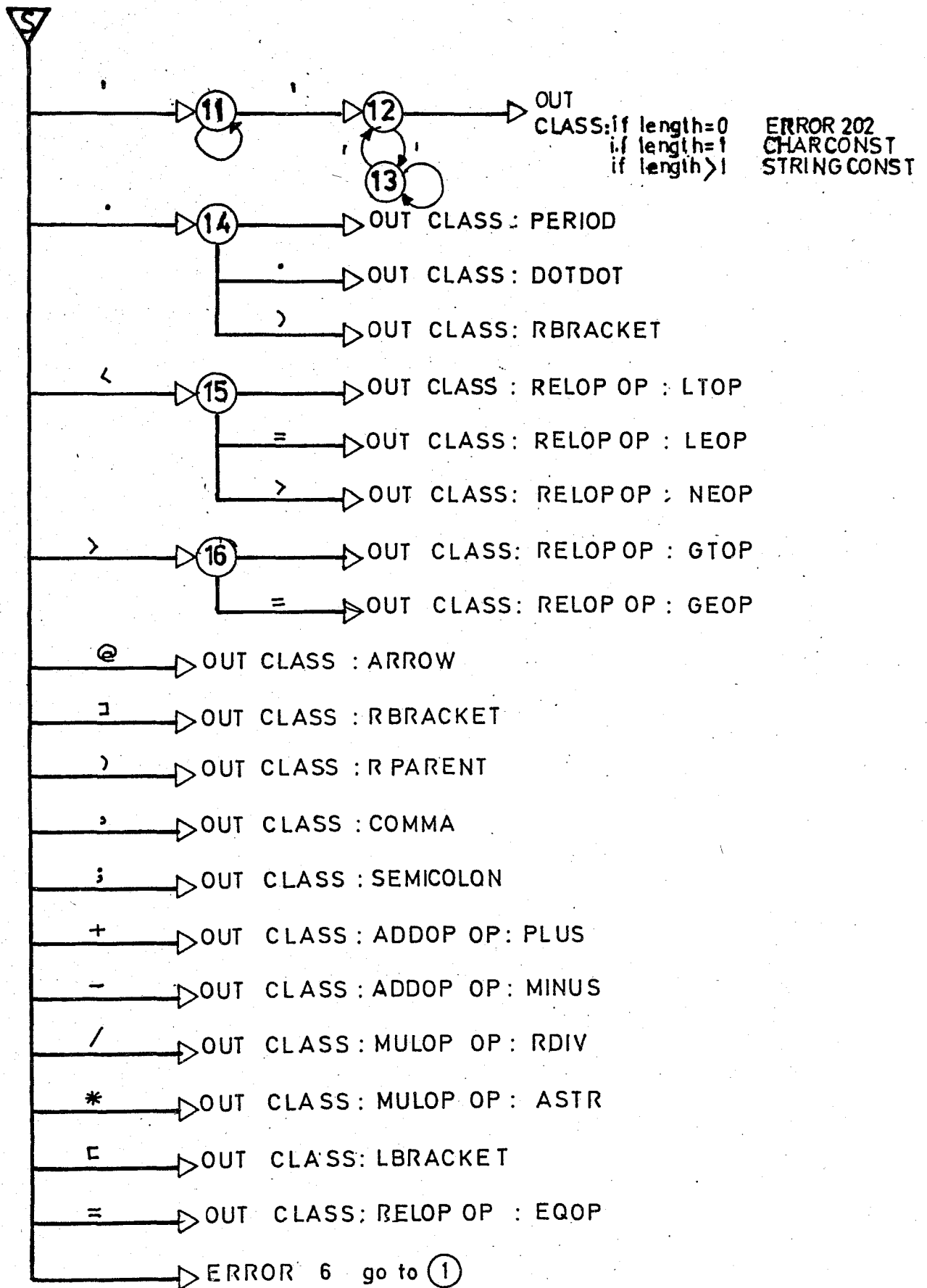


Figure 2.3- Continued.

All other characters will have a CLASS field of value OTHERSY. These are characters that have not been included in PASCAL character set, so if met an error message should be given.

Also note that a class code, EOFPGM is included within the user defined scalar type SYMBOL, which indicates there are no more symbols to the syntactic analyzer.

The procedure SCAN has the algorithm of table 2.2 (Assume result is returned in the global variable TOKEN).

Treatment of case labels in algorithm of table 2.2 are easy and straightforward. That is, once starting character of a symbol is found, it is the task of the corresponding case label to complete the syntax of the symbol. (Note that, SCAN assures that when it returns, the next character will always be scanned).

Only the case label related to identifiers, is specially treated. Identifiers and keywords both have the same syntax. So, the keywords of the language could be initially classified as identifiers and the correct symbol asserted after consulting a table of keywords. Due to the additional table search, this method is slower than the direct recognition of keywords through state transition tables.

Thus assuming we have two keywords ENDE, ELSE the arc labelled L in figure 2.3 will be modified to include keywords ENDE, ELSE (figure 2.4. Note that keyword ENDE used in this figure is not a PASCAL keyword).

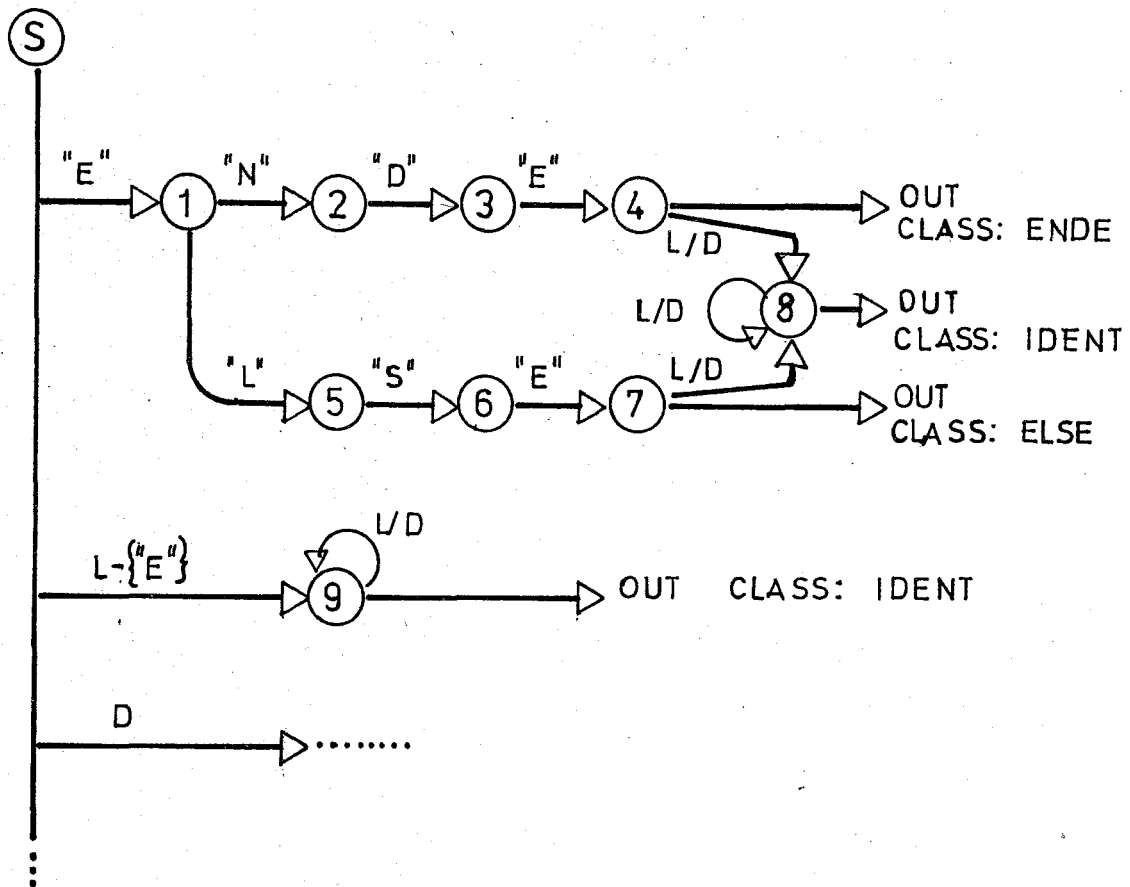


Figure 2.4- State Diagram to Recognize Keywords.

In figure 2.4 symbols L,D represent the set of letters and digits respectively as in fig. 2.3. The label L/D means letter or digit and L-{"E"} denotes the set of letters excluding the character "E". This diagram actually contains all PASCAL keywords but for simplicity I assumed only the existence of two keywords.

A set of tables is used to represent this transition diagram. The set consists of four arrays indexed by state numbers. These arrays are declared as;

```

var  DEFAULT      : array [0..171] of TOKENS;
     NEXT, CHECK  : array [0..168] of INTEGER;
     BASE         : array [0..171] of INTEGER;
  
```

BASE array is used to determine the base location of the entries for each state stored in NEXT and CHECK arrays.

DEFAULT array indicates the class of symbol fetched.

Assuming we are at case label 'A'..'Z' in the procedure SCAN, (thus CH contains a letter initially) table 2.3 gives the algorithm used (BAZ and S are variables declared of type integer).

To compute the transition for state S on input character CH, the pair of arrays NEXT and CHECK is first consulted. In particular, the algorithm finds their entries for state S in location $BAZ := BASE [S] + ORD(CH)$. $NEXT [BAZ]$ is taken to be the next state for S on input CH if $CHECK [BAZ] = S$. This procedure is repeated within a loop until $CHECK [BAZ] \neq S$ for some state S.

When we exit from the loop, the DEFAULT entry corresponding to the last state S, shows whether an identifier or a keyword (if so which one) is met.

For example, suppose state 1 in figure 2.4 is indexed 5 in numbering states (that is $S = 5$ initially). Then a value is chosen for $BASE [5]$ and value 5 is entered into $CHECK [BASE [5] + ORD ('N')]$ and $CHECK [BASE [6] + ORD ('L')]$. The next states on 'N' and 'L' for 5, are entered into the corresponding entries of the NEXT array.

TABLE 2.3- Algorithm to Compute Transitions

```

S := ORD(CH) ; (* INITIAL STATE FOR BASE ARRAY *)
CH := NEXTCHAR ;
BAZ := BASE [S] + ORD (CH) ;

WHILE S := CHECK [BAZ] DO BEGIN
  S := NEXT [BAZ] ;
  CH := NEXTCHAR ;
  BAZ := BASE [S] + ORD (CH) ;
END ; (* WHILE *)

TOKEN := DEFAULT [S] ;
IF CH IN [ 'A'..'Z', '0'..'9' ] THEN
  TOKEN.CLASS := IDENT ;

(* ... RETURN ... *)

```

DEFAULT [5]. CLASS will be IDENT. The algorithm of Table 2.3 will make the right transitions on "N" and "L", otherwise loop will be terminated. Therefore, if input CH is any character, but "N" or "L", we shall not find $CHECK [BASE[5] + ORD(CH)] = 5$ and exit from the loop and token will be DEFAULT 5, which is an IDENT.

Only, the DEFAULT entries corresponding to states 4 and 7 in figure 2.4 will contain classes representing keywords ENDE and ELSE respectively.

BASE values are initialized so that BAZ values in table 2.3 for different S values, do not conflict with the existing CHECK entries. Sizes of NEXT and CHECK arrays are highly dependent on these BASE values, so that they must be chosen carefully.

Assume that the input string is ELSEA. According to the algorithm we exit from the loop with DEFAULT class for the keyword ELSE. The last if statement in table 2.3 used to correct such errors. That is, as soon as we exit from the loop the last character scanned is tested. If it is a letter or digit, that means we have met an identifier, not a keyword.

Algorithm of table 2.3 is also used in recognition of operators such as mod, and, or, div etc. in addition to keywords and identifiers.

The transition tables for the diagram in figure 2.4 is shown in figure 2.5. Here it is assumed that ordinal values of letters range from 1 to 26, in alphabetical order. That is $ORD('A') = 1$, $ORD('B') = 2, \dots, ORD('Z') = 26$, and so on.

Base values are all zero except for entry 31. This is because when S is 31, letter "E" is expected to recognize the keyword ENDE and if BASE 31 were 0, baz value $(BASE[31] + ORD('E'))$ would correspond to CHECK entry 5, which have already been utilized.

	DETAULT	BASE		NEXT	CHECK
1	IDENT	0	1	0	0
2	IDENT	0	2	0	0
3	IDENT	0	3	0	0
4	IDENT	0	4	31	30
5	IDENT	0	5	29	28
⋮	IDENT	0	6	32	31
26	IDENT	0			
27	IDENT	0	12	27	5
28	IDENT	0	13	0	0
29	ELSE	0	14	30	5
30	IDENT	0	18	0	0
31	IDENT	1	19	28	27
32	ENDE	0	20	0	0

Figure 2.5- Transition Tables to Recognize Keywords.

E. Lexical Time Errors

SCAN is responsible of detection and reporting lexical time errors. Lexical time errors of PASCAL are defined as follows;

- ERROR 6 : Illegal symbol
- ERROR 201 : error in real const; digit expected
- ERROR 202 : integer constant exceeds range
- ERROR 205 : null string not allowed
- ERROR 206 : integer part of real constant exceeds range.

The detection points of ERRORS 6,201,205 are shown in the transition diagram of fig. 3.3. Errors 202 and 206 can be detected during conversion of numbers from character code to binary form.

F. Conclusion

SCAN is designed as a hand-written lexical analyzer. For this algorithm GRIES (2) can be referred. The set of transition tables used for identifiers and keywords are explained in AHO-ULLMAN(1).

SCAN can work independent of the machine it is run with simple modifications. It is designed to work in FIELDATA character set, because the transition tables used for the recognition of identifiers and keywords are initialized using FIELDATA character codes. But the algorithm can be made to work in any character set, by addition of a single array.

Conversion algorithms from character code to binary form (in case of numbers) are general but errors 202 and 206 depend on the word size of the machine.

III. SYNTACTIC ANALYSIS

A. Definition

A syntactic analyzer (or PARSER) for a grammar G is a program that takes as input a string W and produces as output either a parse tree for it, if W is a sentence of G , or an error message indicating W is not a sentence of G . In case of translators, W corresponds to programs of a programming language whose grammar is defined by G .

Parsing algorithms can be classified into two categories bottom-up and top-down. The terms refer to the way in which parse trees are built. A bottom-up parser builds parse trees from the bottom (terminal nodes) to the top (root node). A top-down parser builds parse trees starting from the root node and works down to terminal nodes. In both cases the input to the parser is being scanned from left to right, one symbol at a time.

B. Bottom-Up Parsing

The bottom-up technique is to start at the string itself and try to reduce it to the distinguished symbol. Consider the sentence 35 of the following grammar for integers.

$$\begin{aligned} \langle N \rangle & ::= \langle D \rangle \mid \langle N \rangle \langle D \rangle \\ \langle D \rangle & ::= 0 \mid 1 \mid 2 \mid \dots \mid 9 \end{aligned}$$

The first step is to reduce the 3 to $\langle D \rangle$, yielding the sentential form $\langle D \rangle 5$. Thus the direct derivation $\langle D \rangle 5 \Rightarrow 35$

is constructed as shown in figure 3.1a. Next step is to reduce $\langle D \rangle$ to $\langle N \rangle$ (figure 3.1b). This proceeds until the last tree (figure 3.1d) has been formed. Note that, in such a parse, at each step a handle (leftmost simple phrase) of the sentential form is reduced.

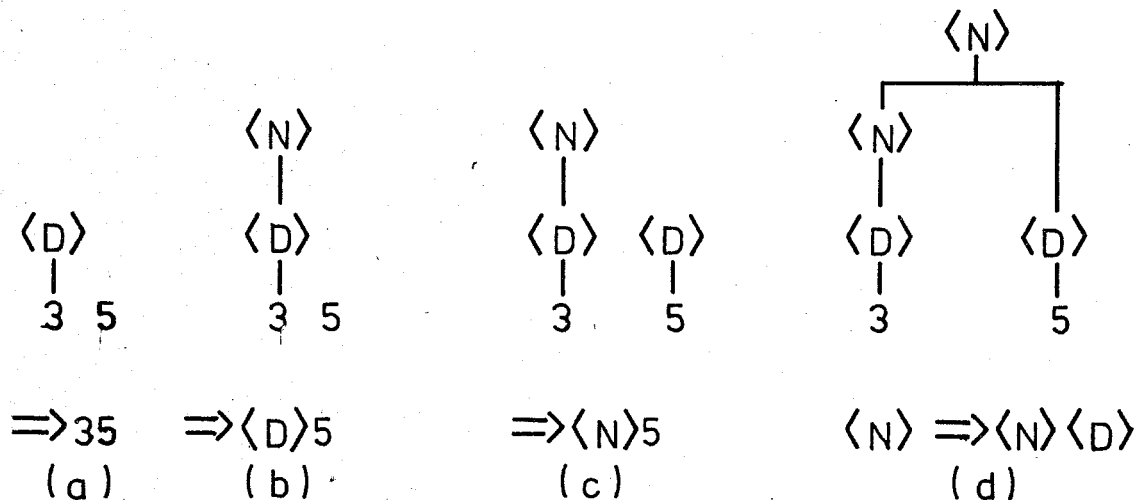


Figure 3.1- Bottom-Up Parse and Derivation of its Constructs.

Bottom-up parsers usually use a stack for reduction process. Symbols are pushed on to the stack until the right side of a production appears on top of the stack. The right side may then be replaced by (reduced to) the symbol on the left side of the production, and the process is repeated.

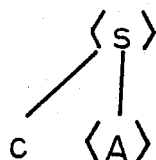
C. Top Down Parsing

Top down parsing can be viewed as an attempt to find a left most derivation for an input string. For example, consider the grammar,

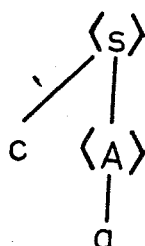
- 1) $\langle S \rangle ::= c \langle A \rangle d$
- 2) $\langle A \rangle ::= b | a$

and let input be $W = cad$. We initially start with the root node $\langle S \rangle$. First symbol of W matches the first symbol (which

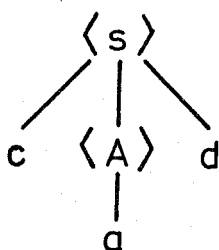
is the terminal c of rule 1. So we advance to the next input symbol, a and try to match it with the successor of terminal c in rule 1. (which is $\langle A \rangle$). Since a non-terminal is reached, we must first expand it before proceeding. Upto here we have constructed the following partial tree.



$\langle A \rangle$ has two alternatives, among which second one matches the input symbol, that is terminal a of rule 2. Thus $\langle A \rangle$ is expanded using the second alternate to obtain the partial tree



We now consider d , third input symbol, and successor of non terminal $\langle A \rangle$ of rule 1 (since its expansion has been completed) which is the terminal d , and this matches with the input symbol.



Since we have now produced a complete parse tree for W , we halt and announce successful completion of parsing.

Parsing algorithm used in this thesis, is a top-down technique. The algorithm is implemented as a PASCAL procedure named PARSE and it will be explained within the rest of this section.

D. Implementation

This section describes the parsing algorithm, and representation of the grammar. It also describes the problems of TOP-DOWN parsing and how they are solved by PARSE.

1. Graphic Notation of Grammars

The algorithm is centered upon the relationship between an element and its successors, and the element and its alternates. To show these relationships schematically, a graphic notation is introduced. Consider the following simple rule;

$$\langle \text{left hand side} \rangle ::= \langle p \rangle \langle q \rangle \mid \langle r \rangle \langle s \rangle$$

This rule can be represented as in figure 3.2. Arcs labelled a show alternates, and arcs labelled s show successors. Finally arc labelled d means "defined as".

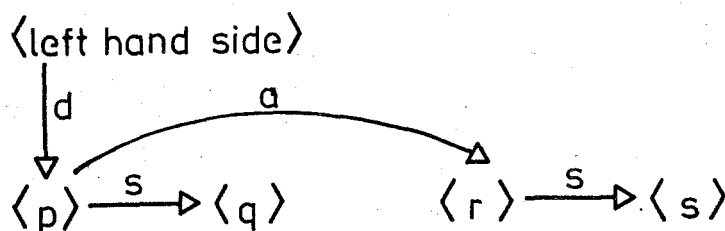


Figure 3.2- Graphic Notation of a Rule.

Tables used by PARSE to represent the grammar in core, use such alternate and successor links. So the graphic notation helps us in designing those tables.

If there is no alternate for an element, we indicate this by the absence of an alternate arc in the graph. While representing in core, alternate link of such an element will contain a dummy constant denoted by the constant identifier FAIL. Trying to access an alternate link whose value is FAIL

causes an error. Similarly to indicate end of successors, we use another dummy constant, denoted by the constant identifier OK.

2. Top Down Problems and Their Solutions

a) Direct left recursion

Assume we have a rule like $\langle x \rangle ::= \langle x \rangle \dots$. Then our first action to expand $\langle x \rangle$ would be to expand $\langle x \rangle$. And next action would also be to expand $\langle x \rangle$, since $\langle x \rangle$ is always the first rule of the expansion of $\langle x \rangle$. Thus we would be in a deadlock creating a loop around $\langle x \rangle$.

Best way to get rid of direct left recursion is to write the rules using iterative notation. Consider the following rule

$$\langle E \rangle ::= \langle E \rangle + \langle T \rangle \mid \langle T \rangle$$

It can be written as;

$$\langle E \rangle ::= \langle T \rangle \{ + \langle T \rangle \}_N$$

where N corresponds to the minimum number of iterations which could be zero. Iteration can be represented in successor alternate graphs easily. In this case, the last element in a $\{ \}$ pair points to the first element in the $\{ \}$ pair, as its successor. And the first element's alternate link will contain OK value. OK value used in an alternate link shows optionality and means end of the expansion of the current non-terminal just as an OK value met in a successor link.

Figure 3.3 shows the graphical representation of the rule $\langle E \rangle ::= \langle T \rangle \{ + \langle T \rangle \}_N$, when N is 0.

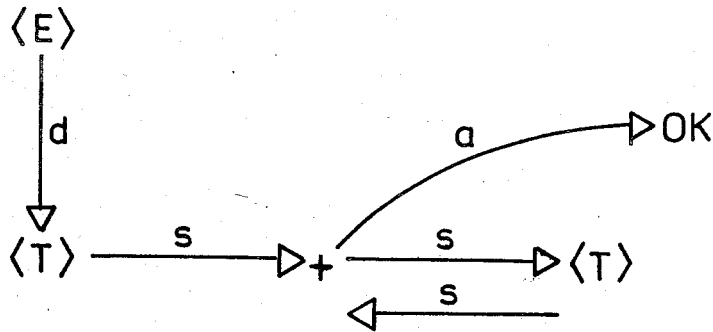


Figure 3.3- Graph of an Iterated Rule (N=0).

If N is greater than 0, the graph gets larger. But N is usually 0, and there are very few rules requiring N greater than 0. Figure 3.4 shows the same graph when N is 1.

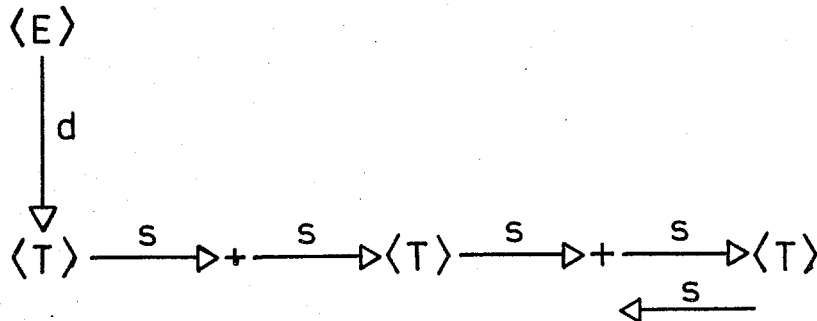


Figure 3.4- Graph of an Iterated Rule (N=1).

b) General left recursion

Assume we have rules like $\langle U \rangle ::= \langle V \rangle x$ and $\langle V \rangle ::= \langle U \rangle y$. These rules yield $\langle U \rangle ::= \langle U \rangle yx$. There is no simple way of getting rid of such rules, except for manually checking.

c) Backups

Backup conditions rise when we are expanding non-terminals. That is, when we met a non-terminal we first go to its definition, and try to parse it. And if its definition also contains non-terminals in its definition we proceed similarly until success has been reached at some level. If

during this process, a failure occurs at any level, we must back-up to a higher level, to try the alternate of the non-terminal that caused the failure. Backup process may continue upto the initial level where we first met a non-terminal.

Backups should be avoided because;

- 1) It is time consuming and inefficient
- 2) If semantics are being performed as each syntactic element is identified, then they have to be undone
- 3) If code generation is being done, the code generated needs to be erased upto a point
- 4) Backups make error recovery very hard.

Backups can be avoided if we impose the restriction, NO NON-TERMINAL CAN HAVE AN ALTERNATE LINK DIFFERENT FROM FAIL VALUE, on the grammar. Thus in case of match fails, alternate link will be referred. If an alternate element exists we proceed with the alternate link. If its value is OK (which means expansion is completed) we move a higher level. But if its value is FAIL, an ERROR routine will be called.

Consider the graph in figure 3.5a. The definition of <repetitive-stmt> does not obey the above stated restriction, since the non-terminal <while-stmt> has an alternate link. Restriction is obeyed, by substituting the definition of <while-stmt> into the definition of <repetitive-stmt> (figure 3.5b).

These sort of substitution transformations on the grammar may lead to very long rules. But it is worth to do so, since only by this way backup problem can be solved

d) Representing empty strings

Consider a rule like $\langle x \rangle ::= a|b|c|d|e|\epsilon$ where ϵ denotes the empty string. The presence of an empty string can be represented graphically by setting the alternate link of last element to OK value (figure 3.6).

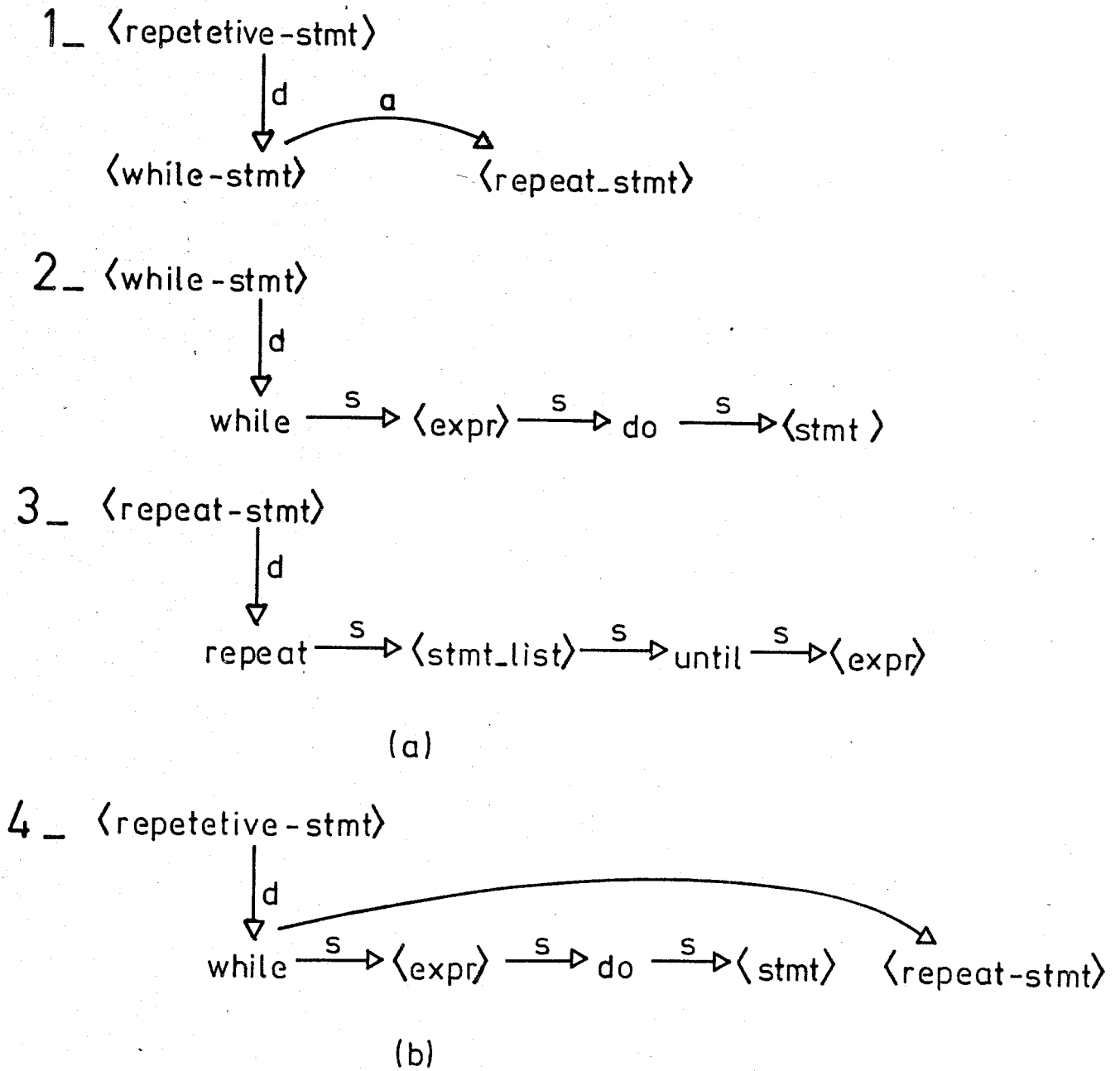


Figure 3.5- An example of Substitution Transformation on a graph.

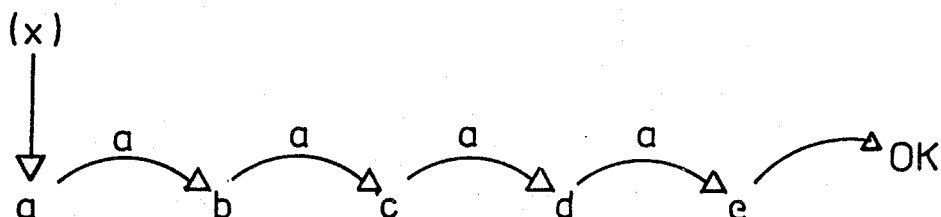


Figure 3.6- Graph of a Rule Containing an Empty String.

Recall that OK value in an alternate link, was used to indicate end of expansion while discussing iterative notation. OK value is used similarly here. Thus, if the alternate link of the terminal has been tried, that means expansion is completed with the empty string.

Representations of empty strings in this way may be time consuming, since all alternate links up to the last terminal should be tried. There is no simple way to get rid of empty strings and sometimes they are unavoidable. Yet some tricks can be used to get rid of them. For example consider the following rules.

$$\begin{aligned} \langle A \rangle & ::= [\langle B \rangle] \\ \langle B \rangle & ::= c \mid \langle D \rangle \\ \langle D \rangle & ::= e \mid \varepsilon \end{aligned}$$

These rules can be modified as follows;

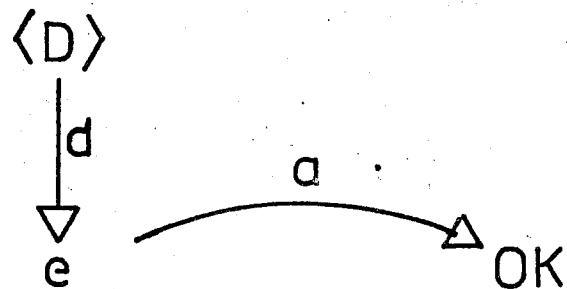
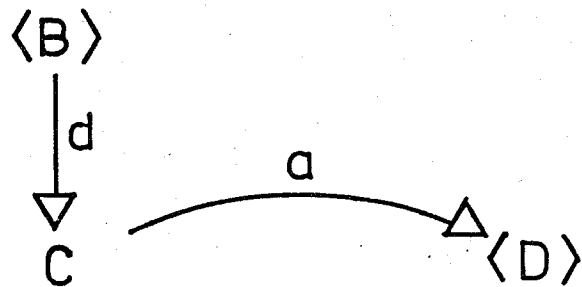
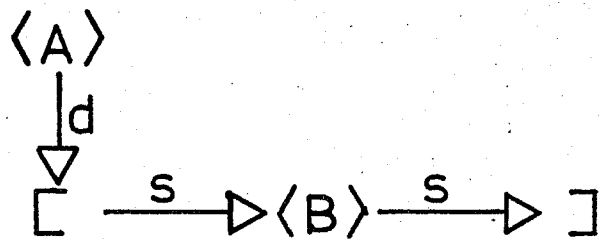
$$\begin{aligned} \langle A \rangle & ::= [\langle B \rangle \\ \langle B \rangle & ::=] \mid c \mid \langle D \rangle] \\ \langle D \rangle & ::= e \end{aligned}$$

That would be easier and faster to parse. Graphical notations of both grammars are given in figure 3.7a and figure 3.7b in that order. Let input be, $w = []$. In figure 3.7a when parsing this input, to see $\langle B \rangle$ is empty we have to move down to the definition of $\langle D \rangle$. Then we check alternate link of the terminal e which is OK. Thus we move up to the first level to match $]$.

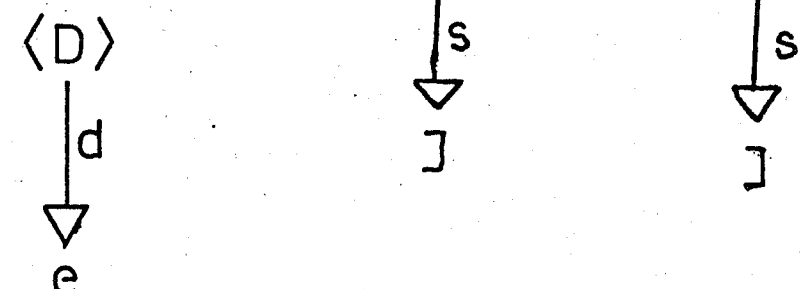
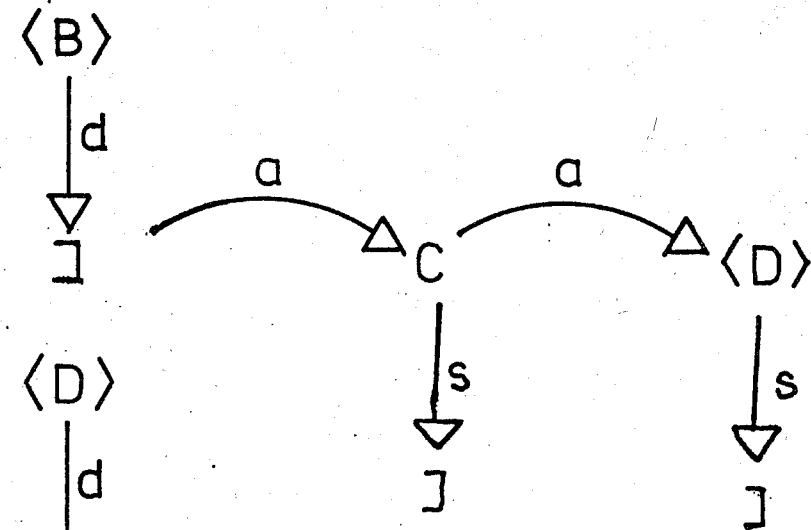
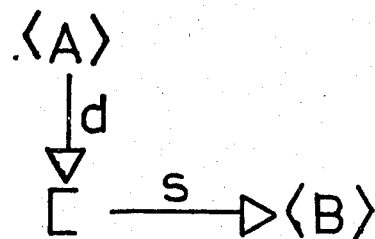
In case of figure 3.7b, for the same input ($w=[]$) it is the second level and first element where a match is found.

3. Data Structures to Represent the Grammar In Core

PARSE uses two static tables to represent the grammar in core.



(a)



(b)

Figure 3.7. An Example to Get Rid of Empty Strings

i) SYNTAX TYPE TABLE (TTABLE)

TTABLE has an entry for each element of the vocabulary
It is declared as;

```

type TPRANGE = 1..TPSIZE ; {#elements in vocabulary}
var TTABLE : array [TPRANGE] of
    record
        case TERMINAL : BOOLEAN of
            TRUE : (CLASS : SYMBOL )
            FALSE : (POINTER : SPRANGE) ;
        end;

```

Tagfield TERMINAL is set to TRUE, if the corresponding element is a terminal and it is set to FALSE if the corresponding element is a non-terminal. Field POINTER is used for non-terminals, and points to an entry in the structure table (second table) where we record the definition of each non-terminal. If the graph notation is considered, this variant corresponds to the arc labelled d.

On the other hand, field CLASS is defined for terminals and is used to communicate with SCAN. In other words, this entry records the internal code of the symbol required.

ii) SYNTAX STRUCTURE TABLE (STABLE)

This table is used represent the rules of the grammar. Thus it has an entry for each symbol of a production, and for all productions. The table is declared as;

```

type SPRANGE = -1.. SPSIZE; {size of grammar}
const FAIL = -1; OK = 0;
var STABLE : array [SPRANGE] of
    record
        SUCCESSOR, ALTERNATE : SPRANGE;
        TPTYPE : TPRANGE
    end;

```

Fields SUCCESSOR and ALTERNATE are used as internal pointers in STABLE and they correspond to arcs labelled s and a in graph notation. As discussed before, ALTERNATE field can take FAIL and OK values, and SUCCESSOR field can take OK value.

So subrange type SPRANGE has a lower bound equal to -1 . Field TPTYPE is a backward pointer to TTABLE. It is used to get information about the current element being tested.

PARSE also uses a stack to keep completed parts of the parse tree being constructed. Two procedures (POP and PUSH) are used to manipulate this stack. An entry is allocated each time a non-terminal is met and deallocated each time the expansion of a non-terminal is completed.

4. Parsing Algorithm

Table 3.1 gives the algorithm of PARSE. Variable SPINX is used to index STABLE. It is initially set to 1, that is, it initially points to the first rule of the grammar. SPINX values are pushed on to the stack as long as they index non-terminal entries.

When a terminal is reached in the definition, it is compared with the input symbol. If they match procedure SUCCESS is called. Otherwise ALTERNATE field is tested. If its value is FAIL then an ERROR routine is called. If its value is OK then, this means end of expansion and again procedure SUCCESS is called. If none of these are true then SPINX is set to the value of ALTERNATE field.

Procedure SUCCESS is called in two cases. Its formal parameter MODE is set to 0, if a match occurs and is set to 1 if value of ALTERNATE link is OK. Function of SUCCESS is to advance SPINX to the successor of current element. It also tests whether expansion is completed by testing $MODE = 1$ or successor of the current element is OK. If so, it moves up a

TABLE 3.1- Parsing Algorithm

```

PROCEDURE PARSE ;
  PROCEDURE POP (ITEM : SPINX) ; BEGIN END ;
  FUNCTION NONTERMINAL : BOOLEAN ; BEGIN END ;
  PROCEDURE SUCCESS (MODE : INTEGER) ;
  BEGIN
    WHILE (STABLE [SPINX].SUCCESSOR=OK) OR (MODE=1) DO
      BEGIN
        MODE := 0 ;
        IF STACKEMPTY THEN
          IF TOKEN.CLASS = EOPFGM THEN
            BEGIN
              WRITEFLN ('ACCEPTED.') ;
              HALT
            END
          ELSE
            BEGIN
              WRITELN ('**END OF ANALYSIS.**') ;
              WRITELN ('..NO MORE SYMBOLS ACCEPTED..') ;
              HALT
            END ;
          SPINX := STACK [STACKTOP] ;
          POP (* MOVE UP A LEVEL *)
        END ; (* WHILE *)
        SPINX := STABLE [SPINX].SUCCESSOR
      END ; (* SUCCESS *)
    BEGIN (* PARSE *)
      SPINX := 1 ; (* INITIALIZE ROOT NODE *)
      WHILE TRUE DO
        WITH STABLE [SPINX] DO BEGIN
          WHILE NONTERMINAL (TPTYPE) DO BEGIN
            PUSH (SPINX) ;
            SPINX := ITABLE [TPTYPE].POINTER
          END ;
          IF ITABLE [TPTYPE].CLASS = TOKEN.CLASS
            THEN (* CURRENT ELEMENT = INPUT SYMBOL *)
              BEGIN
                SCAN : (* GET NEXT TOKEN *)
                SUCCESS (0)
              END
            ELSE IF ALTERNATE = FAIL THEN
              BEGIN ERROR ; RECOVERY END
            ELSE IF ALTERNATE = OK THEN
              SUCCESS (1)
            ELSE SPINX := ALTERNATE
          END
        END (* FOREVER LOOP *)
      END
    END
  END

```

level by a POP operation on stack.

Function NONTERMINAL returns true if its argument points to a non-terminal TTABLE entry. STACK-EMPTY is assumed to be a global variable set true by the procedure pop when the stack becomes empty.

Table 3.2a gives an example grammar, where as table 3.2b and 3.2c gives the internal representation of this grammar in core. A trace of the algorithm, for this grammar using input $W = A > B * C \neq$ is given in table 3.3. The input symbol \neq , is used to indicate end of string. In this table, existence of two functions SUCC and ALT are assumed which return successor and alternate values of their parameters.

E. Conclusion

Procedure PARSE is a predictive parse technique. Predictive parsers are efficient ways of implementing recursive descent parsing by handling the stack of activation records explicitly, and by keeping a set of tables to represent the grammar in core.

PARSE uses two static tables (TTABLE and STABLE) for grammar representation. Both tables have the prefix PACKED which requests a compact storage representation for values of the tables. Sizes of these tables are 97 words and 544 words respectively. Thus the space they occupy is not much.

The complete syntax of PASCAL written in BNF obeying the restrictions stated in this section is given in appendix C.

Various parsing techniques are explained in detail in GRIES (2) and AHO-ULLMAN (1).

TABLE 3.2- Parse Tables for an Example Grammar

GRAMMAR:

- 1) <expression> ::= <simple exp> RELOP <simple exp>
- 2) <simple exp> ::= <term> ADDOP <term>
- 3) <term> ::= <factor> MULOP <factor>
- 4) <factor> ::= IDENT | (<expression>)

(a)

<u>INDEX</u>	<u>EXPLANATION</u>	<u>TERMINAL</u>	<u>POINTER</u>	<u>CLASS</u>
1	<expression>	FALSE	1	
2	<simple exp>	FALSE	4	
3	<term>	FALSE	7	
4	<factor>	FALSE	10	
5	RELOP	TRUE		RELOP
6	ADDOP	TRUE		ADDOP
7	MULOP	TRUE		MULOP
8	IDENT	TRUE		IDENT
9	(TRUE		LPARENT
10)	TRUE		RPARENT
11	≠	TRUE		EOFPGM

Contents of TTABLE

(b)

TABLE 3.2- Continued.

<u>INDEX</u>	<u>EXPLANATION</u>	<u>SUCCESSOR</u>	<u>ALTERNATE</u>	<u>TPTYPE</u>
EXPRESSION				
1	<simple exp>	2	FAIL	2
2	RELOP	3	OK	5
3	<simple exp>	2	FAIL	2
SIMPLE EXPRESSION				
4	<term>	5	FAIL	3
5	ADDOP	6	OK	6
6	<term>	5	FAIL	3
TERM				
7	<factor>	8	FAIL	4
8	MULOP	9	OK	7
9	<factor>	8	FAIL	4
FACTOR				
10	IDENT	OK	11	8
11	(12	FAIL	9
12	<expression>	13	FAIL	1
13)	OK	FAIL	10

Contents of STABLE

(c)

TABLE 3.3- Continued

STEP#	SPINX	IS TERMINAL?	OPERATION	STACK
11.	10	Yes	<p style="text-align: center;">IDENT IDENT</p> 1. Current element $\stackrel{?}{=}$ input symbol 2. Yes 2.1. SCAN input symbol := MULOP 2.2. SUCCESS(0) 3. In SUCCESS with MODE = 0 3.1. (SUCC(10) $\stackrel{?}{=}$ OK) OR (MODE $\stackrel{?}{=}$ 1) 3.2. Yes, 3.2.1. STACK not EMPTY, POP 3.2.2. SPINX := 7 3.3. SPINX := SUCC(7) = 8 3.4. RETURN	→ 4 3
12.	8	Yes	<p style="text-align: center;">MULOP MULOP</p> 1. Current element $\stackrel{?}{=}$ input symbol 2. Yes, 2.1. SCAN {input symbol := IDENT} 2.2. SUCCESS(0) 3. In SUCCESS with MODE = 0 3.1. (SUCC(8) $\stackrel{?}{=}$ OK) OR (MODE $\stackrel{?}{=}$ 1) 3.2. No, SPINX := SUCC(8) = 9 3.3. RETURN	
13.	9	No	1. PUSH (SPINJ) 2. SPINX := 10	→ 9 4 3
14.	10	Yes	<p style="text-align: center;">IDENT IDENT</p> 1. Current elmt = input symbol 2. Yes, 2.1. SCAN {input symbol}:# 2.2. SUCCESS(0) 3. In SUCCESS with MODE = 0 3.1. (SUCC(10) $\stackrel{?}{=}$ OK) OR (MODE $\stackrel{?}{=}$ 1) 3.2. Yes, 3.2.1. Stack not empty, POP 3.2.2. SPINX := 9 3.3. SPINX := SUCC(9) = 8 3.4. RETURN	→ 4 3

TABLE 3.3- Continued

<u>STEP</u>	<u>SPINX</u>	<u>IS TERMINAL?</u>	<u>OPERATION</u>	<u>STACK</u>
15	8	Yes	<p style="text-align: center;">MULOP</p> <ol style="list-style-type: none"> 1. Current element \neq input symbol 2. No, ALT(8) \neq FAIL 3. No, ALT(8) \neq OK 4. Yes, SUCCESS(1) 5. In SUCCESS with MODE = 1 <ol style="list-style-type: none"> 5.1. (SUCC(8) \neq OK) OR (MODE \neq 1) 5.2. Yes, <ol style="list-style-type: none"> 5.2.1. Stack not empty, POP 5.2.2. SPINX := 4 5.3. SPINX := SUCC(4) = 5 5.4. RETURN 	<div style="border: 1px solid black; width: 40px; height: 40px; margin: 0 auto; display: flex; align-items: center; justify-content: center;"> → 3 </div>
16	5	Yes	<p style="text-align: center;">ADDOP</p> <ol style="list-style-type: none"> 1. Current element \neq input symbol 2. No, ALT(5) \neq FAIL, 3. No, ALT(5) \neq OK, 4. Yes, SUCCESS(1) 5. In success with MODE = 1 <ol style="list-style-type: none"> 5.1. (SUCC(5) \neq OK) OR (MODE \neq 1) 5.2. Yes, <ol style="list-style-type: none"> 5.2.1. STACK not empty, POP 5.2.2. SPINX := 3 5.3. (SUCC(3) \neq OK) OR (MODE \neq 1) 5.4. Yes, <ol style="list-style-type: none"> 5.4.1. STACK is empty, 5.4.2. Input symbol \neq 5.4.3. Yes, <ol style="list-style-type: none"> 5.4.3.1. PRINT('ACCEPTED'); 5.4.3.2. HALT 5.5. RETURN 	<div style="border: 1px solid black; width: 40px; height: 40px; margin: 0 auto; display: flex; align-items: center; justify-content: center;"> → </div>

IV. ERROR RECOVERY

A. Definition

Programs submitted to a translator often have errors of various kinds. A good translator, therefore, should find as many errors as possible.

While discussing parsing in the previous chapter, our purpose was to determine whether a sentence of a given language was accepted or not. It was assumed that detection of the first error stops parsing with an "UNACCEPTED SENTENCE" message. Yet, as stated here, even in the presence of errors, a translator should be able to continue parsing and scan the entire program trying to analyze all of it. The term "error recovery" is used for the process of determining how to continue analyzing a source program when an error is found.

Errors are classified as syntactic and semantic errors. Syntactic errors are those detectable by the lexical or syntactic phase of a translator. Other errors detectable by a translator are classified as semantic errors. In this thesis semantic errors are limited to errors of declaration. Lexical phase errors are outlined in section 3. Their recovery is simple and performed by SCAN. Recovery of semantic errors, here, is defined as to suppress extra error messages and is described in Section 5. So this section, describes the algorithm used to recover syntactic errors.

Recovery algorithm modifies the input, so the correct portions of the program can be pieced together and successfully processed.

B. Syntactic Errors

Typical of these are the following;

- . the insertion of an extraneous symbol
- . the deletion of a required symbol
- . the replacement of a correct symbol by an incorrect symbol.
- . the transposition of two adjacent symbols

Note that a replacement error and a transposition error can each be treated as special cases of an insertion error followed by a deletion error.

Here are a few common examples of syntactic errors.

1. Missing right parenthesis

```
MIN(A,2*(3+B);
```

2. Missing semicolon

```
A:=3
```

```
B:=4;
```

3. ":" in place of ";" or "=" in place of ":=

```
A:=3:
```

```
A = 3;
```

```
B:=4;
```

```
B:=4;
```

4. Misspelled keyword

```
PORCEDURE A;
```

5. Extra blank

```
(* COMMENT * )
```

(1) and (2) are examples of deletion errors, (3) replacement error (5) an insertion error and (4) a transposition error.

Quite frequently we can not detect that an error has occurred until long after it has taken place. For example, consider the PASCAL program fragment

```
FORI := 1+K[20] TO 100 DO L:=L+1;
```

The obvious error is a missing blank between the keyword FOR and the name I. No error is discernible, however, until the keyword TO has been read, since PARSE treats FORI:=1+K[20] as an assignment statement.

This example shows that the detection of an error may occur an arbitrarily long distance after the place where the error actually occurred.

Error recovery strategy tries to change the small portion of the program containing the error into a string that is legal, by making minimum number of insertions, deletions and symbol modifications necessary. Because of the distance problem, recovery algorithm may generate several error messages for a single error. If we do happen to generate a few, it doesn't really matter.

C. Recovery From Syntactic Errors

At any point of a parse of a source program, the program has the form

xTt

where x represents the part already processed, T is the next symbol to be scanned, and t is the rest of the program. Suppose an error occurs with T. In the TOP-DOWN method, this means that the partial tree built to cover x can not be extended to cover T.

At this point we must determine how to change the program to "fix" the error. It can be changed most easily in the following ways (or perhaps combinations of them.).

1. Delete T and try to parse again
2. Insert a string of terminals q between x and T (yielding xqTt) and begin parsing using the head of

qT . This insertion should allow us to process all of qT before another error occurs.

3. Delete some symbols from the tail of x .

Deleting part of x , we must change the semantic information accordingly, and this is not easy to do so. Methods (1) and (2) will be our main methods of recovering.

In figure 4.1 the incomplete branch named P corresponds to an application of the rule $\langle P \rangle ::= \langle A \rangle$; and ";" is the incomplete part of the branch. Similarly, the incomplete branch named E corresponds to an application of the rule $\langle E \rangle ::= \langle T \rangle \{ + \langle T \rangle \}$. To complete the branch, we need a single $\langle T \rangle$, followed by a number of "+" $\langle T \rangle$ "s. The incomplete part is therefore $\langle T \rangle \{ + \langle T \rangle \}$.

Incomplete parts of branches play a large role in error recovery; they tell us, in effect, what can or should appear next in the source program.

Now let us suppose that an error occurs during a parse; no partially constructed syntax tree can further be built. The following recovery algorithm is performed.

1. A list L of the symbols in the incomplete parts of the incomplete branches is constructed.
2. The head symbol of Tt is repeatedly examined and discarded (yielding a new string Tt) until a T is found such that $U \Rightarrow * \dots T \dots$ for some U in L (either $U=T$ or $U \Rightarrow + \dots T \dots$).
3. An incomplete branch which caused U of step 2 to be put in L is determined.
4. A terminal string q is determined such that if inserted just before T , the continuation of the parse will cause T to be correctly linked to the incomplete branch of step 3.
5. q is inserted just before T and the parse is continued,

beginning with the head symbol of q as the incoming symbol.

Consider, for example, the parse as indicated by fig. 4.1. An error has occurred with $)$ as the incoming symbol. We have $L = \{ \langle T \rangle, +, ; \}$. By step 2 we see that $\langle T \rangle \Rightarrow + \dots$. That is, $\langle T \rangle \Rightarrow \langle F \rangle \Rightarrow (\langle E \rangle)$. This notation can be expressed as $)$ is reachable from the nonterminal $\langle T \rangle$.

D. Implementation

1. Formal Definition of the Algorithm

I will illustrate the algorithm using the following grammar

$$\begin{aligned} \langle P \rangle &::= \langle A \rangle ; \\ \langle A \rangle &::= i := \langle E \rangle \\ \langle E \rangle &::= \langle T \rangle \{ + \langle T \rangle \}_o \\ \langle T \rangle &::= \langle F \rangle \{ * \langle F \rangle \}_o \\ \langle F \rangle &::= i | (\langle E \rangle) \end{aligned}$$

At any step of a parse, one or more syntax trees have been constructed, with some incomplete branches. An incomplete branch named U corresponds, to an application of a rule

$$U ::= x_1 x_2 \dots x_{i-1} x_i \dots x_n$$

where $x_1 \dots x_{i-1}$ is the completed part of the branch and $x_i \dots x_n$ the incomplete part of the branch. Figure 4.1 gives the tree for the sentence $i := i +) ;$ of the grammar, stated above. In this figure solid lines show a partially completed tree while the dotted lines show how the branches named $\langle P \rangle$ and $\langle E \rangle$ might be completed.

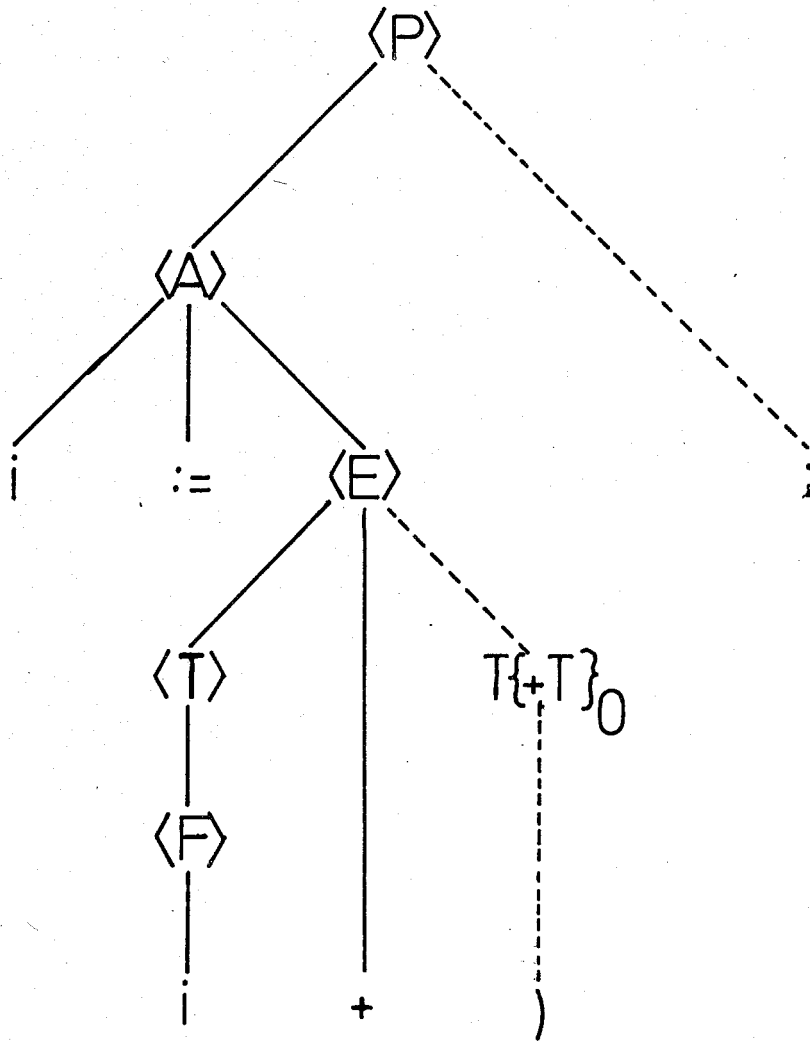


Figure 4.1- Top Down Tree Before Error Recovery.

The incomplete branch which caused $\langle T \rangle$ to be put in L is $\langle E \rangle ::= T\{+T\}_0$. We must therefore insert a string q to complete this branch. To cause the ")" to be associated with this branch, q must include $\langle E \rangle$. Since q will be a string of terminals $\langle E \rangle$ must be expanded. Simplest expansion of $\langle E \rangle$ is the symbol i . Therefore q will be $(i$. Thus, in step 5 qT becomes $(i$ beginning with head symbol "(" and this recovers the error (Figure 4.2).

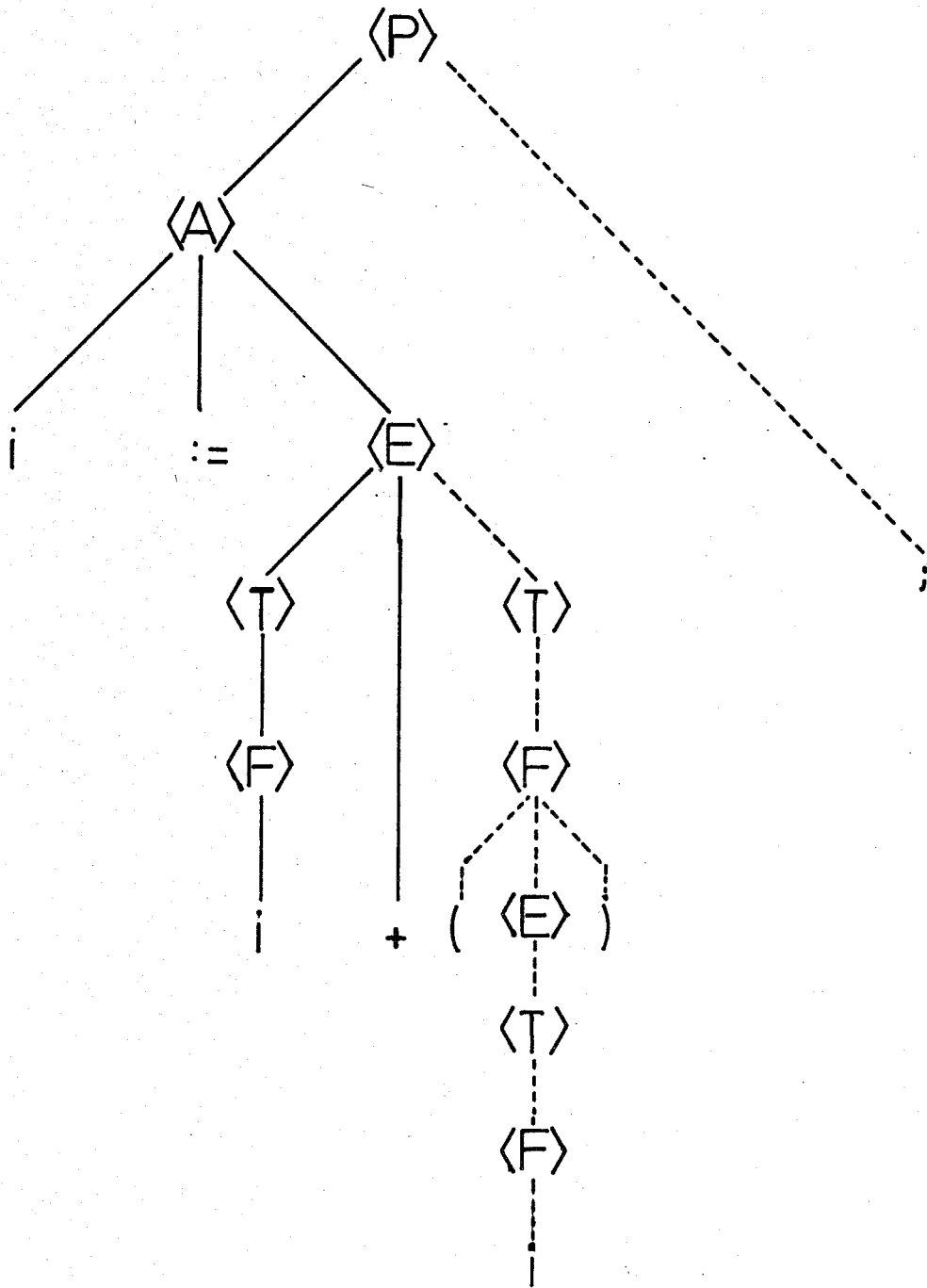


Figure 4.2. Top Down Free After Error Recovery

Figure 4.3. illustrates how the algorithm works with the input $i := (i+);$. In this example $L = \{;,), +, \langle T \rangle\}$. This time $"")"$ is immediately selected by step 2 ($U = ""$).

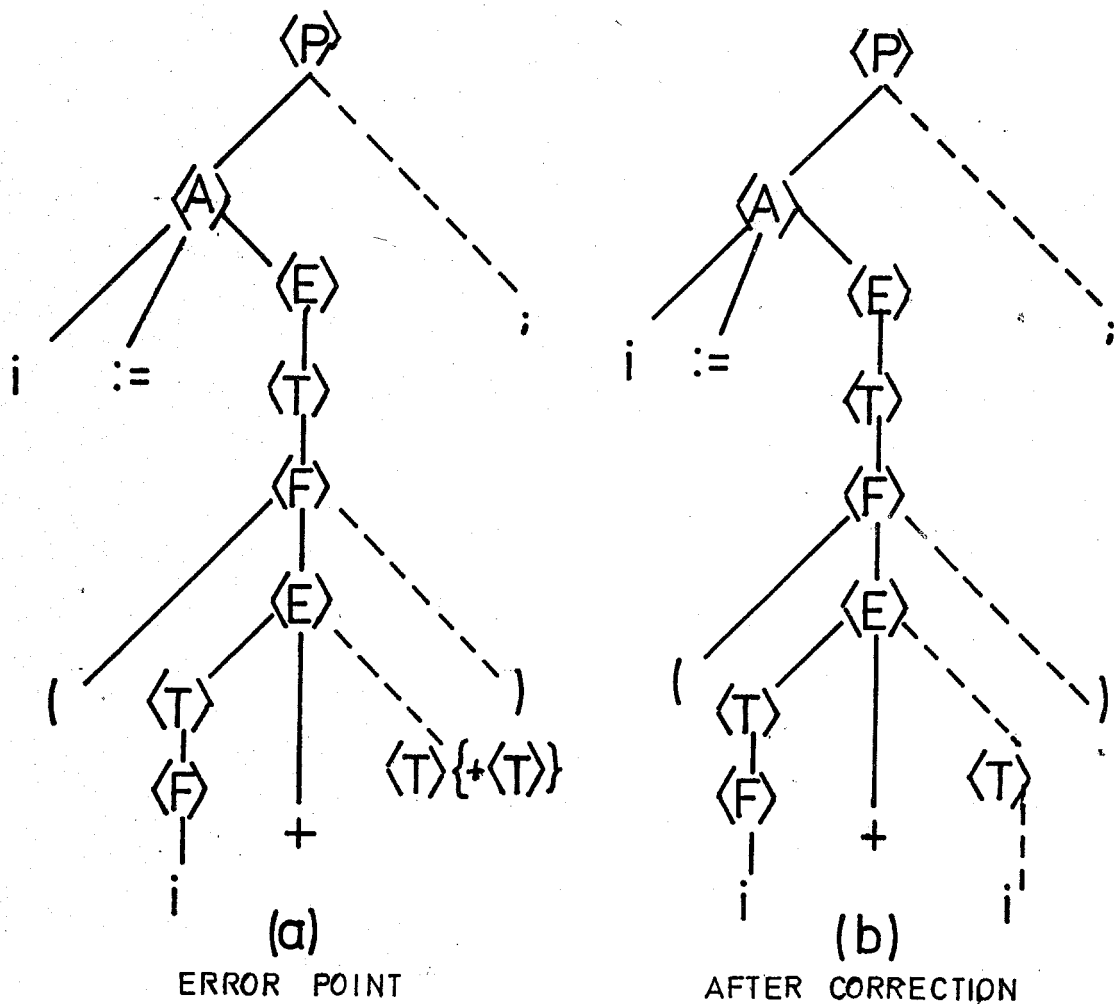


Figure 4.3- Top Down Error Recovery

The incomplete branch which caused ")" to be put in L is $\langle F \rangle ::= (\langle E \rangle)$. To cause ")" to be associated with this branch, we must insert a string to complete the branch $\langle E \rangle ::= \langle T \rangle \{ + \langle T \rangle \}$. Recovery algorithm always tries to insert simplest string possible, and this again is identifier i .

2. Recovery Algorithm

The recovery technique described in this section is quite applicable to our parsing algorithm, since we already have incomplete branches kept in our parse stack.

Step 1 of the recovery algorithm requires preparing a list of the symbols in the incomplete parts of the incompleted

branches, to be used by step 2. By the help of the parse stack it is easy to prepare this list, but do we need such a list explicitly? Actually no! Using iteration, incompleted parts can be reached and transmitted to step 2 one by one. Yet, since it's convenient to assume the existence of such a list, we do it so.

Purpose of step 2, is to search the list, for some member U, such that $U \Rightarrow *...T...$ where T corresponds to the last scanned source symbol. It is possible that no member U, satisfying this condition exists. Then we just discard T, get next source symbol and repeat the process with this symbol as being T. Finally, the desired member U will be reached. A BOOLEAN function named FIND, is declared to perform step 2. FIND is repeatedly called for each member of the list, until a member satisfying the above condition is found.

Search of step 2 starts from current STABLE position (or SPINX) to test symbols of "local" context. Then stack must be searched (incomplete branches) to determine symbols in the "global" context, starting from top of the stack, going downwards since it is most probable that the error is due to an incomplete branch near the stacktop.

Steps 1 and 2 are combined into a single step and expressed algorithmically in table 4.1. In this algorithm FOUND is a boolean variable and STIND is an integer variable used to index STACK. As in table 3.3, SUCC is assumed to be a function that returns successor of its parameter.

Every source symbol discarded means an illegal symbol so an associated ERROR routine is called. When we exit from search loop, STIND indicates the incomplete branch which accepts last scanned token, T as a member of its definition. Thus step 3 of the formal algorithm is also covered.

TABLE 4.1- Recovery Algorithm--Part 1

```

REPEAT
  FOUND := FIND(SUCC (SPINX)) ; (* TEST LOCAL CONTEXT *)
  STIND := STACKTOP ;
  (*****TEST GLOBAL CONTEXT*****)
  WHILE (NOT FOUND) AND (STIND > 0) (* STACK NOT EMPTY *) DO
    BEGIN
      FOUND := FIND (SUCC (STACKESTINDJ)) ;
      STIND := STIND - 1
    END ;

  IF NOT FOUND THEN (* DISCARD T AND GET NEXT TOKEN *)
    BEGIN
      ERROR(6) ;
      SCAN (* GET NEXT TOKEN *)
    END
  UNTIL FOUND ;

```

So, we come to step 4 of the initial algorithm where insertion string q is determined. On exit from the loop of table 4.1. first, the value of $STIND$ is tested. If its value is less than the value of $STACKTOP$, that means there are incomplete branches in between which must be manually completed. In other words, there are non-terminals indicated by $STACK$ entries between $STIND$ and $STACKTOP$, each of which, requires a terminal string generated for them. This test is implemented as a loop which calls a procedure, named $STRING$, repeatedly (Table 4.2).

TABLE 4.2- Recovery Algorithm--Part 2

```

STEMP := STACKTOP ;
WHILE STIND < STEMPT DO
  BEGIN
    STRING (STABLE [STACK[STEMPT]]. [PTYPE]) ;
    STEMPT := STEMPT - 1
  END ;

```

$STRING$ is a procedure that receives a non-terminal as its parameter. It generates simplest possible terminal string for the non-terminal denoted by its parameter and catenates it to q .

Treatment of the incomplete branch where the error is found (the $STACK$ entry indexed by $STIND$) will be different. In this case a string will be generated (and catenated to q)

until the last scanned token T, is reached in the definition, unlike the way STRING works where, the string is generated until the expansion is completed. This last part is accomplished by a call to a routine named INSERT. This routine receives successor of the STABLE entry indexed by STACK[STIND].

When q is completely determined, every symbol in it indicates a missing symbol in the position where the error is detected. So for every symbol in q, an error message of the form ".....EXPECTED" is printed.

3. Description of Routines

i) FUNCTION FIND

This function receives as input a structure table (STABLE) entry. It will try to reach the last scanned source symbol (or T) by moving through the links associated with its input. While doing this, it must visit all the successors and alternates of its input. But if its input is a non-terminal it must look to its definitions successors and alternates also and so on. So function FIND is designed as a recursive routine.

Calling sequence of FIND is based on its input and defined as follows.

```

if input is a NON-TERMINAL
    call itself with DEFINITION of input;
    call itself with SUCCESSOR of input

else
    call itself with ALTERNATE of input;
    call itself with SUCCESSOR of input
  
```

By this way all the successors and alternates for an STABLE entry can be visited. Table 4.3. gives the complete algorithm of FIND. Here parameter SPINX is

the input value. Since there are three distinct calls of FIND there must be three stopping criterias to control them;

- a) When calling FIND with SUCC(SPINX) as argument, if SUCC(SPINX)=OK or SPINX>SUCC(SPINX), that indicates end of successors for SPINX, so call must be prohibited.
- b) When calling FIND with ALT(SPINX) as argument, if ALT(SPINX)=OK or ALT(SPINX)=FAIL, that means there are no alternates for SPINX, so call is not performed.

TABLE 4.3- Algorithm FIND

```

FUNCTION FIND (SPINX : INTEGER) : BOOLEAN ;
VAR FOUND : BOOLEAN ;
BEGIN
  FOUND := FALSE ;
  WITH STABLE [ SPINX ], TTABLE [ TPTYPE ] DO
    CASE TERMINAL OF
      FALSE :
        BEGIN
          IF UNMARKED (FINDCTRL, TPTYPE) THEN
            FOUND := FIND (POINTER) ;
          IF NOT FOUND THEN
            IF (SUC <> OK) AND (SUC > SPINX) THEN
              FOUND := FIND (SUC) ;
            END ;
          TRUE :
            BEGIN
              IF CLASS = TOKEN.CLASS
                THEN FOUND := TRUE
              ELSE
                BEGIN
                  IF (ALT <> FAIL) AND (ALT <> OK) THEN
                    FOUND := FIND (ALT) ;
                  IF NOT FOUND THEN
                    IF (SUC <> OK) AND (SUC > SPINX) THEN
                      FOUND := FIND (SUC) ;
                    END
                END
            END
        END (* CASE *) ; FOUND := FOUND
    END ;
  END (* FIND *)

```

- c) Calling FIND with the definition of a non-terminal may cause infinite loops, since the grammar is recursively defined (directly or indirectly). Consider a rule like $\langle \text{exp} \rangle ::= i | (\langle \text{exp} \rangle)$. Call of FIND with $\langle \text{exp} \rangle$ may cause another call of FIND with $\langle \text{exp} \rangle$ and so on, since $\langle \text{exp} \rangle$ contains $\langle \text{exp} \rangle$ in its definition.

Thus, in order not to refer a non-terminal more than once, a boolean function UNMARKED is defined which returns true if its argument has already been tried.

ii) PROCEDURE STRING

This procedure receives an STABLE entry corresponding to a non-terminal as its input. It will try to generate simplest possible terminal string for its input. Procedure STRING is also recursively defined, since definition of its input may also contain non-terminals.

When a successor link with value OK is met (which indicates end of expansion) or an alternate link with OK value is tried (which shows optionality), the execution of STRING will be terminated. Table 4.4. shows the algorithm of STRING. To concatenate the current symbol to the end of q, a procedure named CATENATE is defined.

TABLE 4.4- Algorithm STRING

```

PROCEDURE STRING ( SPINX : INTEGER ) ;
BEGIN
  WITH STABLE [ SPINX ], TTABLE [ TPTYPE ] DO
    CASE TERMINAL OF
      FALSE :
        BEGIN
          STRING ( POINTER ) ;
          IF SUC <> OK THEN
            STRING ( SUC )
          END ;
        TRUE
        IF ( ALT = FAIL ) OR ( SUC = OK )
        THEN
          BEGIN
            CATENATE ( Q, CLASS ) ;
            IF SUC <> OK THEN
              STRING ( SUC )
            END
          ELSE
            IF ALT <> OK THEN
              STRING ( ALT )
            END
          END ;
        END ;
      END ;
    END ;
  END ;
  (* CASE *)
  (* STRING *)

```

iii) FUNCTION INSERT

This routine is the last portion of the recovery algorithm. It receives the STABLE entry, that contains

last scanned source symbol, T in its definition as input. INSERT is recursively defined as the previous two routines. It generates a terminal string and catenates it to q. Execution of INSERT will terminate when T is reached. INSERT is designed as a boolean function. If it returns false that means a compiler error exists. Table 4.5. gives the algorithm of INSERT. This algorithm is a simplified one. Actual algorithm is much more complex, to take care of infinite loops that frequently occur due to recursive definition of the grammar. For actual algorithm refer to appendix A.

TABLE 4.5- Algorithm INSERT

```

FUNCTION INSERT (SPIXK : SPRANGE) : BOOLEAN ;
BEGIN
  WITH STABLE [SPIXK], TABLE [TPTYPE] DO
    CASE TERMINAL OF
      TRUE : IF CLASS = TOKEN.CLASS THEN (* T IS REACHED *)
              REACHED := TRUE
            ELSE IF (ALTERNATE > FAIL) AND (FIND(ALTERNATE)) THEN
              REACHED := INSERT (ALTERNATE)
            ELSE BEGIN
                  CATENATE (Q, CLASS) ;
                  IF (SUCCESSOR <> OK) AND (SUCCESSOR > SPIXK) THEN
                    REACHED := INSERT (SUCCESSOR)
                  END ;
                FALSE : IF NOT (FIND(POINTER)) THEN BEGIN
                        STRING (POINTER) ;
                        IF (SUCCESSOR <> OK) THEN
                          REACHED := INSERT (SUCCESSOR) END
                        ELSE IF (SUCCESSOR > OK) AND (FIND(SUCCESSOR)) THEN
                          BEGIN STRING (POINTER) ;
                                REACHED := INSERT (SUCCESSOR)
                          END
                        ELSE REACHED := INSERT (POINTER)
                      END ;
                    (* CASE *)
                    (* INSERT *)
                END ;
            END ;
  END ;

```

Table 4.6 gives the trace of the RECOVERY algorithm using the grammar of table 3.2. Input is assumed to be w=A>B*)#. This sentence is analyzed exactly same as table 3.3. until step 13. So table 4.6. outlines steps after 13.

E. Conclusion

Error recovery requires, we can continue analyzing without too much possibility of generating several error

messages for a single error. If we do happen to generate few, it does not really matter. This can be provided by keeping the insertion string q , as small as possible. So procedure `STRING` and `FUNCTION insert` are designed, to minimize the length of q .

Sometimes, recovery algorithm causes extra errors in the incoming source symbols, because symbols of q may conflict with the actual purpose of the programmer.

The recovery algorithm used in this thesis is described in `GRIES(2)`.

TABLE 4.6- continued

<u>STEP#</u>	<u>SPINX</u>	<u>IS TERMINAL</u>	<u>OPERATION</u>	<u>INPUT</u>	<u>STACK</u>
			4.5.4. No, SUCC(11) <> OK and 11 < (SUCC(11)=12)?		
			4.5.5. Yes, FIND(SUCC(11))		
12			4.6. IN FIND with SPINX=12		
			4.6.1. Is TERMINAL(12)		
			4.6.2. No, UNMARKED (TPTYPE(12))		
			4.6.3. No, SUCC(12) <> OK OR 12 <SUCC(12)=13?		
			4.6.4. Yes, FIND(SUCC(12))		
13			4.7. In FIND with SPINX=13		
			4.7.1. Is TERMINAL(13)		
))		
			4.7.2. Yes, current elmt [?] =input symbol		
			4.7.3. Yes, FOUND:=TRUE		
			4.7.4. FIND:=TRUE		
			{FIND RETURNS TRUE}		
			4.8. STIND := STACKTOP		
			4.9. q := empty string		
			4.10. REACHED:=INSERT(11)		
11			4.11. In INSERT with SPINX=11		
			4.11.1. Is TERMINAL(11)		
			4.11.2. Yes, Current symbol [?] =input symbol		
			4.11.3. No, ALT(11) FAIL		
			4.11.4. No.		
			4.11.4.1. q:=CATENATE(q,"")		
			4.11.4.2. (SUCC(11) <> OK) and (12>11)?		
			4.11.4.3. Yes, REACHED:= INSERT(12)		

TABLE 4.6- continued

STEP#	SPINX	IS TERMINAL	OPERATION	INPUT	STACK
	12		4.12. In INSERT with SPINX=12		
			4.12.1. Is TERMINAL(12)		
			4.12.2. No, FIND(12)?{refer operation 4.1}		
			4.12.3. Yes, SUCC(12) <> OK		
			4.12.4. Yes, FIND(SUCC(12))		
	13		4.12.5. In FIND with SPINX=13 FIND returns true, immediate success		
	1		4.12.6. In STRING with SPINX=1		
	1		4.12.6.1. Is TERMINAL(1)		
			4.12.6.2. No, STRING(5)		
	5		4.12.7. In STRING with SPINX=5		
			4.12.7.1. Is TERMINAL(5)		
			4.12.7.2. No, STRING(7)		
	7		4.12.8. In STRING with SPINX=7		
			4.12.8.1. Is TERMINAL(7)		
			4.12.8.2. No, STRING(10)		
	10		4.12.9. In STRING with SPINX=10		
			4.12.9.1. Is, TERMINAL(10)		
			4.12.9.2. Yes, (ALT(10)= FAIL)or(SUCC(10)= OK)?		
			4.12.9.3. Yes, q:=q IDENT= (IDENT : ↓ return form STRING ↓		
	12		4.2.10. REACHED:=INSERT(SUCC(12))		
	13		4.13. In INSERT with SPINX=13) 4.13.1. Current elmt=input symbol 4.13.2. Yes, return from INSERT with input		
				A>B*(ident) ↑ head of q	
16	11		1. current element=input symbol 2. Yes ...		

9
4
3

V. SYMBOL TABLE

A. Definition

A translator needs to collect and use information about names appearing in the source program. This information is entered into a data structure called a SYMBOL table. The information collected about a name includes the string of characters by which it is denoted, its type, its structure etc.

Each time a name is encountered, the symbol table is searched to see whether that name has been seen previously. If it is new, it is entered into the symbol table. Information about a name is entered, by syntactic analyzer while parsing declarations.

Information collected in the symbol table, is used in semantic analysis, (that is, in checking uses of identifiers are, consistent with their declarations) and in code generation.

Symbol table can be used to aid in error detection and correction. For example, we can record whether an error message such as "variable A undefined" has been printed out before, and refrain from doing so more than once.

In block structured languages the same identifier can be used to represent distinct names with nested scopes. In such languages, the symbol table mechanism must make sure that the innermost occurrence of an identifier is always found first and that names are removed from the active portion of

the symbol table when they are no longer active.

Symbol table mechanism, thus should allow us;

1. Determine whether a given name is in the table,
2. add a new name to the table,
3. access the information associated with a given name,
4. add new information for a given name,
5. delete a name or groups of names from the table.

B. Symbol Table Organizations

This section describes the ways of representing symbol tables in general.

1. Unsorted and Sorted Tables

The easiest way to organize a table is to add entries in the order they arrive. A search requires $N/2$ comparisons on the average, for a match if N is greater than 20, and this is inefficient.

Searching can be performed more efficiently if the table entries are sorted according to string of characters denoting the name. Efficient search techniques such as binary search can be used in this case.

Another method of accessing symbols in a table is using hash-addressing. This is a technique for converting symbols to indexes of entries in the table (the indexes are numbered $0, 1, 2, \dots, N-1$ where the table has N entries). The index is obtained by "hashing" the symbol, i.e. by performing some simple arithmetic or logical operation on the symbol. As long as, two symbols do not hash to the same index, we have no problem. Trouble occurs, however if two symbols hash to the same index. This is called a collision, and the hash algorithm must take care of it.

2. Block Structured Tables

Algol-like languages have a nested block and procedure structure. The same identifier may be declared and used many times in different blocks and procedures and each such declaration must have a unique symbol table entry associated with it. Given an identifier, the problem is then to discover the correct symbol table entry for it.

The rule is to look first in the current block, then the surrounding block and so on, until a declaration of that identifier is found. Such a search can be implemented by keeping all the symbol table entries for each block contiguous, and by using a block list (Figure 5.1).

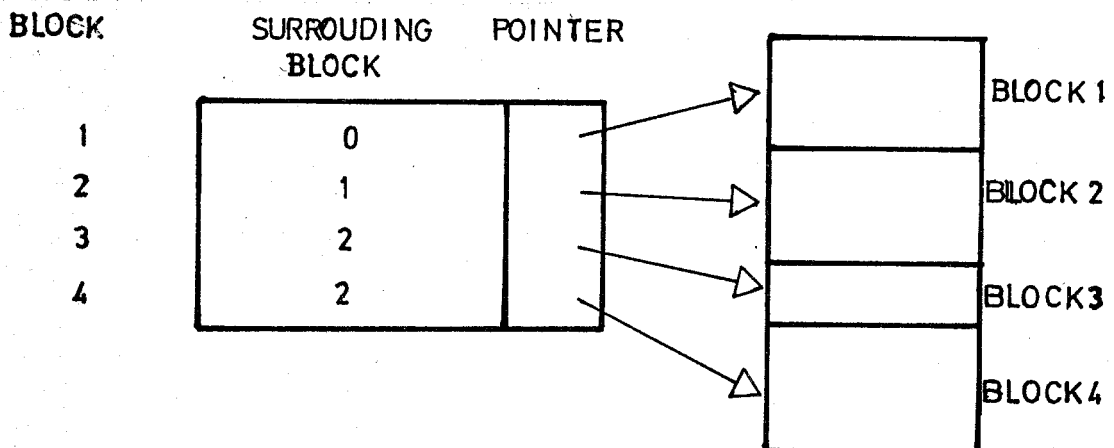


Figure 5.1- Block Structure.

Once an associated block is found, searching it would be simple.

3. Tree Structured Tables

This organization strategy is used in the implementation. The method uses a binary tree to order the entries. Each node of the tree represents a filled entry of the table, the root node being entry 1. Figure 5.2a shows the table with one entry for identifier G. Suppose now that the identifier D is to be entered. A branch is drawn for it to the left, since $D < G$ (figure 5.2b). Now let the M be

entered. Since $G < M$ a branch is drawn for it to the right from G (figure 5.2c). Finally let the identifier E be entered. $E < G$, so we travel down the left branch from G, and to the right of D (figure 5.2b). Figure 5.2e shows the tree after identifiers A, B and F have been added in that order.

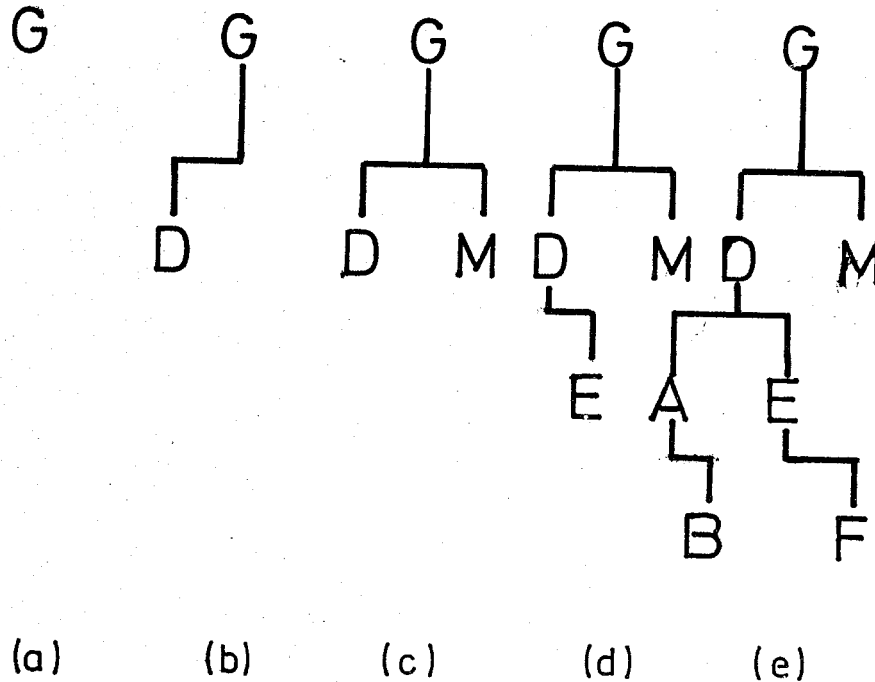


Figure 5.2. Binary Tree Illustration.

One can implement this, by having two pointer fields with each entry, one for the left and one for the right branches.

C. Implementation

In this section data structures and routines used for symbol table manipulation in this thesis will be explained.

1. Organization of Symbol Table

As mentioned before symbol table is organized in a tree structured manner. Since, PASCAL has a nested procedure structure, each procedure has its own tree structure to store its local variables. In other words, symbol table is

organized as an unbalanced binary tree at each level.

It is also necessary to access variables of surrounding procedure (i.e. global variables). This means access to the trees of higher levels. So, keeping trace of levels (with their tree pointers) is necessary. This is accomplished by a stack. (Section 6.3.4, describes implementation of this stack).

2. Data Structures to Represent the Symbol Table

There are two Basic Tables

- i) STRUCTURE: Used to keep information about the attributes of variables. That is, it describes the structure of data types.
- ii) IDNTFR : Holds the identifiers declared throughout the program.

Dynamic allocation facilities and variant records of PASCAL are ideal to represent these tables.

Consider table 5.2. Possible structures of PASCAL data types are given by the enumerated type STRUCTFORM. Similarly IDCLASS shows classes of identifiers.

Field NAME of record IDNTFR is used to keep the actual symbol itself and it is used as a search key. Fields LLINK, RLINK in the same record are left and right branch pointers of the tree. However left and right links are not sufficient. Sometimes sequential linking becomes necessary, for example in a parameter list (or in a user defined scalar) where the parameters (scalar identifiers) must be linked in the order which they are declared. IDTYPE is the STP pointer indicating the structure of the identifier recognized.

Variants of IDNTFR which are dependent on the tagfield KCLASS are introduced to indicate level counters, to differentiate actual-formal declarations and to state values of constants.

TABLE 5.2- Data Structures for Symbol Table

```

CONST MAXLEVEL = 20 ;

TYPE
STRUCTFORM = (SCALAR, SUBRANGE, POINTER, POWER, ARRAYS,
RECORDS, DAPAMLIST, FILES, TAGFIELD, VARIANT);
DECLKIND   = (STANDARD, DECLARED);
STP        = STRUCTURE ;
IDP        = IDENTIFR ;
STRUCTURE  = PACKED RECORD
CASE FORM : STRUCTFORM OF
  SCALAR : ( CASE SCALKIND : DECLKIND OF
DECLARED : (FSTCONST : IDP);
STANDARD : ( ) ) ;
  SUBRANGE : ( RANGETYPE : STP ;
MIN, MAX : INTEGER ) ;
  POINTER : ( ELTYPE : STP ) ;
  POWER : ( ELSET : STP ) ;
  ARRAYS : ( PACKO : BOOLEAN ;
INXTYPE, AELTYPE : STP ) ;
  RECORDS : ( SPACK : BOOLEAN ;
FSTFLD : IDP ;
RECVAR : STP ) ;
  DAPAMLIST : ( FSTPAR : IDP ) ;
  FILES : ( FILTYPE : STP ) ;
  TAGFIELD : ( TAGFIELDP : IDP ;
TAGTYPE : STP ;
FSTVAR : STP ) ;
  VARIANT : ( NXTVAR, SJ3REC : STP ;
VARVAL : INTEGER ) ;
END ;

IDCLASS = (TYPES, KONST, VARS, FIELD, PARAMS, FUNC, PROC, PROG) ;
IDKIND = (ACTUAL, FORMAL) ;
LEVRANGE = 0 ... MAXLEVEL ;

IDENTIFR = PACKED RECORD
NAME : PACKED ARRAY [1..12] OF CHAR ;
IDTYPE : STP ;
NEXT : IDP ;
LLINK, RLINK : IDP ;
CASE CLASS : IDCLASS OF
  KONST : (VALUES : CONSTANT);
  VARS : (VKIND : IDKIND ;
VLEV : LEVRANGE);
  PROC, FUNC :
(CASE PFDECLKIND : DECLKIND OF
STANDARD : ( ) ;
DECLARED :
(PFLEV : LEVRANGE ;
PARAMPTP : STP ;
CASE PFKIND : IDKIND OF
ACTUAL : (FORWDECL : BOOLEAN)))
END ;

```

Now consider STRUCTURE record. Here a tagfield of type STRUCTFORM is used. If the FORM is (user defined) SCALAR we need to know IDNTFR pointer of the first enumerated constant, (field FSTCONST). If it is SUBRANGE, then we need to know RANGETYPE (it must be SCALAR) which points to an another STRUCTURE entry. Fields MIN and MAX contain the lower and upper bounds of the range respectively. When the FORM is ARRAYS we need to know if it is packed, how it is indexed and type of its element. Multi-dimensional arrays are represented as arrays of arrays, so there is no need to have a field indicating how many dimensions the array has.

Other variants except for RECORDS, I believe are clearly understandable from table 5.2. Record structures must have their own trees like procedures, since fields of a record are in accessible if the name of the record is not specified. Thus field FSTFLD of the variant RECORDS should point to the root node of the associated tree. Field RECVAR is used to point to the TAG information if the record contains variants otherwise it is NIL.

A record type may be considered as a road map to an area of memory. It defines how the memory is to be interpreted. A variant record type provides several different road maps for the same area of memory, and a tag field value determines which road map is currently in use. So it is reasonable to consider each variant as a subrecord, activated according to the value of tag field. (If there is no tag field any of these subrecords can be activated arbitrarily).

Field TAGFIELDP of variant TAGFIELD is an IDNTFR pointer indicating the tag, and is NIL if there is no tag field. Field TAGTYPE indicates the type of tag which must be a scalar type. Field FSTVAR points to the first variant of the record.

TABLE 5.3- Symbol Table Dump for a Variant Record

```

TYPE  A = RECORD
      N,ZZ : REAL ;
      CASE M : INTEGER OF
        1,2 : ( K : REAL ) ;
        3   : ( J : INTEGER )
      END ;

**** IDP1 @ :
      NAME = 'A'      ; IDTYPE = STP1 ;   KLASS = TYPES ;
**** STP1 @ :
      FORM = RECORDS ; FSTFLD = IDP2   ; RECVAR = STP2 ;
**** IDP2 @ :
      NAME = 'N'      ; IDTYPE = REALPTR ; KLASS = FIELD ;
**** IDP3 @ :
      NAME = 'ZZ'     ; IDTYPE = REALPTR ; KLASS = FIELD ;
**** STP2 @ :
      FORM           = TAGFIELD ; TAGFIELDOP = IDP4 ;
      TAGTYPE       = INTPTR  ; FSTIVAR   = STP3 ;
**** IDP4 @ :
      NAME = 'M'      ; IDTYPE = STP2   ; KLASS = FIELD ;
**** STP3 @ :
      FORM = VARIANT ; NXTVAR = STP4 ;
      SUBREC = STP5 ; VARVAL = 1 ;
**** STP4 @ :
      FORM = VARIANT ; NXTVAR = STP6 ;
      SUBREC = STP5 ; VARVAL = 2 ;
**** STP5 @ :
      FORM = RECORDS ; FSTFLD = IDP5 ; RECVAR = NIL ;
**** IDP5 @ :
      NAME = 'K'      ; IDTYPE = REALPTR ; KLASS = FIELD ;
**** STP6 @ :
      FORM = VARIANT ; NXTVAR = NIL ;
      SUBREC = STP7 ; VARVAL = 3 ;
**** STP7 @ :
      FORM = RECORDS ; FSTFLD = IDP6 ;
      RECVAR = NIL ;
**** IDP6 @ :
      NAME = 'J'      ; IDTYPE = INTPTR ; KLASS = FIELD ;

```

Field SUBREC of variant VARIANT points to the subrecord which will be activated if the tag field value is equal to the value of the field VARVAL. (VARVAL is unused if there is no tag field). Eventually field NXTVAR points to the next variant of the record if there are any.

To understand fully the representation of a variant record it is necessary to examine Table 5.3. In this example, variables IDP1 to IDP6 are assumed to be pointer constants of type IDP, whereas variables STP1 to STP7 are pointer constants of type STP. Similarly INTPTR and REALPTR are STP pointer constants, pointing to the definition of standard PASCAL types INTEGER and REAL (Figure 5.2 gives the linked list representation of table 5.3).

3. Predefined PASCAL Identifiers

The following identifiers are entered to the tree of level 0, with their STRUCTURE definitions during initialization.

1. MAXINT

is a constant whose value is dependent on the machine,

2. INTEGER

is a standard scalar type. Its value ranges between -MAXINT ... MAXINT.

3. REAL

is a standard scalar type. Its values are an implementation dependent finite subset of real numbers.

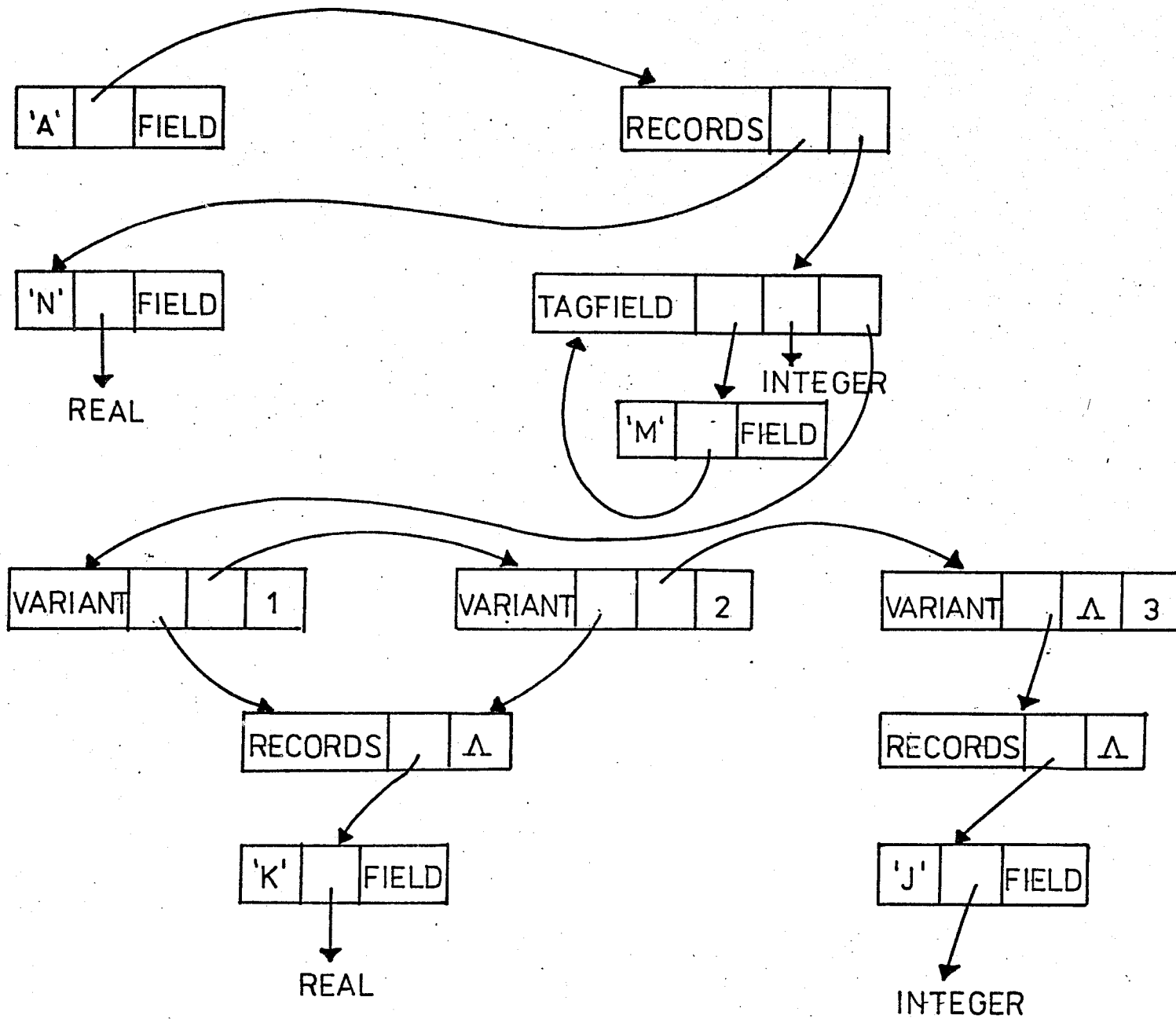
4. CHAR

is a standard scalar type whose values are a set of implementation dependent characters.

5. BOOLEAN

is a pre-declared scalar type which is defined as
type BOOLEAN = (FALSE, TRUE);

Figure 5.3-- Linked List Representation of TABLE 5.3.



Therefore FALSE and TRUE are constants of type BOOLEAN enumerated as 0,1 in that order.

6. TEXT

is a predefined file type which is declared as
type TEXT = file of CHAR;

7. NIL

is a implementation dependent constant used for pointers to indicate null entry.

8. OTHER IDENTIFIERS

These include standard PASCAL procedures and functions such as READ, WRITE, EDLN, EOF etc.

4. Local Tree Pointers

A single dimensional array indexed by level counter is used as a stack, to point local trees. This stack is declared as;

```
var DISPLAY: array [0..20] of IDP:
      TOP      : 0..20; {level counter which is initially
                       zero}
```

TOP is incremented each time a procedure is met. First variable will be entered into the DISPLAY [TOP]. Other variables of the procedure will be entered (or searched) taking DISPLAY [TOP] as the root node.

Similarly TOP is decremented each time a procedure is terminated. Thus the tree of that level will be deallocated, since it won't be used anymore.

Current tree is pointed by DISPLAY [TOP] every time. Trees of surrounding procedures can be reached by referring DISPLAY entries whose indexes are lower than the value of TOP.

5. Symbol Table Manipulation Routines

1. PROCEDURE ENTERID (IDPTR:IDP);
Enters identifier pointed by IDPTR (according to the algorithm stated in section 5.B.3) into the symbol table into the innermost level.
2. PROCEDURE SEARCHID (KLASS: set of IDCLASS; IDPTR:IDP);
Searchs identifier pointed by IDPTR whose class is included in the set KLASS. The search includes whole symbol table, that is trees of all higher levels.

If it fails that means an undeclared identifier is met. So SEARCHID enters it to the symbol table by calling a procedure DECLARE to suppress extra error messages.

3. PROCEDURE DECLARE (IDPTR:IDP);
Used to suppress multiple printout of error messages as mentioned above.
4. PROCEDURE SEARCHSECTION (FCP:IDP; var FCP1=IDP);
Searchs the tree whose root node is pointed by FCP and returns result in FCP1. This kind of search is necessary for instance in searching fields of a record or searching forward declared pointer types.
5. PROCEDURE ENTERSTDIDS;
Used to enter predefined PASCAL identifiers at level 0.
6. PROCEDURE PRINTABLES;
Dumps symbol table, separately at each level.
7. PROCEDURE SEMANTICS (ACTION:INTEGER);
As mentioned before BUPASCAL makes semantic analysis at the declaration level. In other words it enters the identifiers declared throughout the program into the symbol table, keeps information about them and detects associated errors. Thus, purpose of this procedure is to perform these tasks. Calling sequences for the procedures ENTERID, SEARCHID and SEARCHSECTION is

controlled by this procedure.

SEMANTICS is called by the parser each time a terminal is scanned or a non-terminal is stacked (unstacked). Assume that the terminal IDENT is scanned and we are processing var declarations. It will be searched within the symbol table (PROCEDURE SEARCHID) to see if it is previously declared if so an error message is given otherwise it will be entered into the symbol table by calling procedure ENTERID. Successive calls of SEMANTICS will be, thus by terminals indicating the structure of the identifier being processed, for instance successive terminal might be the keyword ARRAYSY. So associated STRUCTURE record, can be modified to have FORM=ARRAYS. Similarly all attributes can be determined. SEMANTICS is called when a non-terminal is stacked, to set some flags, to take some initiative actions etc. or when a non-terminal is unstacked to reset some flags, to decrement a level counter etc.

SEMANTICS will be called using current STABLE entry (SPINX) as argument. Thus SPINX uniquely determines the action to be taken. For non-terminals, SPINX will be negated to indicate the non-terminal is unstacked. SEMANTICS uses a case statement based on its formal parameter ACTION, to perform necessary action. If a case label whose value is equal to the value of ACTION does not exist SEMANTICS will do nothing.

D. Conclusion

The symbol table mechanism of this implementation uses dynamic trees. Thus size of the symbol table is not restricted. On the other hand searching technique used (binary search) has an order of $\log N$ (where N is number of nodes in the tree) is quite efficient unless the number of nodes is very few.

GRIES (2) and AHO-ULLMAN (1) describes several symbol

organization techniques. HOROWITZ-SAHNI (3) explains dynamic allocation and deallocation of binary trees.

Symbol table will be dumped if the option \$PRINTABLES is specified. This dump is similar to that of table 5.3, except actual pointer constants are used. Variables local to a block (program or procedure) are printed just before the first "begin" is scanned (predefined identifiers are also included while printing variables of PROGRAM .. i.e. global variables).

VI. CONCLUSION

The finalized version of BUPASCAL has been successfully tested on a number of input programs, four of which are given in the appendix. These sample inputs contain arbitrary compile time errors to test error recovery capability of BUPASCAL.

Even though BUPASCAL accepts programs written in STANDARD PASCAL as input, it is designed to accept any language whose grammar is given by parse tables. BUPASCAL uses two static tables (TTABLE and STABLE) for grammar representation. The size of TTABLE depends on the size of the vocabulary of the language under consideration, where as the size of STABLE depends on the number of rules used to define the grammar.

Therefore recovery algorithm is also designed, to respond to any language. Because of recursive and iterative representations of grammars, I believe recovery algorithm was the most important and difficult part of this thesis. Test runs showed that, even the most terrible errors can be recovered by the algorithm presented here.

BUPASCAL does not include semantic analysis at the statement level and code generation. But the design of parse tables allows subroutines (or coroutines) to be associated with the productions of the grammar. These routines are to be intended to perform semantic analysis and to generate intermediate code when called at appropriate times by the syntactic analyzer. When a severe error is detected, code generation

terminates where as semantic analysis should continue until the whole program is processed. An associated routine is called each time a syntactic primitive (or terminal) is met and each time a stack operation is performed.

Procedure SEMANTICS defined in section 5 is designed in this way. In this procedure ACTION values correspond to the coroutines described above. This routine can be modified to include semantics analysis at the statement level and code generation by enlarging the range of ACTION values. (i.e. by increasing case labels to respond all entries of STABLE).

APPENDIX A
SOURCE PROGRAM

```

*****
*) PROGRAM BUPASCAL (INPUT,OUTPUT) ;
*
* PURPOSE
* IMPLEMENTATION OF AN ANALYZER FOR PASCAL.
* STANDARD PASCAL IS USED.
* INPUT, OUTPUT MODE IS FIELDATA.
* DESIGNED BY ..
* CEM ATAC
* COMPUTER ..
* UNIVAC 1106
* DATE ..
* SEPTEMBER 16, 1982
* CENTER ..
* BOSPHOROUS UNIVERSITY COMPUTER CENTER
*
*****

```

LABEL 1111 ; (* MAIN RETURN ADDRESS *)
CONST

```

CHARMAX      = 12 ; BUFMAX      = 90 ; MAXERR      = 10 ;
FIRSTCHAR    = '3' ; LASTCHAR   = '9' ; MAXSET      = 71 ;
STRGLGTH     = 78 ; SETSIZE    = 72 ; NSETS       = 12 ;
FAIL         = 268 ; OK        = 92 ; NOACT       = 0 ;
SPSIZE       = 268 ; TPSIZE    = 93 ;
MAXREAL      = 0.898846567431157951E+307 ;
SQRTREAL     = 0.670390396497129853E+154 ;
DISPLIMIT    = 20 ; MINMIN     = 5000 ; MAXMIN      = 5000 ;
MAXLEVEL     = 20 ;

```

TYPE

```

SYMBOL = ( IDENT, INTCONST, REALCONST, CHARCONST,
STRINGCONST, EOFPGM, MULOP, ADDOP,
RELOP, LPARENT, RPARENT, LBRACKET,
RBRACKET, COMMA, SEMICOLON, PERIOD,
DOTDOT, ARROW, COLON, ASGNOP,
ARRAYSY, BEGINSY, CASESY, CONSTSY,
DOSY, DOWNTOSY, ELSESY, ENDSY,
FILESY, FORSY, FORWARDSY, FUNCTIONSY,
GOTOSY, IFSY, LABELSY, NOTSY,
OFSY, PACKEDSY, PROCSY, PROGRAMSY,
RECORDSY, REPEATSY, SETSY, THENSY,
TOSY, TYPESY, UNTILSY, VARSY,
WHILESY, WITHSY, OTHERSY ) ;

```

```

OPERATOR = ( ASTR, RDIV, IDIV, ANDOP,
IMOD, PLUS, MINUS, OROP,
LTOP, LEOP, GEOP, GTOP,
NEOP, EQOP, INOP ) ;

```

```

STRGRANGE = 0 ; STRGLGTH ;
STRGINDEX = 1 ; STRGLGTH ;
SETRANGE = 0 ; MAXSET ;
SPRANGE = FAIL ; SPSIZE ;
TPRANGE = FAIL ; TPSIZE ;
NNONTRMS = 52 ; 93 ;
CHSIZE = 0 ; 168 ;
DISPRANGE = 0 ; DISPLIMIT ;

```

```

CSTCLASS = (INTGR, REEL, PSET, STRING, KHAR) ;
CHARSTRING = PACKED ARRAY [STRGINDEX] OF CHAR ;
CHARARRAY = ARRAY [STRGINDEX] OF CHAR ;

```

```

CSTADDR = CONSTANT ;
CONSTANT = PACKED RECORD
INTVAL : BOOLEAN
CASE CLASS : CSTCLASS OF

```

```

INTEGR : (IVAL : INTEGER )
REAL : (RVAL : REAL )
PSET : (PVAL : SET OF SETRANGE )
KHAR : (ORDCH : INTEGER )
STRING : (SLGTH : STRGRANGE )
          (SVAL : CHARSTRINGS )
END ;

```

```

TKNCLASS = (NOTNEEDED, OPRTR, DRADRES, INDRADRES) ;
TOKENS = RECORD
        CLASS : SYMBOL ;
        CASE TKNCLASS OF
          OPRTR : ( OP : OPERATOR ) ;
          DRADRES : ( CSTADR : CSTADDR ) ;
          NOTNEEDED : ( ) ;
        END ;

```

```

IDSTRING = ARRAY [1..CHARMAX] OF CHAR ;
PACKEDID = PACKED ARRAY [1..CHARMAX] OF CHAR ;
SETS = ARRAY [3..NSETS] OF
        PACKED SET OF SETRANGE ;
LINKS = 2 PARSESTACK ;
PARSESTACK = RECORD
          PIR : SPRANGE ;
          PREVIOUS : LINKS ;
        END ;

```

```

STRUCTFORM = (SCALAR, SUBRANGE, POINTER, POWER, ARRAYS,
              RECORDS, PARAMLIST, FILES, TAGFIELD, VARIANT) ;
DECLKIND = (STANDARD, DECLARED) ;
STP = 2 STRUCTURE ;
IDP = 2 IDNTR ;
STRUCTURE = RECORD
          SIZE : INTEGER ;
          MARKED : BOOLEAN ;
          CASE FORM : STRUCTFORM OF
            SCALAR : ( CASE SCALKIND : DECLKIND OF
                      DECLARED : (FSTCONST : IDP) ;
                      STANDARD : ( ) ) ;
            SUBRANGE : ( RANGETYPE : STP ;
                        MIN, MAX : INTEGER ) ;
            POINTER : ( ELTYPE : STP ) ;
            POWER : ( ELSET : STP ) ;
            ARRAYS : ( PACKD : BOOLEAN ;
                     INXTYPE, AELTYPE : STP ) ;
            RECORDS : ( RPACK, NOFIX : BOOLEAN ;
                       STFLO : IDP ;
                       RECVAR : STP ) ;
            PARAMLIST : ( ANON : BOOLEAN ;
                          FSTPAR : IDP ) ;
            FILES : ( FILTYPE : STP ) ;
            TAGFIELD : ( TAGFIELDP : IDP ;
                       TAGTYPE : STP ;
                       FSTVAR : STP ) ;
            VARIANT : ( NXTVAR, SUBREC : STP ;
                       VARVAL : INTEGER ) ;
          END ;

```

```

IDCLASS = (TYPES, KONST, VARS, FIELD, PARAMS, FUNC, PPROC, PROG) ;
SETOFIDS = SET OF IDCLASS ;
IDKIND = (ACTUAL, FORMAL) ;
LEVRANGE = 0..MAXLEVEL ;

```

```

IDNTR = PACKED RECORD
        NAME : PACKEDID ; LLINK, RLINK : IDP ;
        IDTYPE : STP ; NEXT : IDP ;
        JNDCL : INTEGER ;
        CASE KCLASS : IDCLASS OF
          KONST : (VALUES : CONSTANT) ;
          VARS : (VKIND : IDKIND ;
                 VLEV : LEVRANGE) ;
        PROC, FUNC :
          (CASE PFDECLKIND : DECLKIND OF
            STANDARD : ( ) ;
            DECLARED :
              (PFLEV : LEVRANGE ;
               PARAMPIR : STP ;
               CASE PFKIND : IDKIND OF

```


END ;

ACTUAL : (FORWDECL : BOOLEAN))

LABELSTATE = (REFERENCED, DEFINED, UNDEFINED) ;
LABELPTR = LABELTAB ;
FORWPPPTR = FORWPF ;

LABELTAB =
RECORD
LABVAL : INTEGER ;
LABNIV : LEVRANGE ;
NEXTLAB : LABELPTR ;
STATUS : LABELSTATE ;
END ;

FORWPF =
RECORD
NAME : PACKEDID ;
LINE : INTEGER ;
NEXT : FORWPPPTR ;
END ;

WHERE = (BLCK, REC) ;
UNIT =

PACKED RECORD
NAME : PACKEDID ; (* NAME OF UNIT *)
FIRST : INTEGER ; (* FIRST LINE OF UNIT *)
ERRS : INTEGER ; (* # OF ERRS BEFORE THIS UNIT *)
TYP : IDCLASS ; (* TYPE OF UNIT *)
EID4 : BOOLEAN ; (* ANY UNDECLARED ID'S ? *)
FORWP : FORWPPPTR ; (* LIST OF FORWARDS *)
FORWC : INTEGER ; (* # OF FORWARDS *)
END ;

VAR

(* SCANNER VARIABLES AND DATA STRUCTURES *)
(* ***** *)

SYMSTART : 1, BUFMAX ; (* START OF CURRENT SYMBOL *)
I, J, K, II, JJ, ; (* COUNTERS *)
CARDCNT, ; (* CARD COUNT *)
BUFLGTH, ; (* NO OF CHARS IN INBUF *)
ABORTLINE : INTEGER ; (* START OF COMMENT OR STRING *)
CH : CHAR ; (* LAST CHAR IN INBUF *)
CHCNT, ; (* INBUF CHARACTER COUNTER *)
LGTH : INTEGER ; (* LENGTH OF LAST STRING *)
ACH, ACH1 : ASCII ; (* ASCII CHARACTERS *)
LCASE : ARRAY ['A'..'Z'] OF ASCII ;
TOKEN : TOKENS ;
NAME : PACKEDID ;

DUMMYID, IDSTR : IDSTRING ;
EMPTYSTORE : PACKED ARRAY [1..80] OF CHAR ;
STRBUF, DUMMYSR : CHARARRAY ;

KONSPTR : CSTADDR ;
DIGITS, IDCHAPS : SET OF CHAR ;

DEFAULT : ARRAY [1..17] OF TOKENS ;
NEXT, CHECK : ARRAY [CHSIZED] OF INTEGER ;
BASE : ARRAY [0..17] OF INTEGER ;

CHPCLASS : ARRAY [CHARS] OF TOKENS ;
INBUF : ARRAY [1..BUFMAX] OF CHAR ;
UNCLOSED : (COMMENT, STRINGS, NONE) ;

(* VARIABLES FOR INDENTING AND DEDENTING *)
(* ***** *)

LNEST, ; (* CHANGING NEST LEVELS *)
RNEST : BOOLEAN ; (* CHANGING NEST LEVELS *)
NESTINCR, ; (* NEST INCREMENT *)
NESTDECR : INTEGER ; (* NEST DECREMENT *)
NEST : INTEGER ; (* NEST COUNTER *)
RTNLEVEL : INTEGER ; (* ROUTINE LEVEL COUNTER *)

(* STRUCTURES FOR ERROR HANDLING *)
(* ***** *)

ERRINX : 0, MAXERR ; (* NO OF CURRENT LINE ERRORS *)
LASTERR, ; (* INDEX OF LAST ERROR *)

```

SAVEDSP      ; (* SAVED SP INDEX *)
ERRCOUNT   ; (* NUMBER OF ERR FIELDS *)
ERRPOS      ; INTEGER ; (* ERROR POSITION *)
ALTFLAG     ; BOOLEAN ; (* FOR LOSTLINKS *)
MET         ; BOOLEAN ; (* TO FIND LOST LINKS *)
ERRLIST     ;
  ARRAY [1..MAXERR] OF
    PACKED RECORD
      POS : 0..80 ; (* POSITION OF ERROR *)
      ERRNUM : 1..1000 ; (* ASSOCIATED ERROR CODE *)
  END ;
LOSTLINKS   ;
  ARRAY [1..50] OF
    RECORD
      SPIND : SPRANGE ; (* POSITION OF LOST LINK *)
      TEMP : LINKS ; (* ASSOCIATED LINK *)
    END ;

```

```

ERPDUPL     ;
EMPTYSET    ; SETS ;
ERRBUF      ; ARRAY [1..BUFMAX] OF CHAR ;
ERRWARNCNT  ; ARRAY [BOOLEAN] OF INTEGER ;
MISSING     ; PACKED ARRAY [SYMBOL] OF 0..1000 ;
NONTRMSG    ; PACKED ARRAY [NONTRMS] OF 0..200 ;
ERRWARN     ; ARRAY [BOOLEAN] OF
  PACKED ARRAY [1..13] OF CHAR ;
ABORTMES    ; ARRAY [BOOLEAN] OF
  PACKED ARRAY [1..25] OF CHAR ;
RECOVERY    ; ARRAY [1..100] OF TOKENS ;

```

(* PARSER VARIABLES AND DATA STRUCTURES *)
 (*****)

```

STACKTOP    ; LINKS ; (* CURRENT STACKTOP *)
RECIND      ; (* RECOVERY STACK *)
RIND        ; INTEGER ; (* INDEXES *)
FOREVER     ; (* LOOP CONTROL FLAG *)
RECOVERED   ; BOOLEAN ; (* TRUE IF RECOVER IS CALLED *)
SAVE        ; (* TEMPORARY *)
LAST        ; LINKS ; (* TEMPORARY *)
SPIND       ; SPRANGE ; (* STABLE INDEX *)
STABLE      ;
  ARRAY [SPRANGE] OF
    PACKED RECORD
      SUC : SPRANGE ; (* SUCCESSOR FIELD *)
      ALT : SPRANGE ; (* ALTERNATE FIELD *)
      TPTYPE : TPRANGE ; (* BACKWARD POINTER *)
    END ;
TTABLE      ;
  ARRAY [TPRANGE] OF
    PACKED RECORD
      CASE TERMINAL : BOOLEAN OF
        FALSE : (* STABLE POINTER *)
          ( POINTER : SPRANGE ) ;
        TRUE : (* CLASS : SYMBOL *)
      END ;

```

(* SEMANTIC VARIABLES AND DATA STRUCTURES *)
 (*****)

```

LSP,LSP1,LSP2,LSP3,LSP4, (* THESE ARE STP VARIABLES *)
SAV1,SAV2,SAV3,SAV4,LSP5, (* USED FOR MULTI PURPOSES *)
FSP,FSP1,RESULT,SAV10, (* RESULT IS USED TO KEEP *)
SAVEDTAG,FRECVAR : STP ; (* THE GOAL POINTER *)
SAV6 : STP ; (* TO SAVE STP POINTERS *)
SAV5,SAV11,SAV12, (* *)
SAV7,SAV8,SAV9 : IDP ; (* TO SAVE IDP POINTERS *)
LVALU,FVALU : CONSTANT ; (* TO SAVE CONSTANTS *)
ARRIND : (* TO INDEX ARRAYST *)
LMIN : (* FOR LOWER BOUNDS *)
LMAX : (* FOR UPPER BOUNDS *)
LCNT : (* INDEX OF USER SCALARS *)
VNO,INO : (* VARIABLE COUNTERS *)
PCASEIND : (* RECORD STACK INDEX *)
OLDS : (* TO SAVE TOP VALUES *)
RECINX : INTEGER ; (* TO INDEX RECST *)
LCP,LCP1,LCP2,LCP3, (* MULTIPLE USE IDENTIFIER *)
FCP,FCP1 : IDP ; (* TABLE POINTERS *)
ALLOWDOTS : (* ..300LEANS.. *)
(* .. ALLOWED ? *)

```

```

EMPTY          : (* EMPTY FIELD LIST ? *)
FORM          : (* IS FORWARD ? *)
PRTER        : (* TO SUPPRESS ERROR MESSAGES *)
PRINTABLE     : (* WHEN NECESSARY *)
KOLON        : (* SYMBOL TABLE DUMP REQUIRED? *)
PCKD         : (* : MET ? *)
LID          : BOOLEAN ; (* PACKED STRUCTURE ? *)
PROGNAME     : (* TEMPORARY IDENTIFIER *)
IDENTIFIER   : PACKEDID ; (* NAME FROM PROGRAM HEADING *)
              : (* KEEPERS *)

FWPTR        : IDP ; (* ***** POINTER HEADS *)
              : (* HEAD OF CHAIN OF FORM DECL *)
LABPTR       : (* TYPE IDS *)
FSTLABPTR    : LABELPTR ; (* TO ALLOCATE LABEL *)
              : (* HEAD OF LABEL CHAIN *)
OLDLEV      : (* ***** DCL LEVELS *)
LEVEL       : LEVRANGE ; (* TO SAVE LEVEL COUNTER *)
TTOP, TOP, OLDTOP, DISK : DISPRANGE ; (* CURRENT STATIC LEVEL *)
              : (* TO INDEX DISPLAY *)
              : (* ***** STATISTICS ***** *)
GUNIT       : UNIT ; (* DESCRIBES THE CURRENT UNIT *)
STARTLINE   : INTEGER ; (* START OF CURRENT UNIT *)
              : (* ***** *)
INTPTR, REALPTR, BOOLPTR, CHARPTR, TEXTPTR, NILPTR : STP ; (* ***** STRUCTURE ***** *)
              : (* ***** TABLE ***** *)
              : (* ***** CONSTANTS ***** *)
INTID, REALID, TRUEID, MAXINTID, CHARID, BOOLID, FALSEID, TEXTID : IDP ; (* ***** IDENTIFIER ***** *)
              : (* ***** TABLE ***** *)
              : (* ***** CONSTANTS ***** *)
UTYPPTR     : (* UNDEFINED TYPE POINTER *)
UVARPTR     : (* UNDEFINED VAR POINTER *)
UKONSPTR    : (* UNDEFINED KONS POINTER *)
UFLDPTTR    : (* UNDEFINED FIELD POINTER *)
UPPCPTR     : (* UNDEFINED PROC POINTER *)
UFCTPTR     : IDP ; (* UNDEFINED FUNC POINTER *)

FSY, FSYL    : SYMBOL ;
LFORM       : FORMWFPTR ; (* TEMPORARY *)
OLDTOPS     : ARRAY [1..30] OF DISPRANGE ;
DISPLAY     :
  ARRAY [DISPRANGE] OF
    PACKED RECORD
      FNAME : IDP ;
      OCCUR : WHERE
    END ;
ARRAYST     : ARRAY [1..20] OF STP ;
RECST      : ARRAY [1..20] OF
  RECORD
    FST, LCP : IDP
  END ;
SIGN       : (NOSIGN, NEG, POS) ;
LKIND     : IDKIND ;
PCASES    : ARRAY [1..20] OF
  RECORD
    FPCKD : BOOLEAN ;
    TAGDEFINITION : STP ;
    FIRSTFIELD : IDP ;
    VARHEAD : STP ;
    INTREC : STP
  END ;
IDLIST    : (USERSCALAR, PFVARPAR, PEPARAM, FILEHEADS) ;
ROUTINES  : ARRAY [0..20] OF
  RECORD
    OLDLEV : LEVRANGE ;
    OLDTOP : DISPRANGE
  END ;

```

(*****
*
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*

P R O C E E D U R E S

```

PROCEDURE STACKDUMP ;
  VAR M : LINKS ;
BEGIN
  WRITELN ('          **DUMP** ') ;
  M := STACKTOP ; WHILE M <> NIL DO
  BEGIN WRITELN ('          ', M @ PTR) ; M := M @ PREVIOUS END ;
  WRITELN ('          ****END DUMP ****') END ;
PROCEDURE INITSCAN ;
  VAR I, J : INTEGER ;
PROCEDURE INITRANGE (VAR FILL : ARRAY [CHSIZE] OF INTEGER ;
                    LOW, HIGH, VALUE : CHSIZE ) ;
  VAR I, K, J : CHSIZE ;
BEGIN
  J := VALUE ;
  FOR I := LOW TO HIGH DO
  BEGIN
    FILL [I] := J ;
    J := J + 1
  END
END ; (* INITRANGE *)
BEGIN
  ACH := ACHR (97) ;
  FOR CH := 'A' TO 'Z' DO BEGIN
  LCASE [CH] := ACH ; ACH := SUCC (ACH) END ;
  FOR I := 0 TO 5 DO
  BEGIN
    NEXT [I] := 01 ; BASE [I] := 0 ;
    CHECK [I] := 01
  END ;
NEXT [ 6 ] := 68 ; NEXT [ 7 ] := 74 ; NEXT [ 8 ] := 99 ;
NEXT [ 9 ] := 65 ; NEXT [ 10 ] := 70 ; NEXT [ 11 ] := 76 ;
NEXT [ 12 ] := 71 ; NEXT [ 13 ] := 90 ; NEXT [ 14 ] := 72 ;
NEXT [ 15 ] := 92 ; NEXT [ 16 ] := 95 ; NEXT [ 17 ] := 94 ;
NEXT [ 18 ] := 101 ; NEXT [ 19 ] := 54 ; NEXT [ 20 ] := 73 ;
NEXT [ 21 ] := 77 ; NEXT [ 22 ] := 78 ; NEXT [ 23 ] := 66 ;
NEXT [ 24 ] := 67 ; NEXT [ 25 ] := 75 ; NEXT [ 26 ] := 79 ;
NEXT [ 27 ] := 80 ; NEXT [ 28 ] := 81 ; NEXT [ 29 ] := 82 ;
NEXT [ 30 ] := 69 ; NEXT [ 31 ] := 84 ; NEXT [ 32 ] := 85 ;
NEXT [ 33 ] := 86 ; NEXT [ 34 ] := 83 ; NEXT [ 35 ] := 87 ;
NEXT [ 36 ] := 88 ; NEXT [ 37 ] := 89 ; NEXT [ 38 ] := 91 ;
NEXT [ 39 ] := 93 ; NEXT [ 40 ] := 97 ; NEXT [ 41 ] := 98 ;
NEXT [ 42 ] := 100 ; NEXT [ 43 ] := 103 ; NEXT [ 44 ] := 104 ;
NEXT [ 45 ] := 96 ; NEXT [ 51 ] := 102 ; NEXT [ 62 ] := 113 ;
NEXT [ 63 ] := 122 ;
INITRANGE (NEXT, 46, 50, 105) ;
INITRANGE (NEXT, 52, 54, 110) ;
INITRANGE (NEXT, 55, 57, 114) ;
INITRANGE (NEXT, 58, 61, 118) ;
INITRANGE (NEXT, 64, 68, 124) ;
INITRANGE (NEXT, 69, 73, 130) ;
NEXT [ 74 ] := 136 ; NEXT [ 75 ] := 123 ; NEXT [ 76 ] := 171 ;
NEXT [ 81 ] := 129 ; NEXT [ 94 ] := 155 ; NEXT [ 95 ] := 145 ;
NEXT [ 96 ] := 157 ; NEXT [ 97 ] := 158 ; NEXT [ 98 ] := 154 ;
INITRANGE (NEXT, 77, 80, 137) ;
INITRANGE (NEXT, 82, 85, 141) ;
INITRANGE (NEXT, 86, 93, 146) ;
INITRANGE (NEXT, 99, 104, 159) ;
NEXT [ 105 ] := 168 ; NEXT [ 106 ] := 165 ; NEXT [ 107 ] := 166 ;
NEXT [ 108 ] := 155 ; NEXT [ 109 ] := 167 ; NEXT [ 110 ] := 169 ;
NEXT [ 111 ] := 170 ; NEXT [ 112 ] := 135 ; NEXT [ 113 ] := 117 ;
FOR I := 114 TO 168 DO
  BEGIN

```

NEXT [I] := I + 1 ;

END ;

CHECK	6	67	CHECK	7	8	CHECK	8	98
CHECK	9	64	CHECK	10	7	CHECK	11	75
CHECK	12	70	CHECK	13	89	CHECK	14	71
CHECK	15	91	CHECK	16	94	CHECK	17	93
CHECK	18	00	CHECK	19	6	CHECK	20	72
CHECK	21	00	CHECK	22	77	CHECK	23	6
CHECK	24	65	CHECK	25	74	CHECK	26	78
CHECK	27	79	CHECK	28	99	CHECK	29	81
CHECK	30	68	CHECK	31	83	CHECK	32	84
CHECK	33	95	CHECK	34	9	CHECK	35	86
CHECK	36	10	CHECK	37	88	CHECK	38	10
CHECK	39	11	CHECK	40	95	CHECK	41	97
CHECK	42	99	CHECK	43	102	CHECK	44	73
CHECK	45	1	CHECK	50	12	CHECK	51	11

INITRANGE (CHECK, 46, 49, 104) ;
 INITRANGE (CHECK, 65, 68, 124) ;
 INITRANGE (CHECK, 69, 73, 129) ;
 INITRANGE (CHECK, 77, 79, 136) ;
 INITRANGE (CHECK, 82, 85, 140) ;
 INITRANGE (CHECK, 86, 88, 145) ;
 INITRANGE (CHECK, 99, 101, 158) ;

CHECK	52	109	CHECK	53	110	CHECK	54	14
CHECK	55	17	CHECK	56	114	CHECK	57	15
CHECK	58	18	CHECK	59	119	CHECK	60	19
CHECK	61	120	CHECK	62	114	CHECK	63	20
CHECK	64	124	CHECK	74	130	CHECK	75	20
CHECK	76	134	CHECK	80	23	CHECK	81	21
CHECK	89	154	CHECK	90	149	CHECK	91	25
CHECK	92	154	CHECK	93	152	CHECK	94	55
CHECK	95	140	CHECK	96	155	CHECK	97	26
CHECK	98	250	CHECK	102	27	CHECK	103	62
CHECK	104	288	CHECK	105	29	CHECK	106	64
CHECK	107	55	CHECK	108	25	CHECK	109	66
CHECK	110	168	CHECK	111	169	CHECK	112	171
CHECK	113	116						

BASE	6	0	BASE	7	0	BASE	8	1
BASE	9	14	BASE	10	19	BASE	11	25
BASE	12	30	BASE	13	0	BASE	14	43
BASE	15	0	BASE	16	0	BASE	17	49
BASE	18	38	BASE	19	40	BASE	20	52
BASE	21	58	BASE	22	0	BASE	23	70
BASE	24	79	BASE	25	73	BASE	26	78
BASE	27	96	BASE	28	91			

FOR I := 29 TO 171 DO
 BASE [I] := 0 ;

BASE	66	1	BASE	74	1	BASE	75	1
BASE	77	31	BASE	78	2	BASE	79	12
BASE	81	3	BASE	83	3	BASE	84	13
BASE	85	28	BASE	86	15	BASE	88	13
BASE	89	33	BASE	91	16	BASE	94	16
BASE	96	17	BASE	97	13	BASE	98	2
BASE	99	19	BASE	100	9	BASE	102	24
BASE	103	36	BASE	104	21	BASE	105	33
BASE	106	55	BASE	107	30	BASE	109	27
BASE	110	33	BASE	111	49	BASE	112	47
BASE	116	96	BASE	118	50	BASE	120	36
BASE	124	57	BASE	125	50	BASE	126	57
BASE	127	59	BASE	129	49	BASE	130	62
BASE	131	61	BASE	132	53	BASE	133	47
BASE	134	53	BASE	136	54	BASE	137	72
BASE	138	61	BASE	140	74	BASE	141	63
BASE	142	61	BASE	143	76	BASE	145	76
BASE	146	81	BASE	147	53	BASE	149	65
BASE	151	82	BASE	152	74	BASE	155	73
BASE	156	84	BASE	158	74	BASE	159	86
BASE	160	84	BASE	162	80	BASE	164	92
BASE	166	90	BASE	166	99	BASE	168	95
BASE	169	98	BASE	171	102	BASE	172	1

FOR I := 1 TO 171 DO
 WITH DEFAULT [I] DO CLASS := IDENT ;

WITH DEFAULT [65] DO
 BEGIN CLASS := MULOP ; OP := ANDOP END ;

```

WITH DEFAULT [ 69 ] DO CLASS ::= ARRAYSY
WITH DEFAULT [ 73 ] DO CLASS ::= BEGINSY
WITH DEFAULT [ 76 ] DO CLASS ::= CASESY
WITH DEFAULT [ 80 ] DO CLASS ::= CONSTSY

WITH DEFAULT [ 82 ] DO
  BEGIN CLASS ::= MULOP ; OP ::= IDIV  END

WITH DEFAULT [ 83 ] DO CLASS ::= DOSY
WITH DEFAULT [ 87 ] DO CLASS ::= DOWNTOSY
WITH DEFAULT [ 90 ] DO CLASS ::= ELSESY
WITH DEFAULT [ 92 ] DO CLASS ::= ENDSY
WITH DEFAULT [ 95 ] DO CLASS ::= FILES Y
WITH DEFAULT [ 97 ] DO CLASS ::= FORSY
WITH DEFAULT [ 101 ] DO CLASS ::= FORWARDSY
WITH DEFAULT [ 108 ] DO CLASS ::= FUNCTIONSY
WITH DEFAULT [ 111 ] DO CLASS ::= GOTOSY
WITH DEFAULT [ 112 ] DO CLASS ::= IFSY

WITH DEFAULT [ 113 ] DO
  BEGIN CLASS ::= RELOP ; OP ::= INOP  END

WITH DEFAULT [ 119 ] DO
  BEGIN CLASS ::= MULOP ; OP ::= IMOD  END

WITH DEFAULT [ 123 ] DO
  BEGIN CLASS ::= ADDOP ; OP ::= OROP  END

WITH DEFAULT [ 117 ] DO CLASS ::= LABELSY
WITH DEFAULT [ 121 ] DO CLASS ::= NOTSY
WITH DEFAULT [ 122 ] DO CLASS ::= OPSY
WITH DEFAULT [ 128 ] DO CLASS ::= PACKEDSY
WITH DEFAULT [ 135 ] DO CLASS ::= PROCSY
WITH DEFAULT [ 139 ] DO CLASS ::= PROGRAMSY
WITH DEFAULT [ 144 ] DO CLASS ::= RECORDSY
WITH DEFAULT [ 148 ] DO CLASS ::= REPEATSY
WITH DEFAULT [ 150 ] DO CLASS ::= SETSY
WITH DEFAULT [ 153 ] DO CLASS ::= THEMSY
WITH DEFAULT [ 154 ] DO CLASS ::= TOSY
WITH DEFAULT [ 157 ] DO CLASS ::= TYPESY
WITH DEFAULT [ 161 ] DO CLASS ::= UNTILSY
WITH DEFAULT [ 163 ] DO CLASS ::= VARSY
WITH DEFAULT [ 167 ] DO CLASS ::= WHILESY
WITH DEFAULT [ 170 ] DO CLASS ::= WITHSY

FOR I := 1 TO STRGLGTH DO
  DUMMYSYSTR [ I ] := ' ' ;

FOR I := 1 TO CHARMAX DO
  DUMMYID [ I ] := ' ' ;

DIGITS := [ '0'..'9' ] ;
IDCHARS := DIGITS + [ 'A'..'Z' ] ;

CH := FIRSTCHAR ;
REPEAT
  WITH CHRCLASS [ CH ] DO
    CLASS ::= OTHERSY ;
  CH := SUCC ( CH )
UNTIL CH = LASTCHAR ;
CHRCLASS [ CH ] CLASS ::= OTHERSY ;

CHRCLASS [ ',' ] CLASS ::= COMMA
CHRCLASS [ '(' ] CLASS ::= RPARENT
CHRCLASS [ '.' ] CLASS ::= PERIOD
CHRCLASS [ '[' ] CLASS ::= LBRACKET
CHRCLASS [ ']' ] CLASS ::= RBRACKET
CHRCLASS [ '>' ] CLASS ::= ARROW
CHRCLASS [ '<' ] CLASS ::= ARROW
CHRCLASS [ ';' ] CLASS ::= SEMICOLON

WITH CHRCLASSE '+' DO
  BEGIN CLASS ::= ADDOP ; OP ::= PLUS  END ;

WITH CHRCLASSE '-' DO
  BEGIN CLASS ::= ADDOP ; OP ::= MINUS  END ;

WITH CHRCLASSE '*' DO
  BEGIN CLASS ::= MULOP ; OP ::= ASTR  END ;

WITH CHRCLASSE '=' DO
  BEGIN CLASS ::= RELOP ; OP ::= EQOP  END ;

```

```

WITH CHRCLASSE '/' DO
  BEGIN CLASS := MULOP ; OP := RDIV END ;

LASTERR := 0 ; ERRPOS := 0
ERRINX := 0 ; CHCNT := BUFMAX ;
CARDCNT := 0 ; CH := ;
UNCLOSED := NONE ;
BUFLGTH := 0 ;

FOR I := 1 TO BUFMAX DO
  ERRBUF [ I ] := ' ' ;

ERRWARN [ TRUE ] := ' *** ERROR ' ;
ERRWARN [ FALSE ] := ' *** WARNING ' ;

ABORTMES [ TRUE ] := ' *** UNCLOSED COMMENT *** ' ;
ABORTMES [ FALSE ] := ' *** UNCLOSED STRING *** ' ;

ERRWARNCNT [ TRUE ] := 0 ;
ERRWARNCNT [ FALSE ] := 0 ;

FOR I := 1 TO NSETS DO
  EMPTYSET [ I ] := [] ;
LNEST := FALSE ; RNEST := FALSE ;
ERRDUPL := FALSE ;
NEST := 0 ; NESTINCR := 0 ; NESTDECR := 0
END ; (* INITSCAN *)

```

```

          SPAGE
PROCEDURE INITPARSE
  VAR I : INTEGER
      INITIAL : SYMBOL
;
PROCEDURE INITSUCC (L,H : SPRANGE) ;
  VAR TEMP : SPRANGE
;
  BEGIN
    FOR TEMP := L TO H DO
      WITH STABLE [ TEMP ] DO
        SUC := TEMP
      ;
    ;
  END ;

```

```

BEGIN

```

```

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*
*
*
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```

```

          INIT GRAMMAR TABLES

```

```

*****
*
*
*
*****

```

```

  FOR I := 1 TO SPSIZE DO
    WITH STABLE [ I ] DO
      ALT := FAIL
    ;
  ;

```

```

(*

```

```

*)

```

```

1) <PROGRAM> ::= <PROGRAM HEADING> <BODY> ;

```

```

  INITSUCC ( 2, 3)
;

```

```

  WITH STABLE [ 1 ] DO TPTYPE := 53 ;
  WITH STABLE [ 2 ] DO TPTYPE := 55 ;
  WITH STABLE [ 3 ] DO
    BEGIN SUC := OK; TPTYPE := ORD (PERIOD ) END ;
;

```

```

(*

```

```

*)

```

```

2) <PROGRAM HEADING> ::= PROGRAM IDENT ( <IDENT LIST> ) ;

```

```

  INITSUCC ( 5, 9)
;

```

```

  WITH STABLE [ 4 ] DO TPTYPE := ORD (PROGRAMSY ) ;
  WITH STABLE [ 5 ] DO TPTYPE := ORD (IDENT ) ;
  WITH STABLE [ 6 ] DO TPTYPE := ORD (LPARENT ) ;
  WITH STABLE [ 7 ] DO TPTYPE := 54 ;
  WITH STABLE [ 8 ] DO TPTYPE := ORD (RPARENT ) ;
  WITH STABLE [ 9 ] DO
    BEGIN SUC := OK; TPTYPE := ORD (SEMICOLON ) END ;
;

```

```

(*

```

```

*)

```

```

3) <IDENT LIST> ::= IDENT / , IDENT / 0/

```

```

  WITH STABLE [ 10 ] DO
    BEGIN SUC := 11; TPTYPE := ORD (IDENT ) END ;
  WITH STABLE [ 11 ] DO
    BEGIN SUC := 12; TPTYPE := ORD (COMMA ) ;
      ALT := OK
    ;
  END ;
  WITH STABLE [ 12 ] DO
    BEGIN SUC := 11; TPTYPE := ORD (IDENT ) END ;
;

```

```

(*

```

```

*)

```

```

4) <BODY> ::= <DECLARATIONS> <ROUTINE DECLARATIONS>
          <COMPOUND STATEMENT>

```

```

  INITSUCC ( 14, 15)
;

```

```

  WITH STABLE [ 13 ] DO TPTYPE := 56 ;
  WITH STABLE [ 14 ] DO TPTYPE := 71 ;
  WITH STABLE [ 15 ] DO
    BEGIN SUC := OK; TPTYPE := 75 ;
  END ;
;

```

```

(*

```

```

5) <DECLARATIONS> ::= <LABEL DECLARATIONS>
                   <CONSTANT DEFINITIONS>
                   <TYPE DEFINITIONS>

```


<VAR DECLARATIONS>

*)

```

INITSUCC ( 17, 19)
WITH STABLE C 16J DO TPTYPE := 57
WITH STABLE C 17J DO TPTYPE := 58
WITH STABLE C 18J DO TPTYPE := 59
WITH STABLE C 19J DO
  BEGIN SUC := OK; TPTYPE := 70
END

```

(*

```

6) <LABEL DECLARATIONS>
   ::= LABEL /INTCONST/ , INTCONST/O/ ; /1/
   <EMPTY>

```

*)

```

INITSUCC ( 21, 23)
WITH STABLE C 20J DO
  BEGIN ALT := OK; TPTYPE := ORD (LABELSY ) END
WITH STABLE C 21J DO TPTYPE := ORD (INTCONST )
WITH STABLE C 22J DO
  BEGIN ALT := 24; TPTYPE := ORD (COMMA ) END
WITH STABLE C 23J DO
  BEGIN SUC := 22; TPTYPE := ORD (INTCONST ) END
WITH STABLE C 24J DO
  BEGIN SUC := OK; TPTYPE := ORD (SEMICOLON ) END

```

(*

```

7) <CONSTANT DEFINITIONS>
   ::= CONST / IDENT = <CONSTANT>/1/
   <EMPTY>

```

*)

```

INITSUCC ( 26, 32)
WITH STABLE C 25J DO
  BEGIN ALT := OK; TPTYPE := ORD (CONSTSY ) END
WITH STABLE C 26J DO TPTYPE := ORD (IDENT )
WITH STABLE C 27J DO TPTYPE := ORD (RELOP )
WITH STABLE C 28J DO TPTYPE := 59
WITH STABLE C 29J DO TPTYPE := ORD (SEMICOLON )
WITH STABLE C 30J DO
  BEGIN ALT := OK; TPTYPE := ORD (IDENT ) END
WITH STABLE C 31J DO TPTYPE := ORD (RELOP )
WITH STABLE C 32J DO
  BEGIN SUC := 29; TPTYPE := 59
END

```

(*

```

8) <CONSTANT> ::= INTCONST REALCONST IDENT
                ADDOP INTCONST ADDOP REALCONST
                ADDOP IDENT STRINGCONST
                CHARCONST

```

*)

```

WITH STABLE C 33J DO
  BEGIN SUC := 34; TPTYPE := ORD (ADDOP )
  ALT := 37
END
FOR I := 34 TO 41 DO
  WITH STABLE C IJ DO SUC := OK
  BEGIN ALT := 35; TPTYPE := ORD (INTCONST ) END
  WITH STABLE C 35J DO
    BEGIN ALT := 36; TPTYPE := ORD (REALCONST ) END
  WITH STABLE C 36J DO TPTYPE := ORD (IDENT )
  WITH STABLE C 37J DO
    BEGIN ALT := 38; TPTYPE := ORD (IDENT ) END
  WITH STABLE C 38J DO
    BEGIN ALT := 39; TPTYPE := ORD (REALCONST ) END
  WITH STABLE C 39J DO
    BEGIN ALT := 40; TPTYPE := ORD (CHARCONST ) END
  WITH STABLE C 40J DO
    BEGIN ALT := 41; TPTYPE := ORD (INTCONST ) END
  WITH STABLE C 41J DO TPTYPE := ORD (STRINGCONST)

```

(*

```

9) <TYPE DEFINITIONS>
   ::= TYPE / IDENT = <TYPE> ; /1/ <EMPTY>

```

*)

```

INITSUCC ( 43, 49)

```

```

WITH STABLE C 42] DO
  BEGIN ALT ::= OK; TPTYPE ::= ORD (TYPESY ) END ;
WITH STABLE C 43] DO TPTYPE ::= ORD (IDENT )
WITH STABLE C 44] DO TPTYPE ::= ORD (RELOP )
WITH STABLE C 45] DO TPTYPE ::= 51
WITH STABLE C 46] DO TPTYPE ::= ORD (SEMICOLON )
WITH STABLE C 47] DO
  BEGIN ALT ::= OK; TPTYPE ::= ORD (IDENT ) END ;
WITH STABLE C 48] DO TPTYPE ::= ORD (RELOP )
WITH STABLE C 49] DO
  BEGIN SUC ::= 46; TPTYPE ::= 51 END ;

```

```

(*) 10) <TYPE > ::= IDENT
      PACKED SET OF <SIMPLE TYPE>
      SET OF <SIMPLE TYPE>
      PACKED ARRAY [ <INDEX LIST> ] OF <TYPE>
      ARRAY [ <INDEX LIST> ] OF <TYPE>
      PACKED RECORD <FIELD LIST> END
      RECORD <FIELD LIST> END
      PACKED FILE OF <TYPE>
      FILE OF <TYPE>
      <SIMPLE TYPE>

```

(*)

```

WITH STABLE C 50] DO
  BEGIN ALT ::= 52; TPTYPE ::= ORD (ARROW )
  SUC ::= 51
  END ;
WITH STABLE C 51] DO
  BEGIN SUC ::= OK; TPTYPE ::= ORD (IDENT ) END ;
WITH STABLE C 52] DO
  BEGIN SUC ::= 53; TPTYPE ::= ORD (PACKEDSY )
  ALT ::= 57
  END ;
WITH STABLE C 53] DO
  BEGIN SUC ::= 58; TPTYPE ::= ORD (ARRAYSY )
  ALT ::= 54
  END ;
WITH STABLE C 54] DO
  BEGIN SUC ::= 64; TPTYPE ::= ORD (RECORDSY )
  ALT ::= 55
  END ;
WITH STABLE C 55] DO
  BEGIN SUC ::= 67; TPTYPE ::= ORD (SETSY )
  ALT ::= 56
  END ;
WITH STABLE C 56] DO
  BEGIN SUC ::= 70; TPTYPE ::= ORD (FILESY ) END ;

INITSUCC ( 58, 62)
WITH STABLE C 57] DO
  BEGIN ALT ::= 63; TPTYPE ::= ORD (ARRAYSY ) END ;
WITH STABLE C 58] DO TPTYPE ::= ORD (LBRACKET )
WITH STABLE C 59] DO TPTYPE ::= 52
WITH STABLE C 60] DO TPTYPE ::= ORD (RBRACKET )
WITH STABLE C 61] DO TPTYPE ::= ORD (OFSY )
WITH STABLE C 62] DO
  BEGIN SUC ::= OK; TPTYPE ::= 51 END ;

INITSUCC ( 64, 65)
WITH STABLE C 63] DO
  BEGIN ALT ::= 66; TPTYPE ::= ORD (RECORDSY ) END ;
WITH STABLE C 64] DO TPTYPE ::= 54
WITH STABLE C 65] DO
  BEGIN SUC ::= OK; TPTYPE ::= ORD (ENDSY ) END ;

INITSUCC ( 67, 68)
WITH STABLE C 66] DO
  BEGIN ALT ::= 69; TPTYPE ::= ORD (SETSY ) END ;
WITH STABLE C 67] DO TPTYPE ::= ORD (OFSY )
WITH STABLE C 68] DO
  BEGIN SUC ::= OK; TPTYPE ::= 53 END ;

INITSUCC ( 70, 71)
WITH STABLE C 69] DO
  BEGIN ALT ::= 72; TPTYPE ::= ORD (FILESY ) END ;
WITH STABLE C 70] DO TPTYPE ::= ORD (OFSY )

```

WITH STABLE C 71J DO
BEGIN SUC := OK; TPTYPE := 51 END ;

WITH STABLE C 72J DO
BEGIN SUC := OK; TPTYPE := 53 END ;

(*
*) 11) <INDEX LIST> ::= <SIMPLE TYPE> /, <SIMPLE TYPE>/O/

WITH STABLE C 73J DO
BEGIN SUC := 74; TPTYPE := 53 END ;

WITH STABLE C 74J DO
BEGIN SUC := 75; TPTYPE := ORD (COMMA) ;
ALT := OK

END ;
WITH STABLE C 75J DO
BEGIN SUC := 74; TPTYPE := 53 END ;

(*
*) 12) <SIMPLE TYPE>
::= (<IDENT LIST>) <CONSTANT>
<CONSTANT> * <CONSTANT>

INITSUCC (77, 82) ;

WITH STABLE C 76J DO
BEGIN ALT := 80; TPTYPE := ORD (LPARENT) END ;

WITH STABLE C 77J DO
TPTYPE := 54 ;

WITH STABLE C 78J DO
BEGIN SUC := OK; TPTYPE := ORD (RPARENT) END ;

WITH STABLE C 79J DO
BEGIN SUC := OK; TPTYPE := ORD (IDENT) ;
ALT := 80

END ;
WITH STABLE C 80J DO
TPTYPE := 59 ;

WITH STABLE C 81J DO
BEGIN ALT := OK; TPTYPE := ORD (DOTDOT) END ;

WITH STABLE C 82J DO
BEGIN SUC := OK; TPTYPE := 59 END ;

(*
*) 13) <FIELD LIST> ::= IDENT /, IDENT/O/ : <TYPE>
/ ; <IDENT LIST> : <TYPE> /O/
CASE <VARIANT PART>
/IDENT /, IDENT/ : <TYPE> ; /
CASE <VARIANT PART>

(*
*) INITSUCC (84, 91) ;

WITH STABLE C 83J DO
BEGIN ALT := 90; TPTYPE := ORD (IDENT) END ;

WITH STABLE C 84J DO
BEGIN ALT := 86; TPTYPE := ORD (COMMA) END ;

WITH STABLE C 85J DO
BEGIN SUC := 84; TPTYPE := ORD (IDENT) END ;

WITH STABLE C 86J DO
TPTYPE := ORD (COLON) ;

WITH STABLE C 87J DO
TPTYPE := 51 ;

WITH STABLE C 88J DO
BEGIN ALT := OK; TPTYPE := ORD (SEMICOLON) END ;

WITH STABLE C 89J DO
BEGIN SUC := 84; TPTYPE := ORD (IDENT) ;
ALT := 90

END ;
WITH STABLE C 90J DO
TPTYPE := ORD (CASESY) ;

WITH STABLE C 91J DO
BEGIN SUC := OK; TPTYPE := 55 END ;

(*
*) 14) <VARIANT PART>
::= IDENT : IDENT OF <VARIANT LIST>
IDENT OF <VARIANT LIST>

(*
*) INITSUCC (93, 96) ;

WITH STABLE C 92J DO
TPTYPE := ORD (IDENT) ;

WITH STABLE C 93J DO
BEGIN ALT := 95; TPTYPE := ORD (COLON) END ;

WITH STABLE C 94J DO
TPTYPE := ORD (IDENT) ;

```

WITH STABLE [ 95] DO TPTYPE := ORD (OFSY ) ;
WITH STABLE [ 96] DO
  BEGIN SUC := OK; TPTYPE := 55 END ;

```

```

(*)
15) <VARIANT LIST> ::= <VARIANT> / ; <VARIANT>/O/
*)

```

```

INITSUCC ( 98, 99) ;
WITH STABLE [ 97] DO TPTYPE := 57 ;
WITH STABLE [ 98] DO
  BEGIN ALT := OK; TPTYPE := ORD (SEMICOLON ) END ;
WITH STABLE [ 99] DO
  BEGIN SUC := 98; TPTYPE := 57 END ;

```

```

(*)
16) <VARIANT> ::= <CONSTANT LIST> : ( <FIELD PART>
*)

```

```

INITSUCC (101,103) ;
WITH STABLE [100] DO TPTYPE := 59 ;
WITH STABLE [101] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [102] DO TPTYPE := ORD (LPARENT ) ;
WITH STABLE [103] DO
  BEGIN SUC := OK; TPTYPE := 58 END ;

```

```

(*)
17) <FIELD PART> ::= ) <FIELD LIST> )
*)

```

```

WITH STABLE [104] DO
  BEGIN SUC := OK; TPTYPE := ORD (RPARENT ) ;
  ALT := 105 ;
  END ;
WITH STABLE [105] DO
  BEGIN SUC := 106; TPTYPE := 54 END ;
WITH STABLE [106] DO
  BEGIN SUC := OK; TPTYPE := ORD (RPARENT ) END ;

```

```

(*)
18) <CONSTANT LIST> ::= <CONSTANT> /, <CONSTANT>/O/
*)

```

```

INITSUCC (108,109) ;
WITH STABLE [107] DO TPTYPE := 59 ;
WITH STABLE [108] DO
  BEGIN ALT := OK; TPTYPE := ORD (COMMA ) END ;
WITH STABLE [109] DO
  BEGIN SUC := 108; TPTYPE := 59 END ;

```

```

(*)
19) <VAR DECLARATIONS> ::= VAR /IDENT /, IDENT/O/ : <TYPE> ; /I/
   ::= <EMPTY>
*)

```

```

INITSUCC (111,121) ;
WITH STABLE [110] DO
  BEGIN ALT := OK; TPTYPE := ORD (VARSY ) END ;
WITH STABLE [111] DO TPTYPE := ORD (IDENT ) ;
WITH STABLE [112] DO
  BEGIN ALT := 114; TPTYPE := ORD (COMMA ) END ;
WITH STABLE [113] DO
  BEGIN SUC := 112; TPTYPE := ORD (IDENT ) END ;
WITH STABLE [114] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [115] DO TPTYPE := 51 ;
WITH STABLE [116] DO TPTYPE := ORD (SEMICOLON ) ;
WITH STABLE [117] DO
  BEGIN ALT := OK; TPTYPE := ORD (IDENT ) END ;
WITH STABLE [118] DO
  BEGIN ALT := 120; TPTYPE := ORD (COMMA ) END ;
WITH STABLE [119] DO
  BEGIN SUC := 118; TPTYPE := ORD (IDENT ) END ;
WITH STABLE [120] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [121] DO
  BEGIN SUC := 116; TPTYPE := 51 END ;

```

```

(*)
20) <ROUTINE DECLARATIONS>

```

```

::= PROCEDURE IDENT
    <FORMAL PARAMETER SECTION> ; FORWARD ;
PROCEDURE IDENT
    <FORMAL PARAMETER SECTION> ; <BODY> ;
FUNCTION IDENT
    <FORMAL PARAMETER SECTION> : IDENT ;
FORWARD ;
FUNCTION IDENT
    <FORMAL PARAMETER SECTION> : IDENT ;
<BODY> ;

```

(*)

```

INITSUCC (123,131) ;
WITH STABLE [122] DO ;
  BEGIN ALT := 125; TPTYPE := ORD (PROCSY ) END ;
WITH STABLE [123] DO TPTYPE := ORD (IDENT ) ;
WITH STABLE [124] DO ;
  BEGIN SUC := 130; TPTYPE := 72 END ;
WITH STABLE [125] DO ;
  BEGIN ALT := OK; TPTYPE := ORD (FUNCTIONSY ) END ;
WITH STABLE [126] DO TPTYPE := ORD (IDENT ) ;
WITH STABLE [127] DO TPTYPE := 72 ;
WITH STABLE [128] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [129] DO TPTYPE := ORD (IDENT ) ;
WITH STABLE [130] DO TPTYPE := ORD (SEMICOLON ) ;
WITH STABLE [131] DO ;
  BEGIN SUC := 133; TPTYPE := ORD (FORWARDSY ) ;
  ALT := 132 ;
END ;
WITH STABLE [132] DO ;
  BEGIN SUC := 133; TPTYPE := 55 END ;
WITH STABLE [133] DO ;
  BEGIN SUC := 122; TPTYPE := ORD (SEMICOLON ) END ;

```

(*)

```

21) <FORMAL PARAMETER SECTION>
    ::= ( <FORMAL PARAMETERS> ) <EMPTY>

```

(*)

```

INITSUCC (135,136) ;
WITH STABLE [134] DO ;
  BEGIN ALT := OK; TPTYPE := ORD (LPARENT ) END ;
WITH STABLE [135] DO TPTYPE := 73 ;
WITH STABLE [136] DO ;
  BEGIN SUC := OK; TPTYPE := ORD (RPARENT ) END ;

```

(*)

```

22) <FORMAL PARAMETERS>
    ::= <FORMAL PARAMETER> /;<FORMAL PARAMETER>/G/

```

(*)

```

INITSUCC (138,139) ;
WITH STABLE [137] DO TPTYPE := 74 ;
WITH STABLE [138] DO ;
  BEGIN ALT := OK; TPTYPE := ORD (SEMICOLON ) END ;
WITH STABLE [139] DO ;
  BEGIN SUC := 138; TPTYPE := 74 END ;

```

(*)

```

23) <FORMAL PARAMETER>
    ::= PROCEDURE <IDENT LIST>
        FUNCTION <IDENT LIST> : IDENT
        VAR <IDENT LIST> : <TYPE>
        <IDENT LIST> : <TYPE>

```

(*)

```

INITSUCC (141,149) ;
WITH STABLE [140] DO ;
  BEGIN ALT := 142; TPTYPE := ORD (PROCSY ) END ;
WITH STABLE [141] DO ;
  BEGIN SUC := OK; TPTYPE := 54 END ;
WITH STABLE [142] DO ;
  BEGIN ALT := 146; TPTYPE := ORD (FUNCTIONSY ) END ;
WITH STABLE [143] DO TPTYPE := 54 ;
WITH STABLE [144] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [145] DO ;
  BEGIN SUC := OK; TPTYPE := ORD (IDENT ) END ;
WITH STABLE [146] DO ;
  BEGIN ALT := 147; TPTYPE := ORD (VARSY ) END ;
WITH STABLE [147] DO TPTYPE := 54 ;
WITH STABLE [148] DO TPTYPE := ORD (COLON ) ;
WITH STABLE [149] DO ;

```

BEGIN SUC := OK; TPTYPE := 61 END ;

(* 24) <COMPOUND STATEMENT>
::= BEGIN <STATEMENT LIST> END

INITSUCC (151,152) ;
WITH STABLE [150] DO TPTYPE := ORD (BEGINSY) ;
WITH STABLE [151] DO TPTYPE := 76 ;
WITH STABLE [152] DO ;
BEGIN SUC := OK; TPTYPE := ORD (ENDSY) END ;

(* 24) <STATEMENT LIST>
::= <STATEMENT> / ; <STATEMENT>/O/

INITSUCC (154,155) ;
WITH STABLE [153] DO TPTYPE := 77 ;
WITH STABLE [154] DO ;
BEGIN ALT := OK; TPTYPE := ORD (SEMICOLON) END ;
WITH STABLE [155] DO ;
BEGIN SUC := 154; TPTYPE := 77 END ;

(* 26) <STATEMENT> ::= INTCONST : <UNLABELLED STATEMENT>
<UNLABELLED STATEMENT>

INITSUCC (157,158) ;
STABLE [156] ALT := 158 ;
STABLE [156] TPTYPE := ORD (INTCONST) ;
WITH STABLE [157] DO TPTYPE := ORD (COLON) ;
WITH STABLE [158] DO ;
BEGIN SUC := OK; TPTYPE := 78 END ;

(* 27) <UNLABELLED STATEMENT>
::= BEGIN <STATEMENT LIST> END
IF <EXPRESSION> THEN <STATEMENT>
IF <EXPRESSION> THEN <STATEMENT>
ELSE <STATEMENT>
CASE <EXPRESSION> OF <CASE LIST> END
WHILE <EXPRESSION> DO <STATEMENT>
REPEAT <STATEMENT LIST> UNTIL
<EXPRESSION>
FOR IDENT := <EXPRESSION> TO
<EXPRESSION> DO <STATEMENT>
FOR IDENT := <EXPRESSION> DOWNTO
<EXPRESSION> DO <STATEMENT>
WITH <RECORD VAR LIST> DO <STATEMENT>
<SIMPLE STATEMENT>

(* INITSUCC (160,193) ;
WITH STABLE [159] DO ;
BEGIN ALT := 162; TPTYPE := ORD (BEGINSY) END ;
WITH STABLE [160] DO TPTYPE := 76 ;
WITH STABLE [161] DO ;
BEGIN SUC := OK; TPTYPE := ORD (ENDSY) END ;
WITH STABLE [162] DO ;
BEGIN ALT := 168; TPTYPE := ORD (IFSY) END ;
WITH STABLE [163] DO TPTYPE := 84 ;
WITH STABLE [164] DO TPTYPE := ORD (THENSY) ;
WITH STABLE [165] DO TPTYPE := 77 ;
WITH STABLE [166] DO ;
BEGIN ALT := OK; TPTYPE := ORD (ELSESY) END ;
WITH STABLE [167] DO ;
BEGIN SUC := OK; TPTYPE := 77 END ;
WITH STABLE [168] DO ;
BEGIN ALT := 173; TPTYPE := ORD (CASESY) END ;
WITH STABLE [169] DO TPTYPE := 84 ;
WITH STABLE [170] DO TPTYPE := ORD (OFSY) ;
WITH STABLE [171] DO TPTYPE := 93 ;
WITH STABLE [172] DO ;
BEGIN ALT := OK; TPTYPE := ORD (ENDSY) END ;

```

WITH STABLE [173] DO
BEGIN ALT ::= 177; TPTYPE ::= ORD (WHILESY ) END ;
WITH STABLE [174] DO
WITH STABLE [175] DO
WITH STABLE [176] DO
BEGIN SUC ::= OK; TPTYPE ::= 77 END ;
WITH STABLE [177] DO
BEGIN ALT ::= 181; TPTYPE ::= ORD (REPEATSY ) END ;
WITH STABLE [178] DO
WITH STABLE [179] DO
WITH STABLE [180] DO
BEGIN SUC ::= OK; TPTYPE ::= 34 END ;
WITH STABLE [181] DO
BEGIN ALT ::= 190; TPTYPE ::= ORD (FORSY ) END ;
WITH STABLE [182] DO
WITH STABLE [183] DO
WITH STABLE [184] DO
WITH STABLE [185] DO
BEGIN SUC ::= 187; TPTYPE ::= ORD (TOSY ) ;
ALT ::= 186

END ;
WITH STABLE [186] DO
WITH STABLE [187] DO
WITH STABLE [188] DO
WITH STABLE [189] DO
BEGIN SUC ::= OK; TPTYPE ::= 77 END ;
WITH STABLE [190] DO
BEGIN ALT ::= 194; TPTYPE ::= ORD (WITHSY ) END ;
WITH STABLE [191] DO
WITH STABLE [192] DO
WITH STABLE [193] DO
BEGIN SUC ::= OK; TPTYPE ::= 77 END ;
WITH STABLE [194] DO
BEGIN SUC ::= OK; TPTYPE ::= 79 END ;

```

(*

28) <SIMPLE STATEMENT>

```

::= GOTO INTCONST
IDENT ( <EXPRESSION LIST> )
IDENT <VAR TAIL> ::= <EXPRESSION>
<EMPTY>

```

*)

INITSUCC (196,203)

```

WITH STABLE [195] DO
BEGIN ALT ::= 197; TPTYPE ::= ORD (GOTOSY ) END ;
WITH STABLE [196] DO
BEGIN SUC ::= OK; TPTYPE ::= ORD (INTCONST ) END ;
WITH STABLE [197] DO
BEGIN ALT ::= OK; TPTYPE ::= ORD (IDENT ) END ;
WITH STABLE [198] DO
BEGIN ALT ::= 201; TPTYPE ::= ORD (LPARENT ) END ;
WITH STABLE [199] DO
WITH STABLE [200] DO
BEGIN SUC ::= OK; TPTYPE ::= ORD (RPARENT ) END ;
WITH STABLE [201] DO
WITH STABLE [202] DO
BEGIN ALT ::= OK; TPTYPE ::= ORD (ASGNOP ) END ;
WITH STABLE [203] DO
BEGIN SUC ::= OK; TPTYPE ::= 84 END ;

```

(*

29) <VARIABLE>

```

::= IDENT IDENT @ IDENT <VARIABLE>
IDENT [ <EXPRESSION LIST> ]
IDENT [ <EXPRESSION LIST> ]. <VARIABLE>
IDENT @. <VARIABLE>

```

*)

INITSUCC (205,210)

```

WITH STABLE [204] DO
WITH STABLE [205] DO
BEGIN SUC ::= 209; TPTYPE ::= ORD (ARROW ) ;
ALT ::= 206

END ;
WITH STABLE [206] DO
BEGIN ALT ::= 209; TPTYPE ::= ORD (LBRACKET ) END ;
WITH STABLE [207] DO
WITH STABLE [208] DO
WITH STABLE [209] DO
BEGIN ALT ::= OK; TPTYPE ::= ORD (PERIOD ) END ;
WITH STABLE [210] DO

```

```

BEGIN SUC := OK; TPTYPE := 30 END ;
(*)
30) <VAR TAIL> ::= [ <EXPRESSION LIST> ]
                 | <VARIABLE>
                 | <EXPRESSION LIST> . <VARIABLE>
                 | <EMPTY>
*)

```

```

INITSUCC (213,216)
WITH STABLE [211] DO
  BEGIN ALT :=212; TPTYPE := ORD (ARROW )
  SUC :=215
  END ;
WITH STABLE [212] DO
  BEGIN ALT :=215; TPTYPE := ORD (LBRACKET ) END ;
WITH STABLE [213] DO TPTYPE := 33
WITH STABLE [214] DO TPTYPE := ORD (RBRACKET )
WITH STABLE [215] DO
  BEGIN ALT := OK; TPTYPE := ORD (PERIOD ) END ;
WITH STABLE [216] DO
  BEGIN SUC := OK; TPTYPE := 30 END ;

```

```

(*)
31) <RECORD VAR LIST> ::= <VARIABLE> /, <VARIABLE>/0/
*)

```

```

INITSUCC (218,219)
WITH STABLE [217] DO TPTYPE := 30
WITH STABLE [218] DO
  BEGIN ALT := OK; TPTYPE := ORD (COMMA ) END ;
WITH STABLE [219] DO
  BEGIN SUC :=218; TPTYPE := 30 END ;

```

```

(*)
32) <EXPRESSION LIST> ::= <EXPRESSION> /, <EXPRESSION>/0/
*)

```

```

INITSUCC (221,222)
WITH STABLE [220] DO TPTYPE := 34
WITH STABLE [221] DO
  BEGIN ALT := OK; TPTYPE := ORD (COMMA ) END ;
WITH STABLE [222] DO
  BEGIN SUC :=221; TPTYPE := 34 END ;

```

```

(*)
33) <EXPRESSION> ::= <SIMPLE EXPRESSION> RELOP
                 | <SIMPLE EXPRESSION>
*)

```

```

INITSUCC (224,225)
WITH STABLE [223] DO TPTYPE := 35
WITH STABLE [224] DO
  BEGIN ALT := OK; TPTYPE := ORD (RELOP ) END ;
WITH STABLE [225] DO
  BEGIN SUC :=224; TPTYPE := 35 END ;

```

```

(*)
34) <SIMPLE EXPRESSION> ::= ADDOP <TERM> /ADDOP <TERM>/0/
                       | <TERM> /ADDOP <TERM> /0/
*)

```

```

INITSUCC (227,229)
WITH STABLE [226] DO
  BEGIN ALT :=227; TPTYPE := ORD (ADDOP ) END ;
WITH STABLE [227] DO TPTYPE := 35
WITH STABLE [228] DO
  BEGIN ALT := OK; TPTYPE := ORD (ADDOP ) END ;
WITH STABLE [229] DO
  BEGIN SUC :=228; TPTYPE := 35 END ;

```

```

(*)
35) <TERM> ::= <FACTOR> / MULOP <FACTOR>/0/
*)

```



```

INITSUCC (231,232)
WITH STABLE [230] DO      TPTYPE := 37
WITH STABLE [231] DO
  BEGIN ALT := OK; TPTYPE := ORD (MULOP) ) END ;
WITH STABLE [232] DO
  BEGIN SUC :=231; TPTYPE := 37
  END ;

```

```

(*) 36) <FACTOR> ::= INTCONST REALCONST STRINGCONST
                CHARCONST IDENT (<EXPRESSION LIST>)
                [ <SUBSET LIST> NOT <FACTOR>
                (<EXPRESSION> ) IDENT <VAR TAIL>
*)

```

```

INITIAL := INTCONST
FOR I := 233 TO 242 DO
  WITH STABLE [ I ] DO
    BEGIN
      SUC := OK
      IF I < 237 THEN
        BEGIN
          ALT := I + 1
          TPTYPE := ORD (INITIAL)
          INITIAL := SUCC (INITIAL)
        END
      END
    END
  ;

```

```

WITH STABLE [237] DO
  BEGIN SUC :=238; TPTYPE := ORD (NOTSY)
  ALT :=239
  END ;

```

```

WITH STABLE [238] DO TPTYPE := 37

```

```

INITSUCC (240,241)

```

```

WITH STABLE [239] DO
  BEGIN ALT :=242; TPTYPE := ORD (LPARENT) ) END ;
WITH STABLE [240] DO TPTYPE := 34
WITH STABLE [241] DO TPTYPE := ORD (RPARENT) )

```

```

WITH STABLE [242] DO TPTYPE := 92

```

```

(*) 37) <SUBSET LIST> ::= ] <ELEMENT LIST> ]
*)

```

```

WITH STABLE [247] DO
  BEGIN SUC := OK; TPTYPE := ORD (RBRACKET)
  ALT :=248
  END ;

```

```

WITH STABLE [248] DO
  BEGIN SUC :=249; TPTYPE := 89
  END ;

```

```

WITH STABLE [249] DO
  BEGIN SUC := OK; TPTYPE := ORD (RBRACKET) ) END ;

```

```

(*) 38) <ELEMENT LIST> ::= <ELEMENT> /, <ELEMENT/O/
*)

```

```

INITSUCC (251,252)
WITH STABLE [250] DO TPTYPE := 90
WITH STABLE [251] DO
  BEGIN ALT := OK; TPTYPE := ORD (COMMA) ) END ;
WITH STABLE [252] DO
  BEGIN SUC :=251; TPTYPE := 90
  END ;

```

```

(*) 39) <ELEMENT> ::= <EXPRESSION>
                <EXPRESSION> *, <EXPRESSION>
*)

```

```

INITSUCC (254,255)
WITH STABLE [253] DO TPTYPE := 84
WITH STABLE [254] DO
  BEGIN ALT := OK; TPTYPE := ORD (DOTDOT) ) END ;
WITH STABLE [255] DO
  BEGIN SUC := OK; TPTYPE := 84
  END ;

```

```

(*) 40) <CASE LIST> ::= <CONSTANT LIST> : <STATEMENT>

```

*)

```

INITSUCC (257,258)
WITH STABLE [256] DO TPTYPE := 59
WITH STABLE [257] DO TPTYPE := ORD (COLON)
WITH STABLE [258] DO
  BEGIN SUC := OK; TPTYPE := 77
END ;

```

(*

41) <EMPTY> ::=

*)

```

INITSUCC (260,264)
WITH STABLE [259] DO
  BEGIN ALT := 261; TPTYPE := ORD (LBRACKET)
END ;
WITH STABLE [260] DO
  BEGIN SUC := OK; TPTYPE := 98
END ;
WITH STABLE [261] DO TPTYPE := ORD (IDENT)
WITH STABLE [262] DO
  BEGIN ALT := 265; TPTYPE := ORD (LPARENT)
END ;
WITH STABLE [263] DO TPTYPE := 33
WITH STABLE [264] DO
  BEGIN SUC := OK; TPTYPE := ORD (RPARENT)
END ;
WITH STABLE [265] DO
  BEGIN SUC := OK; TPTYPE := 81
END ;
FOR INITIAL := IDENT TO WITHSY DO
  BEGIN
    MISSING [ INITIAL ] := 0 ;
    TTABLE [ ORD ( INITIAL ) ] CLASS := INITIAL
  END ;

```

INITSUCC (267,268)

```

WITH STABLE [266] DO TPTYPE := 91
WITH STABLE [267] DO
  BEGIN ALT := OK; TPTYPE := ORD (SEMICOLON)
END ;
WITH STABLE [268] DO
  BEGIN SUC := 267; TPTYPE := 91
END ;

```

MISSING [IDENT]	J	::=	2;	MISSING [INTCONST]	J	::=	15;
MISSING [LPARENT]	J	::=	9;	MISSING [RPARENT]	J	::=	14;
MISSING [LBRACKET]	J	::=	11;	MISSING [RBRACKET]	J	::=	12;
MISSING [COMMA]	J	::=	20;	MISSING [SEMICOLON]	J	::=	14;
MISSING [COLON]	J	::=	5;	MISSING [ASSIGNOP]	J	::=	51;
MISSING [BEGINSY]	J	::=	17;	MISSING [EOSY]	J	::=	54;
MISSING [DOWNIOSY]	J	::=	55;	MISSING [ENDSY]	J	::=	53;
MISSING [FILES]	J	::=	57;	MISSING [IFSY]	J	::=	56;
MISSING [TOSY]	J	::=	55;	MISSING [PROGRAMSY]	J	::=	53;
MISSING [ETHENSY]	J	::=	52;	MISSING [UNTILSY]	J	::=	53;
MISSING [OFSY]	J	::=	8;	MISSING [RELOP]	J	::=	16;
MISSING [EMULOP]	J	::=	21;	MISSING [ADDOP]	J	::=	22;
MISSING [REPEATSY]	J	::=	23;	MISSING [LABELSY]	J	::=	24;
MISSING [CONSTSY]	J	::=	25;	MISSING [VARSY]	J	::=	26;
MISSING [TYPESY]	J	::=	27;	MISSING [DOTDOT]	J	::=	28;

```

FOR I := 52 TO 93 DO
  BEGIN
    TTABLE [ I ] TERMINAL := FALSE ;
    NONTRMSGSS [ I ] := 0
  END ;

```

```

NONTRMSGSS [63] := 1; NONTRMSGSS [61] := 10 ;
NONTRMSGSS [70] := 18; NONTRMSGSS [64] := 19 ;

```

```

NONTRMSGSS [56] := 18; NONTRMSGSS [57] := 18 ;
NONTRMSGSS [58] := 18; NONTRMSGSS [60] := 18 ;
NONTRMSGSS [59] := 50; NONTRMSGSS [87] := 58 ;
NONTRMSGSS [80] := 59;

```

```

FOR I := 0 TO ORD (WITHSY) DO
  TTABLE [ I ] TERMINAL := TRUE ;

```

TTABLE [52] POINTER	::=	1;	TTABLE [53] POINTER	::=	4;
TTABLE [54] POINTER	::=	10;	TTABLE [55] POINTER	::=	13;
TTABLE [56] POINTER	::=	16;	TTABLE [57] POINTER	::=	20;
TTABLE [58] POINTER	::=	25;	TTABLE [59] POINTER	::=	33;
TTABLE [60] POINTER	::=	42;	TTABLE [61] POINTER	::=	50;
TTABLE [64] POINTER	::=	83;	TTABLE [65] POINTER	::=	92;
TTABLE [62] POINTER	::=	73;	TTABLE [63] POINTER	::=	76;
TTABLE [66] POINTER	::=	97;	TTABLE [67] POINTER	::=	100;
TTABLE [68] POINTER	::=	104;	TTABLE [69] POINTER	::=	107;
TTABLE [70] POINTER	::=	110;	TTABLE [71] POINTER	::=	122;
TTABLE [72] POINTER	::=	134;	TTABLE [73] POINTER	::=	137;

```

TTABLE [74] : POINTER := 150 ; TTABLE [75] : POINTER := 150 ;
TTABLE [76] : POINTER := 159 ; TTABLE [77] : POINTER := 159 ;
TTABLE [78] : POINTER := 204 ; TTABLE [79] : POINTER := 204 ;
TTABLE [80] : POINTER := 217 ; TTABLE [81] : POINTER := 217 ;
TTABLE [82] : POINTER := 223 ; TTABLE [83] : POINTER := 223 ;
TTABLE [84] : POINTER := 230 ; TTABLE [85] : POINTER := 230 ;
TTABLE [86] : POINTER := 237 ; TTABLE [87] : POINTER := 237 ;
TTABLE [88] : POINTER := 247 ; TTABLE [89] : POINTER := 247 ;
TTABLE [90] : POINTER := 253 ; TTABLE [91] : POINTER := 253 ;
TTABLE [92] : POINTER := 259 ; TTABLE [93] : POINTER := 259 ;

```

```

OKCOUNT := 0 ; RECOVERED := FALSE ;
RECIND := 0 ; LAST := NIL ;
RIND := 0 ; ERROUPL := EMPTYSET ; ALTFLAG := FALSE ;

```

```

FOR TOP := 0 TO DISPLIMIT DO WITH DISPLAY [TOP] DO BEGIN

```

```

    FNAME := NIL ;
    OCCUR := BLCK ;

```

```

END ;

```

```

IDLIST := FILEHEADS ; OLDS := 0 ; PRNTABLE := FALSE ;
TOP := 0 ; FWPTR := NIL ; PRTErr := TRUE ;
SIGN := NOSIGN ; SAVIO := NIL ; PCKD := FALSE ;
ARRPIND := 0 ; FSTLARPTR := NIL ; ALLOWDOTS := FALSE ;
LEVEL := 0 ; RECINX := 0 ; KOLON := FALSE ;
LKIND := ACTUAL ; RCASEIND := 0 ; EMPTY := FALSE ;

```

```

END ; (* INITPARSE *)

```

```

PROCEDURE ENTERID (LCP : IDP) ; FORWARD ;
FUNCTION ISTRING (FSP : STP) : BOOLEAN ; FORWARD ;
PROCEDURE GETBOUNDS (FSP : STP ; VAR FMIN, FMAX : INTEGER) ; FORWARD ;
PROCEDURE ERROR (ERRNO : INTEGER) ; FORWARD ;

```

PROCEDURE ENTERSTDIDS ;

(*

.. THIS ROUTINE INITIALIZES STANDARD NAMES AND THEIR
 ... RELATED STRUCTURE.

*)

```

BEGIN
  NEW (INTPTR, SCALAR, STANDARD) ; (* S *)
  WITH INTPTR @ DO (* T *)
    BEGIN SIZE := 1 ; MARKED := FALSE END ; (* P *)

  NEW (INTID, TYPES) ; (* I *)
  WITH INTID @ DO BEGIN (* N *)
    NAME := 'INTEGER' ; (* T *)
    LLINK := NIL ; RLINK := NIL ; (* E *)
    IDTYPE := INTPTR ; NEXT := NIL ; (* G *)
    UNDCL := 0 ; (* F *)
  END ; ENTERID (INTID) ; (* R *)

  NEW (MAXINTID, KONST) ;
  WITH MAXINTID @ DO BEGIN (* M *)
    NAME := 'MAXINT' ; (* A *)
    LLINK := NIL ; RLINK := NIL ; (* X *)
    IDTYPE := INTPTR ; NEXT := NIL ; (* I *)
    UNDCL := 0 ;
    WITH VALUES DO BEGIN (* T *)
      INTVAL := TRUE ; CLASS := INTGR ;
      IVAL := MAXINT
    END
  END ; ENTERID (MAXINTID) ;

  NEW (REALPTR, SCALAR, STANDARD) ; (* S *)
  WITH REALPTR @ DO (* T *)
    BEGIN SIZE := 2 ; MARKED := FALSE END ; (* P *)

  NEW (REALID, TYPES) ; (* R *)
  WITH REALID @ DO BEGIN (* E *)
    NAME := 'REAL' ; (* A *)
    LLINK := NIL ; RLINK := NIL ; (* L *)
    IDTYPE := REALPTR ; NEXT := NIL ;
    UNDCL := 0 ;
  END ; ENTERID (REALID) ;

  NEW (CHARPTR, SCALAR, STANDARD) ;
  WITH CHARPTR @ DO
    BEGIN SIZE := 1 ; MARKED := FALSE END ;

  NEW (CHARID, TYPES) ; (* C *)
  WITH CHARID @ DO BEGIN (* H *)
    NAME := 'CHAR' ; (* A *)
    LLINK := NIL ; RLINK := NIL ; (* R *)
    IDTYPE := CHARPTR ; NEXT := NIL ;
    UNDCL := 0 ;
  END ; ENTERID (CHARID) ;

  NEW (TRUEID, KONST) ; (* B *)
  NEW (FALSEID, KONST) ; (* O *)
  NEW (BOOLPTR, SCALAR, DECLARED) ; (* O *)
  (* L *)
  WITH BOOLPTR @ DO BEGIN (* E *)
    SIZE := 1 ; MARKED := FALSE ; (* A *)
    FSTCONST := TRUEID (* N *)
  END ;

  NEW (BOOLID, TYPES) ;
  WITH BOOLID @ DO BEGIN
    NAME := 'BOOLEAN' ;
    LLINK := NIL ; RLINK := NIL ; NEXT := NIL ;
    UNDCL := 0 ; IDTYPE := BOOLPTR
  END ; ENTERID (BOOLID) ;

  WITH TRUEID @ DO BEGIN (* T *)
    NAME := 'TRUE' ; (* R *)
    LLINK := NIL ; RLINK := NIL ; (* U *)
    IDTYPE := BOOLPTR ; NEXT := NIL ; (* E *)
    UNDCL := 0 ;
    WITH VALUES DO BEGIN
      INTVAL := TRUE ; CLASS := INTGR ;
      IVAL := 1
    END
  END ; ENTERID (TRUEID) ;

```

```

NAME := 'FALSE'; DO BEGIN
  LLINK := NIL ; RLINK := NIL ;          (* F *)
  IDTYPE := BOOLPTR ; NEXT := NIL ;      (* A *)
  UNDCL := 0 ;                            (* L *)
  WITH VALUES DO BEGIN                  (* S *)
    INTVAL := TRUE ; CLASS := INTGR ;    (* F *)
    IVAL := 0
  END
END ; ENTERID (FALSEID) ;

NEW (TEXTPTR, FILES) ;
WITH TEXTPTR @ DO BEGIN                (* S *)
  SIZE := MAXINT ; MARKED := FALSE ;    (* T *)
  FILTYPE := CHARPTR                    (* P *)
END ;

NEW (TEXTID, TYPES) ;
WITH TEXTID @ DO BEGIN                  (* T *)
  NAME := 'TEXT' ;                       (* E *)
  LLINK := NIL ; RLINK := NIL ;          (* X *)
  NEXT := NIL ; IDTYPE := TEXTPTR ;     (* T *)
  UNDCL := 0
END ; ENTERID (TEXTID) ;

NEW (NILPTR, POINTER) ;
WITH NILPTR @ DO BEGIN                  (* S *)
  SIZE := 1 ; MARKED := FALSE ;          (* T *)
  ELTYPE := NIL                          (* P *)
END ;

NEW (NILID, KONST) ;
WITH NILID @ DO BEGIN                   (* N *)
  NAME := 'NIL' ;                         (* I *)
  LLINK := NIL ; RLINK := NIL ;          (* L *)
  NEXT := NIL ; UNDCL := 0 ;
  WITH VALUES DO BEGIN
    INTVAL := TRUE ; CLASS := INTGR ;
    IVAL := 'NILVAL' ; *****
  END
END ; ENTERID (NILID) ;

NEW (UTYPPTR) ;
WITH UTYPTR @ DO BEGIN
  NAME := ' ' ;
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;
  NEXT := NIL ; UNDCL := 0 ; KLASS := TYPES
END ;

NEW (UVAPPTR) ;
WITH UVAPPTR @ DO BEGIN
  NAME := ' ' ;
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;
  NEXT := NIL ; UNDCL := 0 ; KLASS := VARS
END ;

NEW (UFLDPTR) ;
WITH UFLDPTR @ DO BEGIN
  NAME := ' ' ;
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;
  NEXT := NIL ; UNDCL := 0 ; KLASS := FIELD
END ;

NEW (UKONSPTR) ;
WITH UKONSPTR @ DO BEGIN
  NAME := ' ' ;
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;
  NEXT := NIL ; UNDCL := 0 ; KLASS := TYPES ;
  WITH VALUES DO BEGIN
    INTVAL := TRUE ; CLASS := INTGR ;
    IVAL := 0
  END
END ;

NEW (UPRCPTR) ;
WITH UPRCPTR @ DO BEGIN
  NAME := ' ' ;
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;
  NEXT := NIL ; UNDCL := 0 ; KLASS := PROC ;
  PFDECLKIND := DECLARED ; PFLEV := 0 ;
  PARAMPTR := NIL ; PFKIND := ACTUAL ;
  FORWDECL := FALSE
END ;

```

```
NEW (UFCTPTR) ;  
WITH UFCTPTR @ DO BEGIN  
  NAME := ' ;  
  LLINK := NIL ; RLINK := NIL ; IDTYPE := NIL ;  
  NEXT := NIL ; UNDECL := 0 ; CLASS := FUNC ;  
  PFDECLKIND := DECLARED ; FLEV := 0 ;  
  PARAMPTR := NIL ; PFKIND := ACTUAL ;  
  FORWDECL := FALSE  
END
```

```
END ; (* ENTERSTDIDS *)
```

PROCEDURE PRINTABLES (ALL : BOOLEAN) ;

VAR I,LIM : DISPRANGE ;
PROCEDURE MARKS ;
MARKS SYMBOL TABLE ENTRIES TO PREVENT
MULTIPLE PRINTOUT

VAR I : INTEGER ;
PROCEDURE MARKCTP (FP : IDP) : FORWARD ;
PROCEDURE MARKSTP (FP : STP) ;
MARK DATA STRUCTURES TO PREVENT CYCLES
VAR I : INTEGER ;

```
BEGIN
  IF FP <> NIL THEN
    WITH FP @ DO
      IF NOT MARKED THEN BEGIN
        MARKED := TRUE ;
        CASE FORM OF
          SCALAR           : MARKSTP (PANGETYPE) ;
          SUBRANGE         : MARKSTP (ELTYPE) ;
          POINTER          : MARKSTP (ELSET) ;
          POWER            : MARKSTP (ELTYPE) ;
          ARRAYS           : MARKSTP (AELTYPE) ;
          BEGIN           : MARKSTP (INXTYPE) ;
          RECORDS          : MARKCTP (FSTFLD) ;
          BEGIN           : MARKSTP (RECVAR) ;
          PARAMLIST        : MARKCTP (FSTPAR) ;
          FILES            : MARKCTP (FILTYPE) ;
          TAGFIELD         : MARKSTP (TAGTYPE) ;
          BEGIN           : MARKSTP (FSTVAR) ;
          VARIANT          : MARKCTP (NXTVAR) ;
          BEGIN           : MARKSTP (SUBREC) ;
        END (* CASE *)
      END (* IF *)
    END ; (* MARKSTP *)
```

```
PROCEDURE MARKCTP (FP : IDP) ;
BEGIN
  IF FP <> NIL THEN
    WITH FP @ DO BEGIN
      MARKCTP (LLINK) ; MARKSTP (IDTYPE) ;
      MARKCTP (RLINK) ;
      IF (KLASS IN [PROC, FUNC]) AND (PFDECLKIND=DECLARED)
      THEN MARKSTP (PARAMPTR)
    END
  END ; (* MARKCTP *)
```

```
BEGIN (* MARKS *)
  FOR I := TOP DOWNTO LIM DO MARKCTP (DISPLAY [I].FNAME)
END ; (* MARKS *)
```

```
PROCEDURE FOLLOWCTP (FP : IDP) : FORWARD ;
PROCEDURE FOLLOWSTP (FP : STP) ;
BEGIN
  IF FP <> NIL THEN
    WITH FP @ DO
      IF MARKED THEN BEGIN
        WRITELN ; WRITELN ;
        WRITE (' ' : 20, '****STRUCTURE****') ;
        WRITELN ;
        MARKED := FALSE ; WRITELN ;
        WRITE (' ' : 7, 'INTERNAL STRUCTURE CODE=', FP,
              ' SIZE=', SIZE) ;
        WRITELN ; WRITE (' ' : 7, 'FORM=') ;
        CASE FORM OF
          SCALAR :
            BEGIN WRITE (' SCALAR ', ' ') ;
                  IF SCALKIND = STANDARD THEN
                    WRITE (' STANDARD ') ;
        END
```

```

ELSE WRITE ('DECLARED
FIRST CONSTANT ID PTR=', #STCONST ) ;
WRITELN
END ;
SUBRANGE :
BEGIN WRITE ('SUBRANGE ', 'RANGE PTR=',
RANGETYPE, 'MIN=', MIN, 'MAX=', MAX) ;
FOLLOWSTP (RANGETYPE)
END ;
POINTER :
BEGIN
WRITE ('POINTER ', 'ELTYPE PTR=',
ELTYPE ) ;
FOLLOWSTP (ELTYPE)
END ;
POWER :
BEGIN WRITE ('SET
'ELSET PTR =', ELSET) ;
FOLLOWSTP (ELSET)
END ;
ARRAYS :
BEGIN
IF PACKD THEN WRITE ('D-ARRAY : ', ' )
ELSE WRITE ('ARRAY
WRITE ('ELEMENT TYPE PTR=', AELTYPE,
'INDEX TYPE PTR=', INXTYPE) ;
WRITELN ;
FOLLOWSTP (AELTYPE) ;
FOLLOWSTP (INXTYPE)
END ;
RECORDS :
BEGIN
WRITE ('RECORD ', 'FIRST FIELD PTR=',
FSTFLD, 'VAR PTR=', RECVAR ) ;
FOLLOWCTP (FSTFLD) ;
FOLLOWSTP (RECVAR)
END ;
PARAMLIST :
BEGIN
WRITE ('PARAM LIST, ',
'FIRST PARAMETER PTR=', FSTPAR) ;
FOLLOWCTP (FSTPAR)
END ;
FILES :
BEGIN WRITE ('FILE
'FILE TYPE PTR=', FILTYPE ) ;
FOLLOWSTP (FILTYPE)
END ;
TAGFIELD :
BEGIN WRITE ('TAGFIELD
'TAGFIELD PTR=', TAGFIELDP,
'TAGTYPE PTR=', TAGTYPE,
'FIRST VAR PTR =', FSTVAR ) ;
FOLLOWSTP (TAGTYPE) ; FOLLOWSTP (FSTVAR)
END ;
VARIANT :
BEGIN
WRITE ('VARIANT
'NEXT VARIANT PTR=', NXTVAR,
'SUBREC: =', SUBREC,
'VARIANT VALUE =', VARVAL) ;
FOLLOWSTP (NXTVAR) ;
FOLLOWSTP (SUBREC)
END

```



```

        OTHERWISE
        END (* CASE *)
    END (* IF *)
END ; (* FOLLOWSTP *)

PROCEDURE FOLLOWCTP (FP : IDR) ;
BEGIN
    IF FP <> NIL THEN
        WITH FP & DO BEGIN
            FOLLOWCTP (LLINK) ;
            WRITELN : WRITELN ;
            WRITE (' ' : 20, '****IDENTIFIER****') ;
            WRITELN :
            WRITE (' ' : 7, ' INTERNAL CODE=', FP,
                NAME = ', NAME ) ;
            WRITELN :
            WRITE (' ' : 7, 'LLINK=', LLINK, ' RLINK=', RLINK,
                ' NEXT=', NEXT, ' IDPTR=', IDTYPE ) ;
            WRITELN :
            WRITE (' ' : 7, 'CLASS=') ;
            CASE CLASS OF
                TYPES      : WRITE ('TYPE      ') ;
                KONST      :
            BEGIN
                WRITE ('CONSTANT ') ;
                IF IDTYPE <> NIL THEN
                    IF ISTRING (IDTYPE) THEN
                        WRITE ('STRING ', 'CONST VALUE=',
                            VALUES.SVAL ) ELSE
                    IF IDTYPE = REALPTR THEN
                        WRITE ('REAL ', 'CONST VALUE=',
                            VALUES.RVAL ) ELSE
                    IF VALUES.INTVAL THEN
                        IF IDTYPE = INTPTR THEN
                            WRITE ('INTEGER ', 'CONST VALUE=',
                                VALUES.IVAL ) ELSE
                        IF IDTYPE = BOOLPTR THEN
                            WRITE ('BOOLEAN ', 'CONST VALUE=',
                                VALUES.IVAL = 1 ) ELSE
                        IF IDTYPE = CHARPTR THEN
                            WRITE ('CHARACTER ', 'CONST VALUE=',
                                CHR (VALUES.IVAL) )
                        ELSE
                            WRITE (' ', 'CONST VALUE=', VALUES.IVAL) ;
                    WRITELN
                END ;
            VARS :
            BEGIN
                WRITE ('VARIABLE ', ' ) ;
                IF VKIND = ACTUAL THEN
                    WRITE ('ACTUAL ', ' ) ;
                ELSE WRITE ('FORMAL ', ' ) ;
                WRITE ('LEVEL=', VLEV : 4) ;
                IF VLEV <> I THEN FOLLOWCTP (NEXT)
            END ;
            FIELD : WRITELN ('FIELD      ') ;
            PARAMS : WRITELN ('PARAM LIST') ;
            PROC, FUNC :
            BEGIN
                IF KCLASS = PROC THEN WRITE ('PROCEDURE ', ' )
                ELSE WRITE ('FUNCTION ', ' ) ;
                IF PFDECLKIND = STANDARD THEN
                    WRITE ('STANDARD ', ' )
                ELSE BEGIN
                    IF PFKIND = ACTUAL THEN WRITE ('ACTUAL ', ' )
                    ELSE WRITE ('FORMAL ', ' ) ;
                    WRITE ('LEVEL=', PFLEV) ;
                    IF FORWDECL THEN WRITE (' ', 'FORWARD') ;
                    WRITE (' ', 'PARAMPTR=', 'PARAMPTR') ;
                END
            END
        END
    END
END ;

```

```
END FOLLOWSTP (PARAMPTR)
END ;
OTHERWISE
END ;
FOLLOWSTP (IDTYPE) ;
FOLLOWCTP (RLINK) ;
END (* WITH *)
END ; (* FOLLOWCTP *)
```

```
BEGIN (* PRINTABLES *)
IF PRINTABLE THEN BEGIN
WRITELN (' *****');
WRITELN (' * SYMBOL TABLE DUMP *');
WRITELN (' *****');
WRITELN ;

IF ALL THEN WRITELN (' .. FOR PROGRAM.. ');
ELSE WRITELN (' .. INTERMEDIATE DUMP.. ');

IF ALL THEN LIM := 0
ELSE
LIM := TOP ;

MARKS ;
WRITELN ;
FOR I := TOP DOWNT0 LIM DO BEGIN
FOLLOWCTP (DISPLAY [I].FNAME) ;
WRITELN
END ; WRITELN ; WRITELN ; WRITELN END
END ; (* PRINTABLES *)
```

```

FUNCTION COMPTYPES (FSP1, FSP2 : STP) : BOOLEAN ;
VAR STRICT : BOOLEAN ;

```

```

(*)
** COMPARE RETURNS TRUE IF FSP1 IS COMPATIBLE
** WITH FSP2
*)

```

```

LABEL 1 ;
VAR NXT1, NXT2 : IDP ; COMP : BOOLEAN ;
    MINI, MIN2, MAX1, MAX2 : INTEGER ;

```

```

BEGIN
IF (FSP1=FSP2) OR (FSP1=NIL) OR (FSP2=NIL) THEN
COMP := TRUE

```

```

ELSE
IF FSP1 @. FORM = FSP2 @. FORM THEN
CASE FSP1 @. FORM OF
SCALAR, (* SCALARS DECLARED ON DIFFERENT
LEVELS ARE INCOMPATIBLE *)
RECORDS, TAGFIELD, VARIANT, FILES :
COMP := FALSE ;
SUBRANGE :
COMP := NOT STRICT AND
COMPARE (FSP1 @. RANGETYPE, FSP2 @. RANGETYPE) ;
POINTER :
COMP := (FSP1 @. ELTYPE=NIL) AND
(FSP2 @. ELTYPE=NIL) ;
POWER :
COMP := NOT STRICT AND (FSP1 @. SIZE = FSP2 @. SIZE)
AND (FSP1 @. ELTYPE = FSP2 @. ELTYPE) ;
ARRAYS :
BEGIN
IF NOT STRICT AND
ISTRING (FSP1) AND ISTRING (FSP2) THEN
BEGIN
GETBOUNDS (FSP1 @. INXTYPE, MINI1, MAX1) ;
GETBOUNDS (FSP2 @. INXTYPE, MINI2, MAX2) ;
COMP := (MIN1 = 1) AND (MIN2 = 1)
AND (MAX1 = MAX2) AND
(FSP1 @. PACKD = FSP2 @. PACKD) AND
COMPARE (FSP1 @. INXTYPE, INTPTR) AND
COMPARE (FSP2 @. INXTYPE, INTPTR)
END
ELSE COMP := FALSE
END ;

```

```

PARAMLIST :
BEGIN
COMP := TRUE ; STRICT := TRUE ;
NXT1 := FSP1 @. FSTPAR ;
NXT2 := FSP2 @. FSTPAR ;
WHILE COMP AND (NXT1 <> NIL) AND (NXT2 <> NIL)
DO
BEGIN
COMP := NXT1 @. KCLASS = NXT2 @. KCLASS ;
IF NXT1 @. KCLASS = VARS THEN
COMP := COMP AND
(NXT1 @. VKIND = NXT2 @. VKIND)
ELSE
COMP := COMP AND
((NXT1 @. PARAMPTR @. FSTPAR =
NXT2 @. PARAMPTR @. FSTPAR) AND
NXT1 @. PARAMPTR @. ANON AND
NXT2 @. PARAMPTR @. ANON AND
COMPARE (NXT1 @. PARAMPTR,
NXT2 @. PARAMPTR)) ;
COMP := COMP AND
COMPARE (NXT1 @. IDTYPE ;
NXT2 @. IDTYPE) ;
NXT1 := NXT1 @. NEXT ;
NXT2 := NXT2 @. NEXT ;
END

```

```

END ;
OTHERWISE
END
ELSE (* FSP1 @. FORM <> FSP2 @. FORM *)
IF FSP1 @. FORM = SUBRANGE THEN
COMP := NOT STRICT AND (COMPARE (FSP1 @. RANGETYPE,
FSP2))
ELSE
IF FSP2 @. FORM = SUBRANGE THEN

```

```
      COMP := NOT STRICT AND (COMPARE (FSP2),RANGETYPE,  
      ELSE  
      COMP := FALSE ;  
1 : COMPARE := COMP ;  
  END ; (* COMPARE *)
```

```
BEGIN  
  STRICT := FALSE ;  
  COMPTYPES := COMPARE (FSP1,FSP2 )  
END ;
```

```
PROCEDURE ENTERID (IDPTR : IDP) ;
```

```
.. ENTERS ID POINTED BY IDPTR INTO THE SYMBOL TABLE WHICH  
ON EACH DECLARATION LEVEL IS ORGANIZED AS AN UNBALANCED  
TREE
```

```
VAR   NAM      : PACKEDID ;  
      LCP, LCP1 : IDP     ;  
      LLEFT    : BOOLEAN  ;
```

```
BEGIN  
  NAM := IDPTR @. NAME ;  
  LCP := DISPLAY [ TOP ], FNAME ;  
  
  IF LCP = NIL  
  THEN DISPLAY [ TOP ], FNAME := IDPTR  
  ELSE BEGIN  
    REPEAT  
      LCP1 := LCP ;  
      IF LCP @. NAME = NAM THEN BEGIN  
        IF PRERR THEN ERROR (101) ;  
        LCP := LCP @. RLINK ;  
        LLEFT := FALSE ;  
      END  
      ELSE IF LCP @. NAME < NAM THEN  
        BEGIN  
          LCP := LCP @. RLINK ;  
          LLEFT := FALSE ;  
        END  
      ELSE BEGIN  
        LCP := LCP @. LLINK ;  
        LLEFT := TRUE ;  
      END (* IF *)  
    UNTIL LCP = NIL ;  
  
    IF LLEFT THEN LCP1 @. LLINK := IDPTR  
    ELSE          LCP1 @. RLINK := IDPTR  
  
  END ; (* IF THEN ELSE *)
```

```
WITH IDPTR @ DO BEGIN  
  LLINK := NIL ;  
  RLINK := NIL ;  
  UNDC1 := 0 ;  
END
```

```
END ; (* ENTERID *)
```

```
PROCEDURE DECLAPE (IDPTR : IDP) ;
```

```
VAR LTOP : DISPRANGE ;  
BEGIN  
  LTOP := TOP ;  
  WHILE DISPLAY [ LTOP ], OCCUR <> BLCK DO TOP := TOP + 1 ;  
  PRERR := FALSE ;  
  ENTERID (IDPTR) ;  
  PRERR := TRUE ;  
  IDPTR @. UNDC1 := CARDCNT ;  
  TOP := LTOP ;  
END ;
```

```
PROCEDURE SEARCHID (FIDCLASS : SETOFIDS; VAR FCP : IDP) ;
```

```
  LABEL 1,3  
  VAR LCP : IDP ;
```

```
BEGIN IF RECOVERED THEN GOTO 3 ;  
  FOR DISX := TOP DOWNT0 0 DO BEGIN  
    LCP := DISPLAY [ DISX ], FNAME ;  
    WHILE LCP <> NIL DO  
      IF LCP @. NAME = IDENTIFIER THEN BEGIN  
        IF LCP @. KLAS IN FIDCLASS THEN GOTO 1 ;  
        IF PRERR THEN ERROR (103) ;  
        LCP := LCP @. RLINK ;  
      END  
      ELSE IF LCP @. NAME < IDENTIFIER THEN  
        LCP := LCP @. RLINK  
      ELSE LCP := LCP @. LLINK
```

```
END ; (* FOR *)
```

```
IF PRIERR THEN BEGIN  
  ERROR_(104) ;  
  GUNIT.E104 := TRUE ;
```

```
3 : NEW (LCP) ;  
  IF TYPES IN FIDCLASS THEN LCP@ := UTYPEPTR @ ELSE  
  IF KONST IN FIDCLASS THEN LCP@ := UKONSPTR @ ELSE  
  IF VARS IN FIDCLASS THEN LCP@ := UVARPTR @ ELSE  
  IF FIELD IN FIDCLASS THEN LCP@ := UFLOPTR @ ELSE  
  IF PROC IN FIDCLASS THEN LCP@ := UPROPTR @ ELSE  
  ELSE LCP@ := UFCTPTR @ ;
```

```
  LCP @. NAME := IDENTIFIER ;  
  DECLARE (LCP)  
END ; (* IF *)
```

```
1 :
```

```
  FCP := LCP  
END ; (* SEARCHID *)
```

```
PROCEDURE CREATSUBRANGE (VAR FSP:STP; FMIN, FMAX : INTEGER ;  
  FSP1 : STP) ;
```

```
BEGIN  
  NEW (FSP, SUBRANGE) ;  
  WITH FSP @ DO BEGIN  
    SIZE := 1 ; MARKED := FALSE ;  
    RANGETYPE := FSP1 ;  
    MIN := FMIN ; MAX := FMAX  
  END  
END ;
```

```
FUNCTION ISTRING (FSP : STP) : BOOLEAN ;  
BEGIN  ISTRING := FALSE ;  
  IF FSP <> NIL THEN  
    IF FSP @. FORM = ARRAYS THEN  
      IF FSP @. ELTYPE = CHARPTR THEN  
        ISTRING := TRUE  
  END ;
```

```
PROCEDURE SEARCHSECTION (FCP : IOP; VAR FCP1 : IOP) ;  
  LABEL 16 ;
```

```
BEGIN  
  WHILE FCP <> NIL DO BEGIN  
    IF FCP @. NAME = IDENTIFIER THEN GOTO 16 ;  
    IF FCP @. NAME < IDENTIFIER THEN  
      FCP := FCP @. RLINK  
    ELSE  
      FCP := FCP @. LLINK  
  END ;  
  FCP1 := FCP  
END ;
```

```
16 :
```

```
FUNCTION LENGTH (FSP : STP) : INTEGER ;  
(* ASSUMES FSP @. FORM <= SUBRANGE AND  
  FSP <> REALPTR  
*)  
  VAR LMIN, LMAX : INTEGER ;  
  BEGIN  
    IF FSP = NIL THEN LENGTH := 0 ELSE BEGIN  
      GETBOUNDS (FSP, LMIN, LMAX) ;  
      LENGTH := LMAX - LMIN + 1  
    END  
  END ;
```

```
PROCEDURE GETBOUNDS (FSP:STP ; VAR FMIN, FMAX : INTEGER) ;
```

```
BEGIN  
  IF (FSP=INTPTR) OR (FSP=NIL) THEN  
    BEGIN  FMIN := MAXINT ;  
          FMAX := MAXINT  
    END  
  ELSE  
    WITH FSP@ DO BEGIN  
      IF FORM = SUBRANGE THEN  
        BEGIN  FMIN := MIN ;  
              FMAX := MAX  
        END  
      ELSE BEGIN  
        FMIN := 0 ;
```

```
IF FSP = CHARPTR THEN FMAX := 63 ELSE
IF FSP < FSTCONST THEN NIL
THEN FMAX := FSP * FSTCONST @ VALUES.IVAL
ELSE FMAX := 0
END (* IF THEN ELSE *)
END (* WITH *)
END ; (* GETBOUNDS *)
```

```

*****
*) PROCEDURE SEMANTICS (ACT : INTEGER) ;
*) THIS ROUTINE PERFORMS THE NECESSARY SEMANTIC
*) ACTION EACH TIME A TERMINAL HAS BEEN PARSED.
*) IT ALSO GIVES ASSOCIATED ERROR MESSAGES IF THERE
*) ARE ANY.
*****

```

```

LABEL 122,22 ;
BEGIN
CASE ACT OF

```

```

    (*****
    (* POINTERS *)
    (*****

```

```

#14 : PRINTABLES (LEVEL = 0) ;

```

```

50 :

```

```

BEGIN
    NEW (LSP , POINTER) ;
    WITH LSP % DO
        BEGIN SIZE := 1 ; ELTYPE := NIL END ;
    RESULT := LSP
END ;

```

```

51 : (* IDENT PART OF POINTER *)

```

```

IF NOT RECOVERED THEN BEGIN
    TTOP := TOP ;
    WHILE DISPLAY [TTOP].OCCUR <> BLCK DO TTOP := TTOP + 1 ;

```

```

    SEARCHSECTION (DISPLAY [TTOP].FNAME, LCP) ;
    IF LCP = NIL THEN

```

```

        BEGIN
            NEW (LCP , TYPES) ;
            WITH LCP % DO BEGIN
                NAME := IDENTIFIER ;
                IDTYPE := LSP ;
                NEXT := FWPTR ;
            END ;
            FWPTR := LCP ;

```

```

        ELSE IF LCP % KCLASS <> TYPES THEN
            ERROR (103) ;
        ELSE LCP % ELTYPE := LCP % IDTYPE ;

```

```

    END ;

```

```

    (*****
    (* ARRAYS *)
    (*****

```

```

52 : PCKD := TRUE ;
58 : RESULT := NIL ;

```

```

#73, #75 :

```

```

BEGIN
    IF ACT = #73 THEN LSP2 := NIL ;
    NEW (LSP, ARRAYS) ;
    WITH LSP % DO BEGIN
        AELTYPE := LSP2 ;
        INXTYPE := NIL ;
        SIZE := 1 ;
        BACKD := PCKD ;
    END ;

```

```

LSP2 := LSP ;

```

```

IF RESULT <> NIL THEN
    IF RESULT % FORM <= SUBRANGE THEN BEGIN
        IF RESULT = REALPTR THEN BEGIN
            ERROR (109) ;
            RESULT := NIL ;
        ELSE IF RESULT = INTPTR THEN BEGIN
            ERROR (149) ;
        END ;
    END ;

```



```

        RESULT := NIL                                END ;
    IF RESULT <> NIL THEN BEGIN
        GETBOUNDS (RESULT, LMIN, LMAX) ;
        IF ((LMIN < MINMIN) OR (LMIN > MAXMIN)) AND
           ((LMIN < MAXINT) OR (LMAX > MAXINT)) THEN
            BEGIN
                ERROR (173) ;
                WITH RESULT @ DO
                    BEGIN MIN := 0 ; MAX := 0 END
            END
        END ; (* IF *)
    LSP @ INXTYPE := RESULT
END (* IF *)
ELSE
    ERROR (113)

```

END ;

61 :

```

BEGIN
    ARRIND := ARRIND + 1 ;
    IF ARRIND > 20 THEN BEGIN
        WRITELN ('**TOO MUCH NESTING OF ARRAYS **') ;
        WRITELN ('**COMPILATION ABORTED **') ;
        GOTO 1111
    END ;
    ARRAYST [ARRIND] := LSP2
END ;

```

62 :

```

BEGIN
    LSP1 := ARRAYST [ARRIND] ;
    REPEAT
        WITH LSP1 @ DO BEGIN
            LSP2 := AELTYPE ;
            AELTYPE := RESULT ;
            IF INXTYPE <> NIL THEN
                SIZE := SIZE * LENGTH (INXTYPE)
            END ;
            RESULT := LSP1 ;
            LSP1 := LSP2
        UNTIL LSP1 = NIL ;
        ARRIND := ARRIND + 1
    END ;

```

```

    (*****
    (*      RECORDS      *)
    (*****

```

63 : (* RECORD *)

```

BEGIN
    OLDS := OLDS + 1 ;
    OLDTOPS [OLDS] := TOP ;
    IF TOP < DISPLIMIT THEN BEGIN
        TOP := TOP + 1 ;
        WITH DISPLAY [TOP] DO BEGIN
            FNAME := NIL ;
            OCCUR := REC ;
            END END
        ELSE ERROR (250) ;
        FRECVAR := NIL ; SAVEDTAG := NIL ;
        FCP := NIL ; LSP := NIL ; LCB := NIL ;
        LCP1 := NIL ; LCP2 := NIL
    END ;

```

END ;

83,85,89 : (** RECORD IDENTIFIERS**)

```

BEGIN
    NEW (LCP, FIELD) ;
    WITH LCP @ DO BEGIN

```

```

        NAME := IDENTIFIER ;
        IDTYPE := NIL ;
        NEXT := NIL ;
    END ;

    ENTERID (LCP) ;

    IF LCP1 = NIL THEN
        LCP1 := LCP ;
        IF LCP2 <> NIL THEN LCP2 @. NEXT := LCP ;
        LCP2 := LCP ;
    END ;

86 :      (****COLON****)

    BEGIN
        IF FCP = NIL THEN FCP := LCP1 ;
        RECINX := RECINX + 1 ;
        WITH RECST [RECINX] DO BEGIN
            LCP := LCP1 ; FST := FCP END
        END ;

87 :

    BEGIN

        WITH RECST [RECINX] DO BEGIN
            LCP1 := LCP ; FCP := FST END ;
            WHILE LCP1 <> NIL DO WITH LCP1 @ DO
                BEGIN
                    IDTYPE := RESULT ;
                    LCP1 := NEXT ;
                END ;
            RECINX := RECINX + 1 ;
            LCP1 := NIL ; LSP := NIL ;
            LCP := NIL ; LCP2 := NIL ;
        END ;

90 :      (* CASESY *)

    BEGIN
        NEW (LSP1, TAGFIELD) ;
        WITH LSP1 @ DO BEGIN
            TAGFIELD := NIL ; TAGTYPE := NIL ;
            FSTVAR := NIL ;
        END ;
        FREQVAR := LSP1 ;
    END ;

92 :

    BEGIN
        PRTErr := FALSE ;
        SEARCHID (CTYPES, LCP1) ;
        PRTErr := TRUE ;

        LID := IDENTIFIER ;
    END ;

93 :      (* COLON *)

    BEGIN
        NEW (LCP, FIELD) ;

        WITH LCP @ DO BEGIN
            NAME := LID ;
            IDTYPE := NIL ; NEXT := NIL ;
        END ;
        KOLON := TRUE ;

        ENTERID (LCP) ;
        IF FCP = NIL THEN FCP := LCP ;
        LCP2 := LCP ;
    END ;

94 :      (*** <TYPE ID> ***)

    BEGIN
        SEARCHID (CTYPES, LCP1) ;
    END ;

```

```

END LSP := LCPI @. IDTYPE
95 :      (***OF***)
BEGIN
  IF KOLON THEN KOLON := FALSE
  ELSE BEGIN
    LCP := NIL ;
    IF LCPI = NIL THEN BEGIN
      IDENTIFIER := LID ;
      SEARCHID (CTYPES, LCP) ;
    END ;
    LSP := LCPI @. IDTYPE
  END ;

  IF LSP <> NIL THEN BEGIN (* IF COLON EXISTS *)
    IF LCP <> NIL THEN LCP @. IDTYPE := LSP ;
    IF (LSP @. FORM <= SUBRANGE) OR (ISTRING(LSP)) THEN
      BEGIN
        IF LSP = REALPTR THEN ERROR (109)
        ELSE IF ISTRING(LSP) THEN ERROR (282) ;

        LSP1 @. TAGFIELDP := LCP ;
        LSP1 @. TAGTYPE := LSP
      END
    ELSE ERROR (110)
  END ; (* IF *)
  LSP2 := NIL ; LSP5 := NIL ;
  LSP := NIL ; SAVEDTAG := LSP1
END ;

96 :      FREQVAR @. FSTVAR := LSP5 ;
97, 109 :
BEGIN
  LSP1 := SAVEDTAG ;
  (* RESULT IS RETURNED BY <CONSTANT> *)
  IF COMPTYPES (RESULT, LSP1 @. TAGTYPE) THEN BEGIN
    LSP3 := LSP5 ;
    WHILE (LSP3 <> NIL) DO BEGIN
      IF LSP3 @. VARVAL = LVALU.IVAL THEN
        ERROR (158) ; (* DOUBLEVARIANT *)
      LSP3 := LSP3 @. NXTVAR
    END
  ELSE EPROR (111) ;
  NEW (LSP3, VARIANT) ;

  WITH LSP3 @ DO BEGIN
    NXTVAR := LSP5 ; SUBREC := LSP2 ;
    IF LVALU.INTVAL THEN VARVAL := LVALU.IVAL
    ELSE VARVAL := 0
  END ;

  LSP5 := LSP3 ;
  LSP2 := LSP3 ;
END ;

104 :      EMPTY := TRUE ;
97, 99 :
BEGIN
  IF EMPTY THEN BEGIN
    RCASEIND := RCASEIND @ 1 ;
    EMPTY := FALSE ;
  END
  ELSE BEGIN
    WITH RCASES [RCASEIND], INTREC @ DO BEGIN
      FSTFLD := FCP ;
      RPACK := PCKD ;
      RECVAR := FREQVAR
    END ;
    WITH RCASES [RCASEIND] DO BEGIN
      PCKD := FPCKD ;
      FCP := FIRSTFIELD ;
      FREQVAR := TAGDEFINITION ;
      LSP3 := VARHEAD ; LSP2 := INTREC
    END ;
  END ;

```

```

LSP5      := LSP3      ;
WHILE LSP3 <> NIL DO
  WITH LSP3 @ DO BEGIN
    LSP4 := SUBREC ;
    SUBREC := LSP2 ;
    LSP3 := LSP4 ;
  END ;
WITH RCASES [RCASEIND] DO
  SAVEDTAG := TAGDEFINITION ;
  RESULT := LSP2 ;
  RCASEIND := RCASEIND # 1 ;
END      (* ELSE *)
END ;

```

102 :

```

BEGIN
  EMPTY := FALSE ;
  RCASEIND := RCASEIND + 1 ;
  WITH RCASES [RCASEIND] DO BEGIN
    FPKD := PCKD ;
    TAGDEFINITION := FRECVAR ;
    FIRSTFIELD := FCP ;
    NEW (INTREC, RECORDS) ;
    VARHEAD := LSP3 ;
  END ;
  FCP := NIL ; LSP := NIL ; LCP := NIL ; LCP1 := NIL ;
  LCP2 := NIL ; PCKD := FALSE ; FRECVAR := NIL ;
END ;

```

65 :

```

BEGIN
  NEW (LSP, RECORDS) ;
  WITH LSP @ DO BEGIN
    FSTFLD := FCP ;
    RECVAR := FRECVAR ;
    RPACK := PCKD ;
  END ;
  RESULT := LSP ;
  TOP := OLDTOPS [OLDS] ;
  OLDS := OLDS # 1 ;
END ;

```

68 : (* SETS *)

```

BEGIN
  IF RESULT <> NIL THEN
    IF RESULT @. FORM > SUBRANGE THEN
      BEGIN ERROR (115) ; RESULT := NIL END
    ELSE BEGIN
      IF (RESULT=INTPTR) OR (RESULT=REALPTR) THEN
        BEGIN ERROR (115) ; RESULT := NIL END
      ELSE BEGIN
        GETBOUNDS (RESULT, LMIN, LMAX) ;
        IF (LMIN < 0) OR (LMAX > 72) THEN
          BEGIN ERROR (159) ; RESULT := NIL END
        END
      END
    END ;
  NEW (LSP, POWER) ;
  LSP @. ELSET := RESULT ;
  RESULT := LSP ;

```

END ;

71 : (* FILES *)

```

BEGIN
  NEW (LSP, FILES) ;
  WITH LSP @ DO BEGIN
    SIZE := MAXINT ;
    FILTYPE := RESULT ;
  END

```

```

END ;

      (*****
      (* SIMPLE TYPE *)
      (*****

76 :   (* USER DEFINED SCALARS *)
      BEGIN
        TTOP := TOP ;
        WHILE DISPLAY [TOP].OCCUR <> BLCK DO TOP := TOP + 1 ;
        NEW (RESULT, SCALAR, DECLARED) ;
        WITH RESULT @ DO SIZE := 1 ;

        LCPI := NIL ; LCNT := 0 ;
        IDLIST := USERSCALAR ; LSP := RESULT
      END ;

78 :
      BEGIN
        WITH RESULT @ DO FSTCONST := LCPI ;
        TOP := TTOP
      END ;

33 :
      BEGIN IF NOT RECOVERED THEN
        IF NOT (TOKEN.OP IN [PLUS, MINUS]) THEN
          ERROR (6)
        ELSE IF TOKEN.OP = PLUS THEN SIGN := POS
        ELSE SIGN := NEG
        END ;

34,40 : (* INTEGER CONSTANT *)
      IF NOT RECOVERED THEN BEGIN
        IF SIGN = NEG THEN
          WITH TOKEN.CSTADR @ DO
            BEGIN INTVAL := TRUE ; IVAL := = IVAL END ;
          RESULT := INTPTR ;
          LVALU := TOKEN.CSTADR @ ;
          SIGN := NOSIGN ; ALLOWDOTS := TRUE
        END ELSE RESULT := NIL ;

35,38 :
      IF NOT RECOVERED THEN BEGIN
        IF SIGN = NEG THEN
          WITH TOKEN.CSTADR @ DO
            BEGIN INTVAL := FALSE ; RVAL := = RVAL END ;
          RESULT := REALPTR ;
          LVALU := TOKEN.CSTADR @ ;
          SIGN := NOSIGN
        END ELSE RESULT := NIL ;

36,37 : (* IDENT *)
      IF NOT RECOVERED THEN BEGIN
        ALLOWDOTS := FALSE ;
        SEARCHID (CTYPES, KONST), LCP ) ;

        IF LCP @, KLAS = KONST THEN
          BEGIN
            ALLOWDOTS := TRUE ;
            RESULT := LCP @. IDTYPE ;
            IF SIGN <> NOSIGN THEN
              IF SIGN = NEG THEN
                WITH LCP @, LCP @, VALUES DO
                  IF CLASS = INTGR THEN IVAL := -IVAL ELSE
                  IF CLASS = REEL THEN RVAL := -RVAL
                  ELSE ERROR (105) ;

                SIGN := NOSIGN ;
                LVALU := LCP @. VALUES
            END
          ELSE
            BEGIN
              RESULT := LCP @. IDTYPE ;
              IF SIGN <> NOSIGN THEN ERROR (105)
            END
          END
        END
      END

```

```

        END
    END ELSE RESULT := NIL ;
41,39 : (* STRING CONSTANTS *)
    IF NOT RECOVERED THEN BEGIN
        LSP1 := CHARPTR ;
        IF LGTH = 1 THEN
            BEGIN RESULT := LSP1 ; ALLOWDOTS := TRUE END
        ELSE BEGIN
            NEW (RESULT, ARRAYS) ;
            WITH RESULT @ DO BEGIN
                AELTYPE := LSP1 ; PACKD := TRUE ;
                SIZE := LGTH ; LSP1 := NIL ;
                CREATE SUBRANGE (LSP1, 1, LGTH, INTPTR) ;
                INXTYPE := LSP1
            END
        END ;
        LVALU := TOKEN.CSTADR @
    END ELSE RESULT := NIL ;
81 : IF (NOT ALLOWDOTS) AND (NOT RECOVERED) THEN
        ERROR (6)
    ELSE
        BEGIN
            IF RESULT <> NIL THEN BEGIN
                IF RESULT @. FORM = SUBRANGE THEN
                    RESULT := RESULT @. RANGETYPE ;
                IF RESULT @. FORM <> SCALAR THEN
                    BEGIN ERROR (148) ; RESULT := NIL END
                ELSE IF (NOT LVALU.IVAL) AND (RESULT <> REALPTR)
                    THEN BEGIN
                        ERROR (29) ; RESULT := NIL
                    END
            END ;
            SAVIO := RESULT ;
            IF NOT RECOVERED THEN FVALU := LVALU
        END ;
82 :
    IF NOT ALLOWDOTS THEN ERROR (6) ELSE BEGIN
        ALLOWDOTS := FALSE ;
        IF RESULT <> NIL THEN BEGIN
            IF RESULT @. FORM = SUBRANGE THEN
                RESULT := RESULT @. RANGETYPE ;
            IF RESULT @. FORM <> SCALAR THEN
                BEGIN ERROR (148) ; RESULT := NIL END
            ELSE IF (NOT LVALU.IVAL) AND (RESULT <> REALPTR)
                THEN BEGIN
                    ERROR (29) ; RESULT := NIL
                END
        END ;
        FSP1 := SAVIO ;
        IF (FSP1=REALPTR) AND (RESULT=REALPTR) THEN
            BEGIN ERROR (30) ;
                RESULT := NIL
            END ;
        IF (FSP1=RESULT) AND (RESULT <> NIL) THEN
            CREATSUBRANGE (LSP, FVALU.IVAL, LVALU.IVAL, FSP1)
        ELSE
            BEGIN
                CREATSUBRANGE (LSP, #MAXINT, MAXINT, FSP1) ;
                IF (FSP1 <> NIL) AND (RESULT <> NIL) THEN
                    ERROR (107)
            END ;
        RESULT := LSP
    END ;
END ;

```

```

    (*****
    (* VAR DECLARATIONS *)
    (*****

```

110 : (***VAR***)

```
BEGIN      VNO := 0 ; INO := 0 ;
           LCP1 := NIL ; LCP2 := NIL
END ;
```

111,113,117,119 : (* IDENT *)

```
IF NOT RECOVERED THEN BEGIN
  NEW (LCP, VARS) ;
  VNO := VNO + 1 ;

  WITH LCP @ DO BEGIN
    NAME := IDENTIFIER ; NEXT := NIL ;
    IDTYPE := NIL ; VKIND := ACTUAL ;
    VLEV := LEVEL
  END ;

  ENTERID (LCP) ;

  IF LCP1 = NIL THEN LCP1 := LCP ;
  IF LCP2 <> NIL THEN LCP2 @. NEXT := LCP ;
  LCP2 := LCP
END ;
```

114,120 : (* COLON *)

```
SAV11 := LCP1 ;
```

121, 115 :

```
BEGIN
  LCP1 := SAV11 ;
  WHILE (LCP1 <> NIL) DO
    WITH LCP1 @ DO BEGIN
      IDTYPE := RESULT ;
      LCP1 := NEXT
    END
  END
```

```
END ;
```

```
(*****
(* TYPE DEFINITIONS *)
*****)
```

43,47 : (* TYPE DEFINITIONS *)

```
IF NOT RECOVERED THEN BEGIN
  NEW (LCP, TYPES) ;
  WITH LCP @ DO BEGIN
    NAME := IDENTIFIER ; IDTYPE := NIL ; NEXT := NIL
  END ; SAV11 := LCP
END ELSE SAV11 := NIL ;
```

44,48 : IF TOKEN.OP <> EQOP THEN ERROR (16) ;

45,49 :

```
IF SAV11 <> NIL THEN BEGIN
  LCP := SAV11 ;
  ENTERID (LCP) ;
  LCP @. IDTYPE := RESULT ;
  LCP := NIL ;
```

```
...HAS ANY FORWARD REFERENCE NOW
...BEEN SATISFIED ?
```

```
LCP1 := FWPTR ;
WHILE (LCP1 <> NIL) DO BEGIN
  IF LCP1 @. NAME = LCP @. NAME THEN BEGIN
    LCP1 @. IDTYPE @. ELTYPE := RESULT ;
    IF LCP1 <> FWPTR THEN
      LCP2 @. NEXT := LCP1 @. NEXT
    ELSE
      FWPTR := LCP1 @. NEXT
  END ; (* IF *)

  LCP2 := LCP1 ;
  LCP1 := LCP1 @. NEXT
END ; (* WHILE *)
```

END ;

(*****
(* CONSTANT DEFINITIONS *)
*****)

26,30 :

```
IF NOT RECOVERED THEN BEGIN
  NEW (LCP,KONST) ;
  WITH LCP @ DO BEGIN
    NAME := IDENTIFIER ;
    IDTYPE := NIL ; NEXT := NIL
  END ; SAVII := LCP
END ELSE SAVII := NIL ;
```

27,31 :

```
IF TOKEN.OP <> EQOP THEN ERROR (15) ;
```

28,32 :

```
IF SAVII <> NIL THEN BEGIN
  LCP := SAVII ;
  ENTERID (LCP) ;

  WITH LCP @ DO BEGIN
    IDTYPE := RESULT ;
    VALJES := LVALU
  END ; LCP := NIL
END ;
```

(*****
(* LABEL DECLARATIONS *)
*****)

21,23 : (**LABEL**)

```
BEGIN
  LABPTR := FSTLABPTR ;
  WHILE LABPTR <> NIL DO
    WITH LABPTR @ DO BEGIN
      IF LABNIV < LEVEL THEN GOTO 22 ;
      IF LABVAL = TOKEN.CSTADR @. IVAL THEN
        BEGIN ERROR (165) ;
              GOTO 122
            END
      ELSE LABPTR := NEXTLAB
    END ;
END ;
```

22 :

```
NEW (LABPTR) ;
WITH LABPTR @ DO BEGIN
  LABVAL := TOKEN.CSTADR @. IVAL ;
  LABNIV := LEVEL ;
  STATUS := DEFINED ;
  NEXTLAB := FSTLABPTR
END ;
```

122 : FSTLABPTR := LABPTR ;

END ;

(*****
(* ROUTINE DECLARATIONS *)
*****)

122 : FSY := PROCSY ;
125 : FSY := FUNCTIONSY ;

123,126 :

```
BEGIN
  SEARCHSECTION (DISPLAY [TOP],FNAME , LCP) ;
```

(*
*)

```
...DECIDE WHETHER IT HAS BEEN  
...DECLARED FORWARD ?
```



```

FORM := FORM AND (FSY = PROCSY)
ELSE IF CLASS = FUNC
THEN FORM := FORM AND (FSY = FUNCTIONSY)
ELSE FORM := FALSE ;

IF FORM THEN
  WITH GUNIT DO BEGIN
    FORMC := FORMC + 1 ;
    LFORM := FORMC ;
    WHILE LFORM @ NAME <> IDENTIFIER DO
      LFORM := LFORM @ NEXT ;
      LFORM @ LINE := 0 ;
    END *****
  ELSE ERROR (160) (* MULTIPLE ROUTINE ID *)
END
ELSE FORM := FALSE ;

IF FORM THEN LSP := LCP @ PARAMPTR
(* TAKE PREVIOUSLY DEFINED PARAMETERS *)
ELSE BEGIN
  IF FSY = PROCSY THEN
    NEW (LCP, PROC, DECLARED, ACTUAL)
  ELSE
    NEW (LCP, FUNC, DECLARED, ACTUAL) ;

  NEW (LSP, PARAMLIST) ;
  WITH LSP @ DO BEGIN
    SIZE := 0 ; MARKED := FALSE ;
    ANON := TRUE ; FSTPAR := NIL
  END ;

  WITH LCP @ DO BEGIN
    NAME := IDENTIFIER ;
    PARAMPTR := LSP ; (* SET PARAMETER HEAD *)
    IDTYPE := NIL ; NEXT := NIL ;

    PFLEV := LEVEL ; FORWDECL := FALSE
  END ;

  ENTERID (LCP) ;
END ;
*****
IF RECOVERED THEN
  IF FSY = PROCSY THEN LCP := UPRCPTR
  ELSE LCP := UFCTPTR ;
*****
IF LEVEL < MAXLEVEL THEN
  LEVEL := LEVEL + 1 (* INCR ROUTINE LEVEL *)
ELSE ERROR (251) ;
WITH ROUTINES [LEVEL] DO BEGIN
  OLDLEV := LEVEL ;
  OLDTOP := TOP
END ;

IF TOP < DISPLIMIT THEN
  BEGIN TOP := TOP + 1 ;
  WITH DISPLAY [TOP] DO BEGIN
    OCCUR := BLCK ; (* START OF A NEW BLOCK *)
    IF FORM THEN
      FNAME := LCP @ NEXT
    ELSE FNAME := NIL
  END
END
ELSE ERROR (250) ;

SAV5 := LCP ; (* SAV5 KEEPS IOP POINTER *)
SAV6 := LSP ; (* SAV6 KEEPS STP POINTER *)

END ;
P124, 127 :
BEGIN
  LCP := SAV5 ;

```

```

        LCPI := SAV6 @. FSTPAR ;
        WHILE LCPI <> NIL DO
            WITH LCPI @ DO BEGIN
                IF KLASS = VARS THEN VLEV := LEVEL
                ELSE PFLV := LEVEL ;
                ENTERD (LCPI ) ;
                LCPI := NEXT
            END ;
        END ;
129 : (* FUNCTION TYPE *)
        BEGIN
            SEARCHID (ETYPES, LCPI) ;
            LSP := LCPI @. IDTYPE ;
            LCP @. IDTYPE := LSP ;
            IF LSP <> NIL THEN
                IF LSP @. FORM >= POWER THEN BEGIN
                    ERROR (120) ;
                    LCP @. IDTYPE := NIL
                END
            END ;
        END ;
131 : (* FORWARDID *)
        IF LCP @. FORWDECL THEN
            ERROR (151)
        ELSE BEGIN
            LCP @. FORWDECL := TRUE ;
            GUNIT.FORWC := GUNIT.FORWC + 1 ;
            NEW (LFORW) ;
            WITH LFORW @ DO BEGIN
                NAME := LCP @. NAME ;
                LINE := CARDCNT ;
                NEXT := GUNIT.FORWP ;
                GUNIT.FORWP := LFORW
            END
        END ;
133 :
        BEGIN
            LCPI := SAV6 @. FSTPAR ;
            WHILE LCPI <> NIL DO
                WITH LCPI @ DO BEGIN
                    LLINK := NIL ; OLINK := NIL ;
                    LCPI := NEXT
                END ;
            WITH ROUTINES [LEVEL] DO BEGIN
                LEVEL := OLDLEV ;
                TOP := OLDTOP
            END
        END ;
134 :
        BEGIN
            IF FORW THEN ERROR (119) ;
            LCPI := NIL ; LCP2 := NIL ;
            IDLIST := PVARPAR ;
            LKIND := ACTUAL ; SAV9 := NIL
        END ;
138 :
        BEGIN
            IF SAV9 = NIL THEN
                SAV6 @. FSTPAR := SAV7
            ELSE
                SAV9 @. NEXT := SAV7 ;
                SAV9 := SAV8
            END ;
146 : LKIND := FORMAL ;
149 :
        BEGIN
            (* SAV7 : FIRST ID IN THE SECTION *)
            (* SAV8 : LAST ID IN THE SECTION *)
            LCPI := SAV7 ;

```

```

        WHILE LCP1 <> NIL DO
          WITH LCP1 @ DO BEGIN
            IDTYPE := RESULT ;
            LCP1 := NEXT
          END
        END
140 : BEGIN FSY1 := PROCSY ; IDLIST := PFPARAM END ;
142 : BEGIN FSY1 := FUNCTIONSY ; IDLIST := PFPARAM END ;
145 : BEGIN (*****<TYPE ID>*****)
        SEARCHID ([TYPES], LCP1) ;
        LSP := LCP1 @. IDTYPE ;
        IF LSP <> NIL THEN
          IF NOT (LSP @. FORM IN [SCALAR, SUBRANGE, POINTER]) THEN
            BEGIN ERROR (120) ;
              LSP := NIL
            END ;
          LCP := SAV7 ;
          IF LSP <> NIL THEN
            WHILE LCP <> NIL DO
              WITH LCP @ DO BEGIN
                IDTYPE := LSP ;
                LCP := NEXT
              END
            END
          END ;
10,12 :

```

```

CASE IDLIST OF
  PFPARAM : (* INITIALLY LCP1 = NIL *)
    BEGIN
      IF FSY1 = PROCSY THEN
        NEW (LCP, PROC, DECLARED, FORMAL)
      ELSE
        NEW (LCP, FUNC, DECLARED, FORMAL) ;
      WITH LCP @ DO BEGIN
        NAME := IDENTIFIER ; IDTYPE := NIL ; UNDCL := 0 ;
        LLINK := NIL ; RLINK := NIL ; NEXT := NIL ;
        PFLEV := 0
      END ;
      LCP @. PARAMPTR := NIL ;
      IF LCP1 = NIL THEN
        BEGIN SAV7 := LCP ;
          LCP1 := LCP
        END
      ELSE LCP1 @. NEXT := LCP ;
      LCP1 := LCP ;
      SAV8 := LCP
    END ;

```

```

  PFPVAPPAR :
    BEGIN
      NEW (LCP, VARS) ;
      WITH LCP @ DO BEGIN
        NAME := IDENTIFIER ; IDTYPE := NIL ;
        UNDCL := 0 ; LLINK := NIL ; RLINK := NIL ;
        NEXT := NIL ; VLEV := 0 ; VKIND := LKIND
      END ;
      IF LCP1 = NIL THEN
        BEGIN SAV7 := LCP ;
          LCP1 := LCP
        END
      ELSE LCP1 @. NEXT := LCP ;
      LCP1 := LCP ;
      SAV8 := LCP1
    END ;

```

```

FILEHEADS : ;

```

```

USERSCALAR : (* INITIALLY LCP1 = NIL *)
  IF NOT RECOVERED THEN BEGIN
    NEW (LCP, KONST) ;

```

```
NAME := IDENTIFIER ; IDTYPE := LSP ;
NEXT := LCPI ;
WITH VALUES DO BEGIN
  INTVAL := TRUE ; CLASS := INTGR ;
  IVAL := LCNT
END
END ;
ENTERID (LCP) ;
LCNT := LCNT + 1 ;
LCPI := LCP
END ;
END ;
OTHERWISE (* NO ACTION *)
END ;
END ; (* SEMANTICS *)
```

PROCEDURE MARKSET (VAR CORRSET : SETS; I : INTEGER); FORWARD ;

```
*****
*) PROCEDURE ERROR (ERRNO : INTEGER) ; (*
* THIS PROCEDURE COLLECTS ERROR AND WARNING CODES *
* TO BE PRINTED AFTER EACH LINE. *
*****
```

VAR ERRORWARN : BOOLEAN ;
BEGIN

```
ERRORWARN := ERRNO > 500 ;
ERRWARNCNT [ERRORWARN] := ERRWARNCNT [ERRORWARN] + 1 ;
IF ERRWARNCNT [FALSE] > 70 THEN
  BEGIN
    WRITELN ('***TOO MUCH SEVERE ERRORS',
            '... COMPILATION ABORTED***') ;
    GOTO 1111
  END ;
```

```
IF ERRINX = MAXERR
THEN ERRNO := 255
ELSE ERRINX := ERRINX + 1 ;
```

IF SYMSTART > LASTERR THEN

```
  BEGIN
    LASTERR := SYMSTART ;
    ERRPOS := ERRPOS + 1 ;
    IF ERRPOS >= 10
    THEN
      ERBUF [SYMSTART] := ' '
    ELSE
      ERBUF [SYMSTART] := CHR (ERRPOS + ORD ('0'))
  END ;
```

WITH ERRLIST [ERRINX] DO

```
  BEGIN
    ERRNUM := ERRNO ;
    POS := ERRPOS
  END ;
```

MARKSET (ERRDUPL, ERRNO)
END ; (* ERROR *)

PROCEDURE PRINTERRORES ;
VAR I : INTEGER ;

BEGIN

```
  WRITE (' ':24) ;
  FOR I := 1 TO BUFMAX DO
    BEGIN
      WRITE (ERBUF [ I ]) ;
      ERBUF [ I ] := ' '
    END
  WRITELN ; WRITELN ;
```

```
  ERPOS := 0 ;
  FOR I := 1 TO ERRINX DO
    WITH ERRLIST [ I ] DO
      WRITELN (ERRWARN [ERRNUM < 500], ERRNUM : 3 ,
              '...POSITION ', POS : 3, ' ***') ;
```

```
  WRITELN
  ERRINX := 0 ;
  LASTERR := 0
END ; (* PRINTERRORES *)
```

```
*****
*) PROCEDURE MARKSET (VAR CORRSET : SETS ; I : INTEGER) ; (*
```

```

*          THIS PROCEDURE PERFORMS UNION OPERATION ON AN          *
*          ARRAY OF SETS                                          *
*****

```

```

VAR J : INTEGER ;
BEGIN
  J := I DIV SETSIZE ;
  CORRSET [J] := CORRSET [J] + [ I MOD SETSIZE ]
END ; (* MARKSET *)

```

```

*****
*) FUNCTION MEMBER (CORRSET : SETS ; I : INTEGER) : BOOLEAN ; (*
*   THIS PROCEDURE IS EQUIVALENT TO IN OPERATION              *
*   BUT FOR AN ARRAY OF SETS.                                  *
*****

```

```

VAR J : INTEGER ;
BEGIN
  J := I DIV SETSIZE ;
  MEMBER := (I MOD SETSIZE) IN CORRSET [ J ]
END ; (* MEMBER *)

```

```

FUNCTION NONTERMINAL (SPIND : SPRANGE) : BOOLEAN ;
BEGIN
  WITH STABLE [ SPIND ], TABLE [ TPTYPE ] DO
    NONTERMINAL := NOT TERMINAL
  END ;

```

```

FUNCTION UNMARKED (VAR CORRSET : SETS ; M : INTEGER) : BOOLEAN ;
  VAR J, K : INTEGER ;
BEGIN
  J := M DIV SETSIZE ;
  K := M MOD SETSIZE ;
  IF K IN CORRSET [ J ]
  THEN UNMARKED := FALSE
  ELSE
    BEGIN
      UNMARKED := TRUE ;
      CORRSET [J] := CORRSET [J] + [ K ]
    END
  END
END ; (* UNMARKED *)

```

```

(*****
*) PROCEDURE EXPLAIN ;
THIS PROCEDURE PRINTS DETAILED EXPLANATION
OF ERROR CODES THAT HAS OCCURED.
*****
)

```

```

VAR I,J : INTEGER ;

BEGIN
PAGE ;
WRITELN ( ' ERROR EXPLANATIONS ' ) ;

FOR I := 1 TO 600 DO BEGIN
IF NOT MEMBER (ERRDUPL,I) THEN
J := 0 ELSE J := I ;
CASE J OF
0 : ;
1 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ERROR IN SIMPLE TYPE ' ) ;
2 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' IDENTIFIER EXPECTED ' ) ;
3 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' PROGRAM EXPECTED ' ) ;
4 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
5 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
6 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ILLEGAL SYMBOL ' ) ;
7 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ERROR IN PARAMETER LIST ' ) ;
8 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' OF EXPECTED ' ) ;
9 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
10 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ERROR IN TYPE ' ) ;
11 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
12 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
13 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
14 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
15 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
16 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
17 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
18 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;
19 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ERROR IN DECLARATION PART ' ) ;
20 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) EXPECTED ' ) ;

21 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' MULOP ASSUMED HERE ' ) ;
22 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ADDOP ASSUMED HERE ' ) ;
23 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) REPEAT EXPECTED ' ) ;
24 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) LABEL ASSUMED HERE ' ) ;
25 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) CONST ASSUMED HERE ' ) ;
26 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) VAR ASSUMED HERE ' ) ;
27 : WRITELN ( ' ***ERROR ' , J : 5, ' . . . ' ,
' ) TYPE ASSUMED HERE ' ) ;

```

```

28  :: WRITELN ( '***ERROR ' , J : 5, '...',
    '...' EXPECTED ' ) ;
50  : WRITELN ( '***ERROR ' , J : 5, '...',
'ERROR IN CONSTANT ' ) ;
51  :: WRITELN ( '***ERROR ' , J : 5, '...',
    '=' EXPECTED ' ) ;
52  :: WRITELN ( '***ERROR ' , J : 5, '...',
    'THEN' EXPECTED ' ) ;
53  : WRITELN ( '***ERROR ' , J : 5, '...',
    'UNTIL' EXPECTED ' ) ;
54  :: WRITELN ( '***ERROR ' , J : 5, '...',
    'DO' EXPECTED ' ) ;
55  : WRITELN ( '***ERROR ' , J : 5, '...',
    'TO'/'DOWNTO' TO EXPECTED ' ) ;
56  :: WRITELN ( '***ERROR ' , J : 5, '...',
    'IF' EXPECTED ' ) ;
57  : WRITELN ( '***ERROR ' , J : 5, '...',
    'FILE' EXPECTED ' ) ;
58  : WRITELN ( '***ERROR ' , J : 5, '...',
'ERROR IN FACTOR ' ) ;
59  : WRITELN ( '***ERROR ' , J : 5, '...',
'ERROR IN VARIABLE ' ) ;
101 : WRITELN ( '***ERROR ' , J : 5, '...',
'IDENTIFIER DECLARED TWICE ' ) ;
102 : WRITELN ( '***ERROR ' , J : 5, '...',
'LOW BOUND EXCEEDS HIGH BOUND ' ) ;
103 : WRITELN ( '***ERROR ' , J : 5, '...',
'IDENTIFIER IS NOT OF APPROPRIATE CLASS ' ) ;
104 : WRITELN ( '***ERROR ' , J : 5, '...',
'IDENTIFIER IS NOT DECLARED ' ) ;
105 : WRITELN ( '***ERROR ' , J : 5, '...',
'SIGN NOT ALLOWED HERE ' ) ;
109 : WRITELN ( '***ERROR ' , J : 5, '...',
'TYPE MUST NOT BE REAL ' ) ;
113 : WRITELN ( '***ERROR ' , J : 5, '...',
'INDEX TYPE MUST BE SCALAR OR SUBRANGE ' ) ;
*FUNCTION RESULT TYPE MUST BE SCALAR OR SUBRANGE ' ) ;
115 : WRITELN ( '***ERROR ' , J : 5, '...',
'BASE TYPE MUST BE SCALAR OR SUBRANGE ' ) ;
148 : WRITELN ( '***ERROR ' , J : 5, '...',
'SUBRANGE BOUNDS MUST BE SCALAR ' ) ;
149 : WRITELN ( '***ERROR ' , J : 5, '...',
'INDEX TYPE MUST NOT BE INTEGER ' ) ;
169 : WRITELN ( '***ERROR ' , J : 5, '...',
'ERROR IN BASE SET ' ) ;
173 : WRITELN ( '***ERROR ' , J : 5, '...',
'TOO MUCH STORAGE REQUIRED FOR AN ARRAY ' ) ;
201 : WRITELN ( '***ERROR ' , J : 5, '...',
'ERROR IN REAL CONSTANT : DIGIT EXPECTED ' ) ;
160 : WRITELN ( '***ERROR ' , J : 5, '...',
'PREVIOUS DECLARATION WAS NOT FORWARD ' ) ;

202 : WRITELN ( '***ERROR ' , J : 5, '...',
'String constant can't exceed source line ' ) ;
203 : WRITELN ( '***ERROR ' , J : 5, '...',
'INTEGER CONSTANT EXCEEDS RANGE ' ) ;
204 : WRITELN ( '***ERROR ' , J : 5, '...',
'REAL CONSTANT EXCEEDS RANGE ' ) ;
206 : WRITELN ( '***ERROR ' , J : 5, '...',
'INTEGER PART OF REAL CONSTANT EXCEEDS RANGE ' ) ;
250 : WRITELN ( '***ERROR ' , J : 5, '...',
'TOO MANY NESTED SCOPE OF IDENTIFIERS ' ) ;
251 : WRITELN ( '***ERROR ' , J : 5, '...',
'TOO MANY NESTED PROCEDURES/FUNCTIONS ' ) ;

255 : WRITELN ( '***ERROR ' , J : 5, '...',
'LINE ERROR LIMIT EXCEEDED ' ) ;

205 : WRITELN ( '***ERROR ' , J : 5, '...',
'NULL STRING NOT ALLOWED : ' ' ' ASSUMED ' ) ;
OTHERWISE
END (* CASE *)
END
END ; (* EXPLAIN *)

PROCEDURE KEEP ;

```


BEGIN

```
MET := FALSE ; K := 1 ;  
WHILE (NOT MET) AND (K <= OKCOUNT) DO  
  BEGIN  
    MET := (SPIND=LOSTLINKS [K].SPIND) ;  
    K := K + 1  
  END ;
```

```
IF NOT MET THEN  
  BEGIN
```

```
    OKCOUNT := OKCOUNT + 1 ;  
    IF OKCOUNT = 1 THEN  
      LOSTLINKS [OKCOUNT].SPIND := SAVEDSP  
    ELSE  
      LOSTLINKS [OKCOUNT].SPIND := SPIND ;
```

```
  WITH LOSTLINKS [OKCOUNT] DO
```

```
    BEGIN  
      NEW (TEMP) ;  
      TEMP @ := STACKTOP@  
    END ;
```

```
  IF OKCOUNT > 1 THEN  
    LOSTLINKS [OKCOUNT-1].TEMP  
    @, PREVIOUS :=  
    LOSTLINKS [OKCOUNT].TEMP
```

```
END ;
```

\$PAGE

```
(*****  
*) PROCEDURE SCAN ; (*  
*) THIS ROUTINE IS THE LEXICAL ANALYZER. THE CLASS (*  
*) AND VALUE (IF ANY) OF THE CURRENT SYMBOL IS (*  
*) RETURNED IN THE RECORD VARIABLE 'TOKEN'. (*  
*) SCAN IS A HAND WRITTEN SCANNER EXCEPT FOR (*  
*) IDENTIFIERS AND KEYWORDS FOR WHICH A SET OF (*  
*) TRANSITION TABLES USED. (*  
*)  
*****)
```

```
LABEL 111 ; (* RETURN ADDRESS *)  
VAR I,J,K ; (* COUNTERS *)  
XPO,IVAL : INTEGER; (* FOR CONSTANT CONVERSION *)  
P,RVAL,FAC : REAL ; (* FOR CONSTANT CONVERSION *)  
SIGN : ; (* FOR NEGATIVE CONSTANTS *)  
FINISHED : ; (* FOR LOOP CONTROL *)  
ERR : BOOLEAN; (* ERROR FLAG FOR CONSTANTS *)  
BAZ,STATE : INTEGER; (* FOR ID RECOGNITION *)
```

```
PROCEDURE PRINTLINE ;  
VAR I : INTEGER ;  
  
BEGIN WRITE ( ' ' ) ;  
IF LNEST THEN  
BEGIN  
LNEST := FALSE ;  
WRITE (NESTINCR : 5)  
END  
ELSE WRITE ( ' ' : 6 ) ;  
IF RNEST THEN  
BEGIN  
RNEST := FALSE ;  
WRITE (NESTDECR : 5)  
END  
ELSE WRITE ( ' ' : 6 ) ;  
WRITE ( ' ' ) ;  
  
WRITE ( ' ', CARDCNT : 5, ' ' ) ;  
FOR I := 1 TO BUFMAX DO  
IF INBUF [ I ] IN [ 'A'..'Z' ] THEN  
WRITE (LCASE [ INBUF [ I ] ] )  
ELSE WRITE ( INBUF [ I ] ) ;  
  
WRITELN  
END ; (* PRINTLINE *)
```

```
PROCEDURE READLINE ;  
VAR I : INTEGER ;  
  
BEGIN  
CHCNT := 0 ;  
IF CARDCNT > 0 THEN PRINTLINE ;  
  
IF ERRPINX > 0 THEN  
PRINTERRORS ;  
  
READLN (INPUT) ;  
IF EOF (INPUT)  
THEN  
BEGIN  
TOKEN.CLASS := EOFPSM ;  
IF UNCLOSED <> NONE THEN  
BEGIN  
WRITELN (ABORTMES [UNCLOSED = COMMENT] ;  
'STARTED AT LINE ',ABORTLINE : 3) ;  
GOTO 1111  
END ;  
IF CARDCNT = 0 THEN GOTO 1111 ELSE GOTO 111  
END  
ELSE  
BEGIN  
BUFLGTH := 0 ;  
WHILE (NOT EOLN (INPUT)) AND (BUFLGTH < BUFMAX) DO
```

```

        BEGIN
            BUFLGTH := BUFLGTH + 1 ;
            READ (INBUF [BUFLGTH])
        END ;

        I := BUFLGTH + 1 ;
        WHILE ( I <= BUFMAX ) DO
            BEGIN
                INBUF [I] := ' ' ;
                I := I + 1
            END ;
        END ;

        CARDCNT := CARDCNT + 1
    END ; (* READLINE *)

```

```

FUNCTION NEXTCH : CHAR ;
BEGIN
    IF CHCNT >= BUFLGTH
    THEN
        BEGIN
            NEXTCH := ' ' ;
            READLINE
        END
    ELSE
        BEGIN
            CHCNT := CHCNT + 1 ;
            NEXTCH := INBUF [CHCNT]
        END
    END ; (* NEXTCH *)

```

```

FUNCTION NEXTCHAR : CHAR ;
BEGIN
    IF CHCNT >= BUFLGTH
    THEN
        NEXTCHAR := ' '
    ELSE
        BEGIN
            CHCNT := CHCNT + 1 ;
            NEXTCHAR := INBUF [CHCNT]
        END
    END ; (* NEXTCHAR*)

```

```

BEGIN (* SCAN *)
    REPEAT
        WHILE ( CH = ' ' ) DO
            BEGIN
                IF CHCNT >= BUFLGTH THEN
                    READLINE ;

                CHCNT := CHCNT + 1 ;
                CH := INBUF [CHCNT]
            END ;

            SYMSTART := CHCNT ;
            FINISHED := TRUE ;

            CASE CH OF
                'A', 'B', 'C', 'D', 'E', 'F', 'G', 'H',
                'I', 'J', 'K', 'L', 'M', 'N', 'O', 'P',
                'Q', 'R', 'S', 'T', 'U', 'V', 'W', 'X',
                'Y', 'Z' :
                    BEGIN (* ID RECOGNITION *)
                        IDSTR := DUMMYID ;
                        STATE := ORD (CH) ;
                        IDSTR [ 1 ] := CH ;

                        CH := NEXTCHAR ;
                        I := 1 ;
                        BAZ := BASE [STATE] + ORD (CH) ;
                    END
            END
        UNTIL FINISHED
    END ;

```

```

WHILE STATE = CHECK [BAZ] DO
BEGIN
  I := I + 1 ;
  IDSTR [ I ] := CH ;
  CH := NEXTCHAR ;
  STATE := NEXT [BAZ] ;
  BAZ := BASE [STATED] + ORD (CH) ;
END ;

```

```

IF CH IN IDCHARS THEN
BEGIN
  STATE := 1 ;
  REPEAT
    I := I + 1 ;
    IDSTR [ I ] := CH ;
    CH := NEXTCHAR ;
  UNTIL NOT (CH IN IDCHARS) ;
END ;

```

```

IDENTIFIER := IDSTR ;
TOKEN := DEFAULT [STATED] ;
CASE TOKEN, CLASS OF
  RECORDSY ; BEGINSY ;
  REPEATSY ; CASESY ;
  IF (STABLE[STACKTOP], PTR], IPTYPE < 75)
  AND (TOKEN, CLASS = CASESY) THEN ELSE
  BEGIN
    NEST := NEST + 1 ;
    IF NOT LNEST THEN BEGIN
      NESTINC := NEST ; LNEST := TRUE END
    END ;
  UNTILSY ; ENDSY ;
  BEGIN
    RNEST := TRUE ;
    NESTDEC := NEST ;
    NEST := NEST - 1 ;
  END ;
  OTHERWISE
  END (* CASE *)
END ;

```

```

'0', '1', '2', '3', '4', '5', '6', '7',
'8', '9', :

```

```

BEGIN (* DIGIT CONSTANTS *)
  TOKEN, CLASS := INTCONST ;
  RVAL := 0 ; ERR := FALSE ;
  REPEAT
    RVAL := RVAL * 10 + (ORD (CH) - ORD ('0')) ;
    CH := NEXTCHAR ;
  UNTIL NOT (CH IN DIGITS) ;
  IF CH = '.' THEN
  BEGIN
    CH := NEXTCHAR ;
    IF (CH = '.' ) OR (CH = ')')
    THEN
    BEGIN
      CHCNT := CHCNT + 1 ;
      CH := '.' ;
    END
    ELSE
    BEGIN
      TOKEN, CLASS := REALCONST ;
      IF NOT (CH IN DIGITS)
      THEN ERROR (201)
      ELSE
      BEGIN
        R := 1 ;
        REPEAT
          R := R * 10 ;
          RVAL := RVAL * 10 +
            (ORD (CH) - ORD ('0')) ;
          CH := NEXTCHAR ;
        UNTIL NOT (CH IN DIGITS) ;
        RVAL := RVAL / R ;
      END
    END
  END
END ;

```

```

IF (CH = 'E') THEN
  BEGIN
    TOKEN.CLASS := REALCONST ;
    SIGN := FALSE ;

    CH := NEXTCHAR ;
    IF CH = '+' THEN
      THEN CH := NEXTCHAR
    ELSE
      IF CH = '-' THEN
        BEGIN
          SIGN := TRUE ;
          CH := NEXTCHAR
        END ;

    R := 0 ;
    IF NOT (CH IN DIGITS)
    THEN ERROR (201)
    ELSE
      REPEAT
        R := R * 10 + (ORD (CH) - ORD ('0')) ;
        CH := NEXTCHAR
      UNTIL NOT (CH IN DIGITS) ;

    IF R > MAXINT
    THEN ERROR (206)
    ELSE
      IF R <> 0 THEN
        BEGIN
          EXPO := TRUNC (R)
          R := 1 ; FAC := 10 ;

          REPEAT
            IF ODD (EXPO) THEN
              R := R * FAC ;
              IF FAC < SQRTREAL
              THEN
                BEGIN
                  FAC := SQR (FAC) ;
                  ERR := MAXREAL /
                    FAC < R
                END
              ELSE ERR := TRUE ;

            EXPO := EXPO DIV 2
          UNTIL (EXPO = 0) OR (ERR) ;

          IF EXPO <> 0
          THEN ERROR (207)
          ELSE
            IF SIGN
            THEN RVAL := RVAL / R
            ELSE
              IF MAXREAL / R > RVAL
              THEN RVAL := RVAL * R
              ELSE ERROR (207)

            END
          END ;

    END ;

IF TOKEN.CLASS = INTCONST
THEN
  BEGIN
    NEW (KONSPTR, INTGR) ;
    KONSPTR 2. IVAL := TRUE ;

    IF RVAL > MAXINT
    THEN BEGIN ERROR (203) ;
      KONSPTR 3. IVAL := MAXINT
    END
    ELSE KONSPTR 3. IVAL := TRUNC (RVAL) ;
    TOKEN.CSTADR := KONSPTR
  END
ELSE
  BEGIN
    NEW (KONSPTR, REEL) ;
    KONSPTR 2. IVAL := FALSE ;
    IF ERR THEN RVAL := 0 ;
    KONSPTR 3. RVAL := RVAL ;
    TOKEN.CSTADR := KONSPTR
  END
END ;

```

```

BEGIN
  LGTH      := 0 ;
  STRBUF    := DUMMYSTR ;
  ABORTLINE := CARDCNT ;
  ERR       := FALSE ;
  UNCLOSED  := STRINGS ;

  REPEAT
    REPEAT
      CHCNT := CHCNT + 1 ;
      LGTH  := LGTH + 1 ;
      CH    := INBUF [CHCNT] ;
      IF NOT ERR THEN
        STRBUF [LGTH] := INBUF [CHCNT]
      UNTIL ( CH = '' ) OR (CHCNT > BUFLGTH)
        OR (LGTH=STRGLGTH) ;

      IF CH = ''
      THEN CH := NEXTCHAR
      ELSE
        BEGIN
          IF NOT ERR THEN
            ERROR (202) ;
          CH := '' ;
          IF CHCNT > BUFLGTH THEN
            READLINE ;
          ERR := TRUE
        END
      UNTIL CH <> '' ;

  UNCLOSED := NOVE ;
  STRBUF [ LGTH ] := '' ;
  LGTH := LGTH - 1 ;

  IF LGTH IN [ 0, 1 ]
  THEN
    BEGIN
      IF LGTH = 0 THEN
        BEGIN LGTH := 1 ;
          STRBUF [ 1 ] := '' ;
          ERROR (205) ;
        END ;
      TOKEN.CLASS := CHARCONST ;
      NEW (KONSPTR) ;
      WITH KONSPTR DO
        CLASS := INTGR ;
        IVAL  := 00 (STRBUF [1]) ;
        INTVAL := TRUE ;
        END ;
      TOKEN.CSTADR := KONSPTR
    END
  ELSE
    BEGIN
      TOKEN.CLASS := STRINGCONST ;
      NEW (KONSPTR, STRING) ;
      IF NOT ERR
      THEN
        WITH KONSPTR DO
          BEGIN INTVAL := FALSE ;
            SLGTH := LGTH ;
            SVAL  := STRBUF
          END
        ELSE
          WITH KONSPTR DO
            BEGIN INTVAL := FALSE ;
              SLGTH := 0 ;
              SVAL  := DUMMYSTR
            END ;
          TOKEN.CSTADR := KONSPTR
        END
    END ;
  END ;

  ... ;

  BEGIN
    CH := NEXTCHAR ;
    IF CH = ''
    THEN
      BEGIN
        TOKEN.CLASS := ASGNOP ;

```

```

        END CH          := NEXTCHAR
    ELSE
        TOKEN.CLASS    := COLON
    END ;

'<' :
BEGIN
    CH := NEXTCHAR ;
    TOKEN.CLASS := RELOP ;
    IF CH = '='
    THEN
        BEGIN
            TOKEN.OP := LEOP ;
            CH := NEXTCHAR
        END
    ELSE
        IF CH = '>'
        THEN
            BEGIN
                TOKEN.OP := NEOP ;
                CH := NEXTCHAR
            END
        ELSE
            TOKEN.OP := LTOP
        END ;
END ;

'>' :
BEGIN
    CH := NEXTCHAR ;
    TOKEN.CLASS := RELOP ;
    IF CH = '='
    THEN
        BEGIN
            TOKEN.OP := GEOP ;
            CH := NEXTCHAR
        END
    ELSE
        TOKEN.OP := STOP
    END ;
END ;

'(' :
BEGIN
    CH := NEXTCHAR ;
    IF CH = '('
    THEN
        BEGIN
            TOKEN.CLASS := LBRACKET ;
            CH := NEXTCHAR
        END
    ELSE
        IF CH = '*'
        THEN
            BEGIN
                UNCLOSED := COMMENT ;
                ABORTLINE := CARDCNT ;

                REPEAT
                    REPEAT
                        UNTIL NEXTCH = '*'
                    UNTIL NEXTCH = ')' ;

                UNCLOSED := NONE ;
                CH := NEXTCHAR ;
                FINISHED := FALSE ;
            END
        ELSE
            TOKEN.CLASS := LPARENT
        END ;
END ;

'.' :
BEGIN
    CH := NEXTCHAR ;
    IF CH = '.'
    THEN
        BEGIN
            TOKEN.CLASS := DOTDOT ;
            CH := NEXTCHAR
        END
    ELSE

```

```

IF CH = ')'
THEN
  BEGIN
    TOKEN.CLASS := LBRACKET ;
    CH := NEXTCHAR ;
  END
ELSE
  BEGIN
    TOKEN.CLASS := PERIOD ;
  END ;

```

```

's' : BEGIN I := 1 ; CH := NEXTCHAR ;
      IDSTR := DUMMYID ;
      WHILE CH IN 'A-Z' DO
        BEGIN IDSTR ID := CH ; I := I+1 ; CH := NEXTCHAR END ;
      FINISHED := FALSE ; NAME := IDSTR ;
      PRINTABLE := NAME = 'PRINTABLES' ;
      END ;

```

```

'.' : BEGIN I := 1 ; CH := NEXTCHAR ;
      IDSTR := DUMMYID ;
      WHILE CH IN 'A-Z' DO
        BEGIN IDSTR ID := CH ; I := I+1 ; CH := NEXTCHAR END ;
      FINISHED := FALSE ; NAME := IDSTR ;
      PRINTABLE := NAME = 'PRINTABLES' ;
      END ;

```

```

BEGIN
  TOKEN := CHRCLASS ECHO ;
  CH := NEXTCHAR ;
END ;

```

OTHERWISE

```

BEGIN
  ERROR (6) ;
  CH := NEXTCHAR ;
  FINISHED := FALSE ;
END

```

END (* CASE *)

UNTIL FINISHED ;

END ; (* SCAN *)

111:


```

(*****
*) PROCEDURE RECOVER ( SPINX : INTEGER ) ; (*
*) THIS PROCEDURE TAKES CONTROL WHEN A SYNTAX ERROR (*
*) OCCURS TO RECOVER IT. THE ALGORITHM DEPENDS ON (*
*) IRON'S TECHNIQUE FOR "TOP-DOWN" PARSERS WITHOUT (*
*) BACKUP. (*
*)
*****

```

```

LABEL 22,56 ;
VAR

```

```

SUCCESSOR : INTEGER ; (* SUCCESSIVE STABLE ENTRY *)
FOUND : (* DEPENDS ON FUNCTION FIND *)
FND : (* DEPENDS ON FUNCTION FIND *)
MET : (* TEMPORARY *)
REACHED : BOOLEAN ; (* LOOP CONTROL FLAG *)
STEMP : (* TEMPORARY PARSE STACK *)
CURRENT : (*
CURRPOS : LINKS ; (* POINTERS
INSCONTRL : (* CONTROLS INFINITE LOOPS
FINDCONTRL : SETS ; (* CONTROLS INFINITE LOOPS

```

```

FUNCTION FIND (SPINX : INTEGER) : BOOLEAN ;

```

```

VAR FOUND : BOOLEAN ;
BEGIN

```

```

FOUND := FALSE ;
WITH STABLE [ SPINX ], TTABLE [ TPTYPE ] DO
CASE TERMINAL OF

```

```

FALSE :
BEGIN

```

```

IF UNMARKED (FINDCONTRL, TPTYPE) THEN
FOUND := FIND (POINTER) ;

```

```

IF NOT FOUND THEN
IF (SUC <> OK) AND (SUC > SPINX) THEN
FOUND := FIND (SUC) ;

```

```

FIND := FOUND
END ;

```

```

TRUE :
BEGIN

```

```

IF CLASS = TOKEN.CLASS
THEN FIND := TRUE
ELSE

```

```

BEGIN
IF (ALT <> FAIL) AND (ALT <> OK) THEN
FOUND := FIND (ALT) ;

```

```

IF NOT FOUND THEN
IF (SUC <> OK) AND (SUC > SPINX) THEN
FOUND := FIND (SUC) ;

```

```

FIND := FOUND
END ;
END ;

```

```

END (* CASE *)
END ; (* FIND *)

```

```

PROCEDURE STRING ( SPINX : INTEGER ) ;

```

```

BEGIN
WITH STABLE [ SPINX ], TTABLE [ TPTYPE ] DO
CASE TERMINAL OF

```

```

FALSE :
BEGIN

```

```

STRING (POINTER) ;
IF SUC <> OK THEN
STRING (SUC) ;

```

```

END ;
TRUE :

```

```

IF (ALT = FAIL) OR (SUC = OK)
THEN

```

```

      BEGIN
        RECIND := RECIND + 1 ;
        RECOVERY [RECIND].CLASS := CLASS ;
      END
    IF SUC <> OK THEN
      STRING (SUC)
    END
  ELSE
    IF ALT <> OK THEN
      STRING (ALT)
    END
  END ;
  (* CASE
  (* STRING
  *)

FUNCTION INSERT (SUCCESSOR : SPRANGE) : BOOLEAN ;
  LABEL 33
  VAR REACHED : BOOLEAN ;
      SAVED : INTEGER ;
  BEGIN
    REACHED := FALSE ;
    FINDCONTRL := EMPTYSET ;
    WITH STABLE [SUCCESSOR], TABLE [TPTYPE] DO
      CASE TERMINAL OF
        TRUE :
          IF CLASS = TOKEN.CLASS
          THEN
            BEGIN
              REACHED := TRUE ;
              RECIND := RECIND + 1 ;
              RECOVERY [RECIND].CLASS := TOKEN.CLASS
            END
          ELSE
            IF (ALT = FAIL) OR (ALT = OK)
            THEN GOTO 33
            ELSE
              IF FIND (ALT)
              THEN
                BEGIN
                  SAVED := RECIND ;
                  REACHED := INSERT (ALT) ;
                  IF NOT REACHED THEN
                    BEGIN
                      RECIND := SAVED + 1 ;
                      RECOVERY [RECIND].CLASS := CLASS ;
                      IF (SUC <> OK) AND
                        (SUC > SUCCESSOR) THEN
                        REACHED := INSERT (SUC)
                      END
                    END
                  END
                ELSE
                  33 :
                    BEGIN
                      RECIND := RECIND + 1 ;
                      RECOVERY [RECIND].CLASS := CLASS ;
                      IF (SUC <> OK) AND (SUC > SUCCESSOR) THEN
                        REACHED := INSERT (SUC)
                      END
                    END
                END
              END
            END
          END
        FALSE :
          IF NOT FIND (POINTER)
          THEN
            BEGIN
              STRING (POINTER) ;
              IF SUC <> OK THEN
                REACHED := INSERT (SUC)
              END
            END
          ELSE
            IF UNMARKED (INSCONTRL, TPTYPE) THEN
              BEGIN
                FINDCONTRL := EMPTYSET ;
                MARKSET (FINDCONTRL, TPTYPE) ;
                IF (SUC <> OK) AND (SUC > SUCCESSOR)
                THEN
                  IF FIND (SUC) THEN BEGIN
                    STRING (POINTER) ;
                    REACHED := INSERT (SUC) END
                  ELSE REACHED := INSERT (POINTER)
                END
              END
            END
          END
        END
      END
    END
  END

```

```

ELSE REACHED := INSERT ( POINTER ) ;
IF NOT REACHED THEN
  IF (SUC <> OK) AND (SUC > SUCCESSOR) TH
    REACHED := INSERT (SUC)
END
ELSE
  BEGIN
    SAVED := RECIND ;
    FINDCTRL := EMPTYSET ;
    IF (SUC <> OK) AND (SUC > SUCCESSOR) THEN
      IF FIND (SUC) THEN
        BEGIN
          STRING ( POINTER ) ;
          REACHED := INSERT (SUC);
          IF NOT REACHED THEN
            BEGIN
              RECIND := SAVED ;
              REACHED := INSERT ( POINTER )
            END
          END
        END
      ELSE
        END
    ELSE
      END
      STRING ( POINTER )
    END
  END ; (* CASE *)
  INSERT := REACHED
END ; (* INSERT *)

```

SPACE

BEGIN (* RECOVER *)

```
(*****  
* EACH MEMBER OF THE PARSE STACK SHOWS AN INCOMPLETE BRANCH. *  
* THE SYMBOL THAT CAUSED ERROR MUST BE IN THE INCOMPLETE PART *  
* OF ONE OF THESE INCOMPLETE BRANCHES, OR IT MUST BE IN ONE OF *  
* THE IMMEDIATE SUCCESSORS OF SPINX. *  
* IF NOT, JUST SKIP THE CURRENT TOKEN AND TRY NEXT UNTIL ONE *  
* IS FOUND. *  
* *****)
```

```
MET := FALSE ;  
IF OKCOUNT > 0 THEN  
  BEGIN  
    FOR K := 1 TO OKCOUNT DO  
      WITH LOSTLINKS [K], TEMP 0 DO  
        IF PTR IN [16,17,18,19] THEN  
          BEGIN MET := TRUE ;  
            FINDCONTRL := EMPTYSET ;  
            WITH STABLE [LOSTLINKS[K], SPIND], TTABLE [TPTYPE] DO  
              IF SUC <> OK THEN IF FIND (SUC) THEN  
                BEGIN  
                  STACKTOP.PTR := 13 ;  
                  PREVIOUS := STACKTOP ;  
                  STACKTOP := LOSTLINKS[K].TEMP ;  
                  SPIND := LOSTLINKS[K].SPIND ;  
                  SUCCESSOR := SPIND ;  
                  OKCOUNT := 0 ;  
                  ALTFLAG := FALSE ;  
                  GOTO 22  
                END  
              END ;  
            IF MET THEN GOTO 56 ;  
            WITH LOSTLINKS [OKCOUNT], TEMP 0, PREVIOUS := STACKTOP ;  
            BEGIN  
              STACKTOP := TEMP ;  
              SPINX := SPIND ;  
            END ;  
            ALTFLAG := FALSE ; OKCOUNT := 0 ;  
            SPIND := SPINX  
          END ;
```

56 :

```
REPEAT  
  FINDCONTRL := EMPTYSET ;  
  WITH STABLE [SPINX], TTABLE [TPTYPE] DO  
    IF SUC <> OK THEN IF FIND (SUC) THEN  
      BEGIN  
        SUCCESSOR := SPINX ;  
        SPIND := SPINX ;  
      END  
    GOTO 22  
  END ;
```

```
CURRPOS := STACKTOP ;  
FOUND := FALSE ;  
FNO := FALSE ;  
CURRENT := STACKTOP ;
```

```
IF STACKTOP = NIL THEN  
  BEGIN  
    WITH STABLE [SPINX], TTABLE [TPTYPE] DO  
      FOUND := CLASS = TOKEN.CLASS ;  
      SUCCESSOR := SPINX  
    END ;
```

```
WHILE (NOT FOUND) AND (CURRENT <> NIL) DO  
  BEGIN  
    WITH STABLE [CURRENT], PTR DO  
      BEGIN  
        FINDCONTRL := EMPTYSET ;  
        IF SUC <> OK THEN BEGIN  
          IF SUC > CURRENT, PTR THEN  
            MARKSET (FINDCONTRL, TPTYPE) ;  
            FOUND := FIND (SUC) END  
        END ;  
      IF FOUND THEN  
        BEGIN
```

```
STRING (SPINX) :
```

```

        END SUCCESSOR := STABLE [CURRENT @, PTR], SUC
        CURRPOS := CURRENT ; CURRENT := CURRENT @.PREVIOUS
    END ; CURRENT := STACKTOP ;

    WHILE (NOT FOUND) AND (CURRENT <> NIL) DO
        BEGIN
            WITH STABLE [CURRENT @, PTR], [TABLE [TTYPE]] DO
            IF ((POINTER = 13) OR (POINTER = 122)) AND (SPINX >= .50) THEN
                ELSE FOUND := FIND (POINTER) ;

                IF FOUND THEN
                    BEGIN
                        STACKTOP := CURRENT ;
                        WITH STABLE [CURRENT @, PTR], [TABLE [TTYPE]] DO
                            BEGIN SPIND := POINTER ;
                                SUCCESSOR := POINTER ; END
                            END ; CURRPOS := CURRENT ;
                        CURRENT := CURRENT @.PREVIOUS
                    END ;

                IF NOT FOUND
                THEN
                    BEGIN
                        ERROR (5) ;
                        SCAN ;

                        IF TOKEN.CLASS = EOFPGM THEN
                            BEGIN
                                WRITELN ('**UNEXPECTED EOF INDICATOR **',
                                    '***COMPILATION ABORTED ***') ;
                                GOTO 1111
                            END
                        END
                    END
                UNTIL FOUND ;

```

```

(*****
*
* THE CURRPOS NOW POINTS TO THE STABLE ENTRY WHICH ACCEPTS
* LAST SCANNED TOKEN, SAY T, AS A MEMBER OF ITS DEFINITION
* SO REMAINING STATEMENTS OF RECOVER DETERMINES A TERMINAL
* STRING AND INSERTS IT JUST BEFORE T, SO THAT CONTINUATION
* OF THE PARSE CAN NOW BE CORRECTLY HANDLED
*
*****

```

```

STEMP := STACKTOP ;

```

```

WHILE (STEMP <> CURRPOS) DO
    WITH STEMP @ DO
        BEGIN
            WITH STABLE [PTR] DO
                IF SUC <> OK THEN
                    STRING (SUC) ;
                STEMP := PREVIOUS
            END ;

```

```

22 : INSCONTRL := EMPTYSET ;
    WITH STABLE [SUCCESSOR], [TABLE [TTYPE]] DO
    IF TERMINAL THEN IF SUC < SUCCESSOR THEN
        BEGIN
            RECIIND := RECIIND + 1 ;
            RECOVERY [RECIIND].CLASS := CLASS ;
            SUCCESSOR := SUC
        END ;

    REACHED := INSERT (SUCCESSOR) ;
    IF NOT REACHED THEN
        BEGIN
            WRITELN ('*** COMPILER ERROR ***') ;
            STACKDUMP ;
            WRITELN ('*SPIND :', SPIND) ; HALT
        END
    END ;
END ; (* RECOVER *)

```

```

PROCEDURE PARSEERROR ( SPINX : SPRANGE ) ;
VAR USEDMSGGS      : SETS ;
    I, ERRNO       : INTEGER ;
    CURRPOS        : LINKS ;

BEGIN
    USEDMSGGS      := EMPTYSET ;
    RECIND := 0 ;
    RECOVAR (SPINX) ;
    IF RECIND > 1 THEN
    FOR I := 1 TO RECIND - 1 DO
        BEGIN
            WITH RECOVERY [ I ] DO
                BEGIN
                    ERRNO := MISSING (CLASS) ;

                    IF (ERRNO > 0) AND (UNMARKED (USEDMSGGS, ERRNO)) THEN
                        ERROR (ERRNO)
                END ;
            END ;
        END ;
    END ;

    CURRPOS := STACKTOP ;
    REPEAT
        WITH CURRPOS & DO
            BEGIN
                WITH STABLE [ PTR ] DO
                    ERRNO := NONTRMSGGS [TPTYPE] ;
                    IF ERRNO <> 0 THEN
                        IF UNMARKED (USEDMSGGS, ERRNO) THEN
                            ERROR (ERRNO)
                        CURRPOS := PREVIOUS
                    END ;
                UNTIL CURRPOS = NIL
            END ;
        (* PARSEERROR *)
    END ;

```

SPAGE

```
PROCEDURE PUSH (ELEMENT : SPRANGE) ;
VAR SAV : LINKS ;
BEGIN
  SAV := STACKTOP ;
  NEW (STACKTOP) ;
  WITH STACKTOP @ DO
    BEGIN
      PTR := ELEMENT ;
      PREVIOUS := SAV ;
    END ;
  SAVE := STACKTOP ;
END ; (* PUSH *)
```

```
PROCEDURE POP ;
VAR I : INTEGER ;
```

```
BEGIN IF ALTFLAG THEN KEEP ;
  SAVE := STACKTOP ;
  STACKTOP := STACKTOP @. PREVIOUS ;
  (* MARK ( I ) ;
  RELEASE ( I ) ; *) DISPOSE (SAVE)
END ; (* PUSH *)
```

```
PROCEDURE SUCCESS ( MODE : INTEGER ) ;
```

```
BEGIN
  WHILE (STABLE [SPIND], SUC = OK) OR (MODE = 1) DO
    BEGIN
      MODE := 0 ;
      IF STACKTOP = NIL THEN
        BEGIN
          IF TOKEN.CLASS <> EOF THEN
            WRITELN ('**EOF EXPECTED**') ;
          GOTO 1111 ;
        END
      SAVEDSP := SPIND ;
      SPIND := STACKTOP @. PTR ;
      (* WRITELN ('***SEMANTICS(' , SPIND : 5, ')') ; *)
      SEMANTICS (SPIND) ;
      POP
    END ; (* WHILE *)
  SPIND := STABLE [SPIND], SUC
END ; (* SUCCESS *)
```

SPAGE

```
BEGIN (* MAIN PROGRAM *)
PAGE ;
INITSCAN ;
INITPARSE ; STACKTOP := NIL ;
ENTERSTDIDS ;
SCAN ; SPIND := 1 ;
```

```
(*****
*
* THE PARSING ALGORITHM USED IN THIS IMPLEMENTATION IS A
* TOPDOWN TECHNIQUE WITHOUT BACKUP. IT DEPENDS ON THE
* ALGORITHM BY CHEATNEM AND SATTLEY
*
*****)
```

```
WHILE FOREVER DO
  BEGIN
```

```
    WHILE NONTERMINAL (SPIND) DO
      WITH STABLE [ SPIND ], TTABLE [ TPTYPE ] DO
        BEGIN
          PUSH (SPIND) ;
          SPIND := POINTER
        END ;
```

```
    WITH STABLE [ SPIND ], TTABLE [ TPTYPE ] DO
      IF CLASS = TOKEN.CLASS
      THEN
```

```
        BEGIN
          WRITELN ('***SEMANTICS (' , SPIND:5, ')') ; (*)
          SEMANTICS (SPIND) ;
```

```
          IF RECIND > 0
          THEN
            IF RIND = RECIND
            THEN
              BEGIN
```

```
                RECIND := 0 ;
                RIND := 0 ;
                RECOVERED := FALSE ;
```

```
                SCAN
```

```
            END
          ELSE
```

```
            BEGIN
              RIND := RIND + 1 ;
              IF RIND = RECIND THEN
                RECOVERED := FALSE ;
                TOKEN := RECOVERY [RIND]
```

```
            END
```

```
          ELSE
```

```
            SCAN ;
            IF (TOKEN.CLASS = EOFPGM) AND
              (STACKTOP <> NIL) THEN BEGIN
              WRITELN ('**UNEXPECTED EOF INDICATOR.**',
                ' COMPILATION ABORTED ***') ;
              GOTO 1111
```

```
            END ;
```

```
            ALTFLAG := FALSE ;
            OKCOUNT := 0 ;
            SUCCESS ( 0 )
```

```
          END
```

```
        ELSE
          IF ALT = FAIL
          THEN
```

```
            BEGIN
              PARSERROR (SPIND) ;
              RIND := 1 ;
              RECOVERED := TRUE ;
              TOKEN := RECOVERY [RIND]
```

```
            END
```

```
          ELSE
```

```
            IF ALT = OK
            THEN BEGIN ALTFLAG := TRUE ;
              SUCCESS ( 1 )
```

```
            END
```


END ; (* FOREVER LOOP *) := ALT

WRITELN ('WRITELN (' , CARD CNT : 5 , ' RECORDS ARE PROCESSED . ') ;
WRITELN (' ' , ERRWARN CNT [FALSE] : 5 , ' ERRORS DETECTED . ') ;
EXPLAIN
END . (* OF PASCAL ANALYZER *)

CONTAINS 4792 LINES AND 44 PROCEDURES

THE ASCII FILES

TIME USED: 237 SECONDS

NO ERRORS NO WARNINGS 15.89 I#BANK 32768 D#BANK

06/27/83 21:01:12 (0)

32537 TIME: 114.465 STORAGE: 5164/14/1/8236

06/27/83 21:03:48

001000 047441 19746 IBANK WORDS DECIMAL
050000 150450
006640

VT 3MAIN9 001000 047441 050000 150450

0(1)	001000	001174			07 AUG 78	09
0(1)	001175	001570			07 AUG 78	09
					20 DEC 78	17
0(1)	001571	001770			07 AUG 78	09
					17 JUL 78	13
0(1)	001771	002241	0(0)	050000 050023	23 APR 82	12
0(1)	002242	002261	0(0)	050024 050025	23 JUL 80	13
0(3)	002262	002265				
0(1)	002266	002305	0(2)	050025 050026	18 MAR 80	14
0(3)	002306	002312				
0(1)	002313	002363	0(2)	050027 050027	19 MAR 80	15
0(3)	002364	002377				
			0(2)	050030 050065	30 AUG 79	10
0(1)	002400	002564	0(2)	050066 050077	23 OCT 78	11
0(3)	002565	002601				
0(1)	002602	003357	0(2)	050100 050107	23 APR 82	10
0(3)	003360	003415				
0(1)	003416	003431			23 OCT 78	11
0(1)	003432	003605	0(2)	050110 050117	17 JUL 78	13
0(1)	003606	003655			10 DEC 79	17
0(3)	003656	003711				
0(1)	003712	004065			23 OCT 78	11
0(3)	004066	004120				
			0(0)	050120 050120	23 JUL 80	13
0(1)	004121	004144			21 MAR 80	17
0(1)	004145	005345	0(2)	050121 050234	11 DEC 79	16
0(3)	005346	005372				
0(1)	005373	005502	0(2)	050235 050242	18 MAR 80	14
0(3)	005503	005525				
0(1)	005527	006316	0(2)	050243 050343	23 APR 82	10
0(3)	006317	006413				
0(1)	006414	006505	0(2)	050344 050344	16 JUL 80	11
0(3)	006506	006543				
0(1)	006544	006613			20 MAR 80	18
0(3)	006614	006631				
0(1)	006632	007770	0(2)	050345 050450	23 JUL 80	13
0(3)	007771	010200				
0(1)	010201	011511			23 APR 82	10
0(3)	011512	011714				
0(1)	011715	045600	0(036)	050451 150450	27 JUN 83	21
0(3)	045601	047441				

TIME: 23.177 STORAGE: 17792/4/040777/073777

*HCEM.PASCODE

06/27/83 21:04:03

APPENDIX B
TEST RUNS

PROGRAM TEST (INPUT,OUTPUT) ;

(*
...THIS DATA CONTAINS NO ERRORS
AND SYMBOL TABLE DUMP IS NOT REQUIRED...
*)

TYPE DAYS = (M,T,W,TH,FR,SA,S) ;
WEEK = SET OF DAYS ;

VAR WK,WORK,FREE : WEEK ;
D : DAYS ; SS: WEEK ;

PROCEDURE CHECK ; RET : (K,L,M) ;
VAR D : DAYS ;
BEGIN WRITE (' ') ;
FOR D := M TO SU DO
IF D IN SS THEN WRITE ('X')
ELSE WRITE ('O') ;
WRITELN
END ;

BEGIN WORK := [] ; FREE := [] ;
WK := (M,SA,SU) ;
D := SA ; FREE := [D] + FREE + (SU) ;
SS := FREE ; CHECK ;

WORK := WK ; FREE ; SS := CHECK ; CHECK ;

IF FREE <= WK THEN WRITE ('O') ;
IF WK >= WORK THEN WRITE ('X') ;
IF NOT (WORK >= FREE) THEN WRITE (' JACK') ;
IF (SA) <= WORK THEN WRITE (' FORGET IT') ;

WRITELN

END.

NO ANALYSTS ***
ARE PROCESSED .
CORRUS ARE
DETECTED.

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```
PROGRAM TEST2 (INPUT,OUTPUT) ;
  PRINTABLES
  (*
    ..THIS DATA IS SAME AS PROGRAM TEST1
    BUT SYMBOL TABLE DUMP IS NOW REQUIRED..
  *)
  TYPE DAYS = (M,T,W,TH,FR,SA,S) ;
  WEEK = SET OF DAYS ;
  VAR WK,WOPK,FREE : WEEK ;
  D : DAYS ; SS: WEEK ;
  PROCEDURE CHECK ;
  VAR D : DAYS ; RET : (K,L,M) ;
```

* SYMBOL TABLE DUMP *

..INTERMEDIATE DUMP..

INTERNAL CODE=147674 IDENTIFIER= D
LLINK= *NIL* RLINK=147674 NEXT=147674 IDPTR=150336
CLASS=VARIABLE ,ACTUAL ,LEVEL= 1

****STRUCTURE****

INTERNAL STRUCTURE CODE=150336 SIZE=1
FORM=SCALAR ,DECLARED ,FIRST CONSTANT ID PTR=150120

INTERNAL CODE=147635 IDENTIFIER= K
LLINK= *NIL* RLINK=147611 NEXT= *NIL* IDPTR=147663
CLASS=CONSTANT ,CONST VALUE=0

****STRUCTURE****

INTERNAL STRUCTURE CODE=147663 SIZE=1
FORM=SCALAR ,DECLARED ,FIRST CONSTANT ID PTR=147565

INTERNAL CODE=147611 IDENTIFIER= L
LLINK= *NIL* RLINK=147565 NEXT=147635 IDPTR=147663
CLASS=CONSTANT ,CONST VALUE=1

INTERNAL CODE=147565 IDENTIFIER= M
LLINK= *NIL* RLINK= *NIL* NEXT=147611 IDPTR=147663
CLASS=CONSTANT ,CONST VALUE=2

INTERNAL CODE=147674 IDENTIFIER= RET
LLINK=147635 RLINK= *NIL* NEXT= *NIL* IDPTR=147663
CLASS=VARIABLE ,ACTUAL ,LEVEL= 1

```
1          15          BEGIN WRITE (' ') ;
           17          FOR D := M TO SU DO
           18            IF D IN SS THEN WRITE ('X') ELSE WRITE (
           19              WRITELN
           20              END ;
           21
```

* SYMBOL TABLE DUMP *

..INTERMEDIATE DUMP..

INTERNAL CODE=150663 IDENTIFIER= BOOLEAN
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150641
CLASS=TYPE

****STRUCTURE****

INTERNAL STRUCTURE CODE=150641 SIZE=1
FORM=SCALAR, DECLARED, FIRST CONSTANT ID PTR=150672

INTERNAL CODE=150716 NAME= CHAR
LLINK=150633 RLINK=150646 NEXT= *NIL* IDPTR=150724
CLASS=TYPE

****STRUCTURE****

INTERNAL STRUCTURE CODE=150724 SIZE=1
FORM=SCALAR, STANDARD

INTERNAL CODE=147755 NAME= CHECK
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR= *NIL*
CLASS=PROCEDURE, ACTUAL, LEVEL=0, PARAMPTR=147750

****STRUCTURE****

INTERNAL STRUCTURE CODE=147750 SIZE=0
FORM=PARAM LIST, FIRST PARAMETER PTR= *NIL*

INTERNAL CODE=150024 NAME= 0
LLINK=147755 RLINK= *NIL* NEXT=150004 IDPTR=150336
CLASS=VARIABLE, ACTUAL, LEVEL= 0

****STRUCTURE****

INTERNAL STRUCTURE CODE=150336 SIZE=1
FORM=SCALAR, DECLARED, FIRST CONSTANT ID PTR=150120

INTERNAL CODE=150347 NAME= DAYS
LLINK=150324 RLINK= *NIL* NEXT= *NIL* IDPTR=150336
CLASS=TYPE

INTERNAL CODE=150646 NAME= FALSE
LLINK=150347 RLINK=150170 NEXT= *NIL* IDPTR=150641
CLASS=CONSTANT, BOOLEAN, CONST VALUE=FALSE

INTERNAL CODE=150170 NAME= FR
LLINK= *NIL* RLINK=150044 NEXT=150214 IDPTR=150336
CLASS=CONSTANT, CONST VALUE=4

INTERNAL CODE=150044 NAME= FREE
LLINK= *NIL* RLINK= *NIL* NEXT=150024 IDPTR=150076
CLASS=VARIABLE, ACTUAL, LEVEL= 0

****STRUCTURE****

INTERNAL STRUCTURE CODE=150076 SIZE=29451204315
FORM=SET, ELSET PTR=150336

INTERNAL CODE=150766 NAME= INTEGER
LLINK=150716 RLINK=150742 NEXT= *NIL* IDPTR=150774
CLASS=TYPE

****STRUCTURE****

INTERNAL STRUCTURE CODE=150774 SIZE=1
FORM=SCALAR, STANDARD

INTERNAL CODE=150310 NAME= M
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150336
CLASS=CONSTANT, CONST VALUE=0

INTERNAL CODE=150742 NAME=MAXINT
LLINK=150310 RLINK=150730 NEXT=*NIL* IDPTR=150774
CLASS=CONSTANT INTEGER ,CONST VALUE=34359738367

INTERNAL CODE=150571 NAME=.NIL
LLINK=*NIL* RLINK=*NIL* NEXT=*NIL* IDPTR=*NIL*
CLASS=CONSTANT

INTERNAL CODE=150730 NAME=REAL
LLINK=150571 RLINK=150672 NEXT=*NIL* IDPTR=150736
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150736 SIZE=2
FORM=SCALAR ,STANDARD

INTERNAL CODE=150120 NAME=S
LLINK=*NIL* RLINK=*NIL* NEXT=150144 IDPTR=150336
CLASS=CONSTANT ,CONST VALUE=6

INTERNAL CODE=150144 NAME=SA
LLINK=150120 RLINK=150004 NEXT=150170 IDPTR=150336
CLASS=CONSTANT ,CONST VALUE=5

INTERNAL CODE=150004 NAME=SS
LLINK=*NIL* RLINK=*NIL* NEXT=*NIL* IDPTR=150076
CLASS=VARIABLE ,ACTUAL ,LEVEL=0

INTERNAL CODE=150264 NAME=T
LLINK=150144 RLINK=*NIL* NEXT=150310 IDPTR=150336
CLASS=CONSTANT ,CONST VALUE=1

INTERNAL CODE=150621 NAME=TEXT
LLINK=150264 RLINK=150214 NEXT=*NIL* IDPTR=150627
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150627 SIZE=34359738367
FORM=FILE ,FILE TYPE PTR=150724

INTERNAL CODE=150214 NAME=TH
LLINK=*NIL* RLINK=*NIL* NEXT=150240 IDPTR=150336
CLASS=CONSTANT ,CONST VALUE=3

INTERNAL CODE=150672 NAME=TRUE
LLINK=150621 RLINK=150240 NEXT=*NIL* IDPTR=150641
CLASS=CONSTANT BOOLEAN ,CONST VALUE=TRUE

INTERNAL CODE=150240 NAME=W
LLINK=*NIL* RLINK=150110 NEXT=150264 IDPTR=150336
CLASS=CONSTANT ,CONST VALUE=2

INTERNAL CODE=150110 NAME=WEEK
LLINK=*NIL* RLINK=150060 NEXT=*NIL* IDPTR=150076
CLASS=TYPE

INTERNAL CODE=150060 NAME=WK
LLINK=*NIL* RLINK=150052 NEXT=150052 IDPTR=150076
CLASS=VARIABLE ,ACTUAL ,LEVEL=0

INTERNAL CODE=150052 IDENTIFIER NAME= WORK
LLINK= *NIL* RLINK= *NIL* NEXT=150044 IDPTR=150076
CLASS=VARIABLE ,ACTUAL ,LEVEL= 0

```
1          22          BEGIN WORK := 0 ; FREE := 0 ;  
          23          WK := EM ; SUJ ;  
          24          D := SA ; FREE := [D] + FREE + [SUJ] ;  
          25  
          26          SS := FREE ; CHECK ;  
          27  
          28          WORK := WK - FREE ; SS := CHECK ; CHECK  
          29  
          30          IF FREE <= WK THEN WRITE ('D') ;  
          31          IF WK >= WORK THEN WRITE ('K') ;  
          32          IF NOT (WORK >= FREE) THEN WRITE (' JACK')  
          33          IF [SA] <= WORK THEN WRITE (' FORGET IT')  
          34  
          35          WRITELN  
          36          END.  
1  
*** END OF ANALYSIS ***  
36 RECORDS ARE PROCESSED  
0 ERRORS DETECTED.
```

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PROGRAM ALLERRORS (INPUT OUTPUT) ;

*** ERROR 20...POSITION 1 ***

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(* THIS SAMPLE RUN CONTAINS ALOT OF ERRORS.
IT IS NOT A MEANINGFULL PROGRAM YET IT IS
DESIGNED TO TEST THE RECOVERY ALGORITHM
OF PUPASCAL, RECALL THAT NO SEMANTIC ANALYSIS
IS PERFORMED FOR STATEMENTS.
)

CONST A = 9 ; B = ' ' ; C = 365499999999999934
1 2

*** ERROR 275...POSITION 1 ***
*** ERROR 203...POSITION 2 ***

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22

TYPE
A : RECORD
N : REAL ;
CASE M : INTEGER OF
1, 2 : (CASE A : CHAP OF
'L' : (NV : C4) ;
1

*** ERROR 104...POSITION 1 ***

23
24

'N', 'M' : (ZZ : 1..3) ;
'K', 'N' : (E : BOOLEAN) ;
1

*** ERROR 158...POSITION 1 ***

25

3, 'A' ; (FF : INTEGER) ;

*** ERROR 111...POSITION 1 ***

26
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30

CC : ARRAY [1..9] OF REAL) ;
END ;
VAR
AX : A ;
T, J, K : TEXT ;
1

*** ERROR 2...POSITION 1 ***
*** ERROR 18...POSITION 1 ***

31

T, AN : BADTYPE ;
1 2

*** ERROR 101...POSITION 1 ***
*** ERROR 104...POSITION 2 ***

32

SETS : SET CHAP ;
1

*** ERROR 3...POSITION 1 ***
*** ERROR 10...POSITION 1 ***
*** ERROR 18...POSITION 1 ***

33

ARRAYS : ARRAY [INTEGER, BOOLEAN, REAL] OF
1 2

*** ERROR 149...POSITION 1 ***
*** ERROR 109...POSITION 2 ***

34

ARRAY [1..3, 9] OF CHAR ;
1

*** ERROR 50...POSITION 1 ***
*** ERROR 1...POSITION 1 ***
*** ERROR 101...POSITION 1 ***
*** ERROR 18...POSITION 1 ***

*** ERROR 52...POSITION -1 ***
*** ERROR 58...POSITION 1 ***

*** 69 END OF ANALYSIS END ***
60 RECORDS ARE PROCESSED
69 ERRORS DETECTED.

INTERNAL STRUCTURE CODE=147542 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147750 VAR PTR= *NIL*

****IDENTIFIER****
INTERNAL CODE=147605 NAME = DAY
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150177
CLASS=FIELD

****IDENTIFIER****
INTERNAL CODE=147750 NAME = MO
LLINK=147605 RLINK=147565 NEXT= *NIL* IDPTR=147737
CLASS=FIELD

****STRUCTURE****

INTERNAL STRUCTURE CODE=147737 SIZE=1
FORM=SCALAR ,DECLARED ,FIRST CONSTANT ID PTR=147615

****IDENTIFIER****
INTERNAL CODE=147565 NAME = YEAR
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150774
CLASS=FIELD

****IDENTIFIER****
INTERNAL CODE=150016 NAME = DIVORCED
LLINK=147737 RLINK=147711 NEXT=150042 IDPTR=150114
CLASS=CONSTANT ,CONST VALUE=2

****STRUCTURE****

INTERNAL STRUCTURE CODE=150114 SIZE=1
FORM=SCALAR ,DECLARED ,FIRST CONSTANT ID PTR=147772

****IDENTIFIER****
INTERNAL CODE=147815 NAME = ETC
LLINK= *NIL* RLINK= *NIL* NEXT=147641 IDPTR=147737
CLASS=CONSTANT ,CONST VALUE=3

****IDENTIFIER****
INTERNAL CODE=147865 NAME = FEB
LLINK=147515 RLINK=147366 NEXT=147711 IDPTR=147737
CLASS=CONSTANT ,CONST VALUE=1

****IDENTIFIER****
INTERNAL CODE=147366 NAME = FEMALE
LLINK= *NIL* RLINK= *NIL* NEXT=147412 IDPTR=147440
CLASS=CONSTANT ,CONST VALUE=1

****STRUCTURE****

INTERNAL STRUCTURE CODE=147440 SIZE=1
FORM=SCALAR ,DECLARED ,FIRST CONSTANT ID PTR=147366

****IDENTIFIER****
INTERNAL CODE=147711 NAME = JAN
LLINK=147865 RLINK=147641 NEXT= *NIL* IDPTR=147737
CLASS=CONSTANT ,CONST VALUE=0

****IDENTIFIER****
INTERNAL CODE=147412 NAME = MALE
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=147440
CLASS=CONSTANT ,CONST VALUE=0

****IDENTIFIER****
INTERNAL CODE=147641 NAME = MARCH
LLINK=147412 RLINK= *NIL* NEXT=147665 IDPTR=147737
CLASS=CONSTANT ,CONST VALUE=2

INTERNAL CODE=1500867 NAME = MARRIED
LLINK=150015 RLINK=150042 NEXT= *NIL* IDPTR=150114
CLASS=CONSTANT ,CONST VALUE=0

INTERNAL CODE=147084 NAME = PEOPLE
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=147076
CLASS=VARIABLE ,ACTUAL ,LEVEL= 1

****STRUCTURE****

INTERNAL STRUCTURE CODE=147076 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147527 VAR PTR=147336

INTERNAL CODE=147356 NAME = BIRTH
LLINK= *NIL* RLINK=147326 NEXT= *NIL* IDPTR=147542
CLASS=FIELD

INTERNAL CODE=147202 NAME = DDATE
LLINK= *NIL* RLINK=147127 NEXT= *NIL* IDPTR=147542
CLASS=FIELD

INTERNAL CODE=147127 NAME = INDEPT
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150641
CLASS=FIELD

INTERNAL CODE=147255 NAME = MDATE
LLINK=147202 RLINK= *NIL* NEXT= *NIL* IDPTR=147542
CLASS=FIELD

INTERNAL CODE=147326 NAME = MS
LLINK=147255 RLINK= *NIL* NEXT= *NIL* IDPTR=150114
CLASS=FIELD

INTERNAL CODE=147522 NAME = NAME
LLINK=147356 RLINK=147451 NEXT= *NIL* IDPTR=147457
CLASS=FIELD

****STRUCTURE****

INTERNAL STRUCTURE CODE=147457 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147510 VAR PTR= *NIL*

INTERNAL CODE=147510 NAME = FIRST
LLINK= *NIL* RLINK=147502 NEXT=147502 IDPTR=150153
CLASS=FIELD

INTERNAL CODE=147502 NAME = LAST
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150153
CLASS=FIELD

INTERNAL CODE=147451 NAME = SEX
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=147440
CLASS=FIELD

****STRUCTURE****

INTERNAL STRUCTURE CODE=147336 SIZE=29451204315
FORM=TAGFIELD ,TAGFIELD PTR=147326 ,TAGTYPE PTR=150114 ,FIRST VAR

****STRUCTURE****

INTERNAL STRUCTURE CODE=147152 SIZE=29451204315

****STRUCTURE****

INTERNAL STRUCTURE CODE=147225 SIZE=29451204315
FORM=VARIANT ,NEXT VARIANT PTR=147300, SUBREC: =147214, VARIANT

****STRUCTURE****

INTERNAL STRUCTURE CODE=147300 SIZE=29451204315
FORM=VARIANT ,NEXT VARIANT PTR=147310, SUBREC: =147267, VARIANT

****STRUCTURE****

INTERNAL STRUCTURE CODE=147310 SIZE=29451204315
FORM=VARIANT ,NEXT VARIANT PTR= *NIL*, SUBREC: =147267, VARIANT

****STRUCTURE****

INTERNAL STRUCTURE CODE=147267 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147255 VAR PTR= *NIL*

****IDENTIFIER****

INTERNAL CODE=147202 NAME = DDATE
LLINK= *NIL* RLINK=147127 NEXT= *NIL* IDPTR=147542
CLASS=FIELD

****IDENTIFIER****

INTERNAL CODE=147127 NAME = INDEPT
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150641
CLASS=FIELD

****IDENTIFIER****

INTERNAL CODE=147255 NAME = MDATE
LLINK=147202 RLINK= *NIL* NEXT= *NIL* IDPTR=147542
CLASS=FIELD

****STRUCTURE****

INTERNAL STRUCTURE CODE=147214 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147202 VAR PTR= *NIL*

****IDENTIFIER****

INTERNAL CODE=147202 NAME = DDATE
LLINK= *NIL* RLINK=147127 NEXT= *NIL* IDPTR=147542
CLASS=FIELD

****IDENTIFIER****

INTERNAL CODE=147127 NAME = INDEPT
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150641
CLASS=FIELD

****STRUCTURE****

INTERNAL STRUCTURE CODE=147141 SIZE=29451204315
FORM=RECORD ,FIRST FIELD PTR=147127 VAR PTR= *NIL*

****IDENTIFIER****

INTERNAL CODE=147127 NAME = INDEPT
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150641
CLASS=FIELD

****IDENTIFIER****

INTERNAL CODE=147534 NAME = PERSON
LLINK=147064 RLINK= *NIL* NEXT= *NIL* IDPTR=147076
CLASS=TYPE

****IDENTIFIER****

INTERNAL CODE=147772 NAME = SINGLE
LLINK=147534 RLINK=150125 NEXT=150016 IDPTR=150114
CLASS=CONSTANT ,CONST VALUE=3

****IDENTIFIER****

INTERNAL CODE=150125 NAME = STATUS
LLINK= *NIL* RLINK= *NIL* NEXT= *NIL* IDPTR=150114
CLASS=TYPE

INTERNAL CODE=150641, IDENTIFIER=150641, MIDWORD
LINK=150772, LINK=150772, NEXT=150646 TOP=150114
CLASS=CONSTANT, CONST VALUE=1

1

SYMBOL TABLE DUMP

34 BEGIN END :

...INTERMEDIATE DUMP...

INTERNAL CODE=150633, IDENTIFIER=150633, BOOLEAN
LINK=150633, LINK=150633, NEXT=NULL, TOP=150641
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150646, IDENTIFIER=150646, FIRST CONSTANT ID OF=150672

INTERNAL CODE=150716, IDENTIFIER=150716, CHAR
LINK=150633, LINK=150646, NEXT=NULL, TOP=150724
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150724, IDENTIFIER=150724, STANDARD
FORM=SCALAR, SIZE=1

INTERNAL CODE=150733, IDENTIFIER=150733, CHARACTER
LINK=150733, LINK=150733, NEXT=NULL, TOP=150774
CLASS=CONSTANT, INTEGER, CONST VALUE=1

STRUCTURE

INTERNAL STRUCTURE CODE=150774, IDENTIFIER=150774, STANDARD
FORM=SCALAR, SIZE=1

INTERNAL CODE=150646, IDENTIFIER=150646, FALSE
LINK=150633, LINK=150646, NEXT=NULL, TOP=150641
CLASS=CONSTANT, BOOLEAN, CONST VALUE=FALSE

INTERNAL CODE=150786, IDENTIFIER=150786, INTEGER
LINK=150742, LINK=150742, NEXT=NULL, TOP=150774
CLASS=TYPE

INTERNAL CODE=150742, IDENTIFIER=150742, MAXINT
LINK=150742, LINK=150733, NEXT=NULL, TOP=150774
CLASS=CONSTANT, INTEGER, CONST VALUE=34359738367

INTERNAL CODE=150571, IDENTIFIER=150571, NULL
LINK=150571, LINK=NULL, NEXT=NULL, TOP=150774
CLASS=CONSTANT, NULL

INTERNAL CODE=150730, IDENTIFIER=150730, REAL
LINK=150672, LINK=150672, NEXT=NULL, TOP=150736
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150736, IDENTIFIER=150736, SIZE=2

INTERNAL CODE=150627, FIRM=, TESTREC
LINK=, NIL*, RLINK=, NIL*, NEXT=, NIL*, TOPTR=, NIL*,
CLASS=PROCEDURE, ACTUAL, LEVEL=0, PARAM=PTR=150277

STRUCTURE

INTERNAL STRUCTURE, PARM=27, PTR=, SIZE=0

INTERNAL CODE=150627, FIRM=, TEXT
LINK=150304, RLINK=, NIL*, NEXT=, NIL*, TOPTR=150627
CLASS=TYPE

STRUCTURE

INTERNAL STRUCTURE CODE=150627, SIZE=34359738367
FORM=FILE, FILE TYPE PTR=150724

INTERNAL CODE=150672, FIRM=, TRJF
LINK=150521, RLINK=, NIL*, NEXT=, NIL*, TOPTR=150641
CLASS=CONSTANT, BOOLEAN, CONST VALUE=TRUE

1 35 BEGIN

*** 1 END OF RECORDS ARE PROCESSED .
35 ERRORS DETECTED.

APPENDIX C
SYNTAX OF STANDARD PASCAL

1. `<program> ::= <program heading><body>`.
2. `<program heading>`:
`::= PROGRAM IDENT (<ident list>);`
3. `<ident list> ::= IDENT {, IDENT}0`
4. `<body> ::= <declarations><routine declarations>`
`<compound statement>`
5. `<declarations> ::= <label declarations>`
`<constant definitions>`
`<type definitions>`
`<var declarations>`
6. `<label declarations>`
`::= LABEL INTCONST {, INTCONST}0; | empty`
7. `<constant definitions>`
`::= CONST {IDENT = <constant>;}1 | empty`
8. `<constant> ::= INTCONST | REALCONST | IDENT | ADDOP INTCONST | ADDOP`
`REALCONST | ADDOP IDENT | STRINGCONST | CHARCONST`
9. `<type definitions> ::= TYPE {IDENT = <type>;}1 | <empty>`
10. `<type> ::= IDENT | PACKED <structured type>`
`SET OF <simple type> |`
`ARRAY [<index list>] OF <type> |`
`RECORD <field list> END |`
`FILE OF <type> |`
`<simple type>`
11. `<structured type> ::= SET OF <simple type> |`
`ARRAY [<index list>] OF <type> |`
`RECORD <field list> END |`
`FILE OF <type>`

12. $\langle \text{index list} \rangle ::= \langle \text{simple type} \rangle \{, \langle \text{simple type} \rangle\}_0$
13. $\langle \text{simple type} \rangle ::= (\langle \text{ident list} \rangle) \mid \text{IDENT} \langle \text{subrange tail} \rangle \mid$
 $\langle \text{constant} \rangle \dots \langle \text{constant} \rangle$
14. $\langle \text{subrange tail} \rangle ::= \dots \langle \text{constant} \rangle \mid \langle \text{empty} \rangle$
15. $\langle \text{field list} \rangle ::= \langle \text{fixed part} \rangle \langle \text{variant part} \rangle$
16. $\langle \text{fixed part} \rangle ::= \text{IDENT} \{, \text{IDENT}\}_0 : \langle \text{type} \rangle \{; \langle \text{ident list} \rangle : \langle \text{type} \rangle\}_0$
 $\mid \langle \text{empty} \rangle$
17. $\langle \text{variant part} \rangle ::= \underline{\text{CASE}} \text{IDENT} \langle \text{tagtype} \rangle \underline{\text{OF}} \langle \text{variant list} \rangle$
18. $\langle \text{tagtype} \rangle ::= : \text{IDENT} \mid \langle \text{empty} \rangle$
19. $\langle \text{variant list} \rangle ::= \langle \text{variant} \rangle \{; \langle \text{variant} \rangle\}_0$
20. $\langle \text{variant} \rangle ::= \langle \text{constant list} \rangle : (\langle \text{field part} \rangle$
21. $\langle \text{field part} \rangle ::=) \mid \langle \text{field list} \rangle$
22. $\langle \text{constant list} \rangle ::= \langle \text{constant} \rangle \{, \langle \text{constant} \rangle\}_0$
23. $\langle \text{var declarations} \rangle ::= \underline{\text{VAR}} \{ \text{IDENT} \{, \text{IDENT} \}_0 : \langle \text{type} \rangle \}_1 \mid \langle \text{empty} \rangle$
24. $\langle \text{routine declarations} \rangle ::=$
 $\underline{\text{PROCEDURE}} \text{IDENT} \langle \text{formal parameter section} \rangle ; \langle \text{routine tail} \rangle ; \mid$
 $\underline{\text{FUNCTION}} \text{IDENT} \langle \text{formal parameter section} \rangle : \text{IDENT} ; \langle \text{routine}$
 $\text{tail} \rangle ; \mid \langle \text{empty} \rangle$
25. $\langle \text{routine tail} \rangle ::= \text{FORWARD} \mid \langle \text{body} \rangle$
26. $\langle \text{formal parameter section} \rangle ::= (\langle \text{formal parameters} \rangle) \mid \langle \text{empty} \rangle$
27. $\langle \text{formal parameters} \rangle ::= \langle \text{formal parameter} \rangle \{; \langle \text{formal parameter} \rangle\}_0$

28. <formal parameter> ::=
 PROCEDURE <identlist> | FUNCTION <parameter group>
 VAR <parameter group> | <parameter group>
29. <parameter group> ::= <identlist>: IDENT
30. <compound statement> ::= BEGIN <statement list> END
31. <statement list> ::= <statement> { ; <statement> }_o
32. <statement> ::= INTCONST: <unlabelled statement> |
 <unlabelled statement>
33. <unlabelled statement> ::=
 BEGIN <statement list> END |
 IF <expression> THEN <statement> <else part> |
 CASE <expression> OF <case list> END |
 WHILE <expression> DO <statement> |
 REPEAT <statement list> UNTIL <expression> |
 FOR IDENT := <expression> <for tail> |
 WITH <record-var list> DO <statements> |
 <simple statement>
34. <else part> ::= ELSE <statement> | <empty>
35. <for tail> ::= TO <expression> DO <statement> |
 DOWNTO <expression> DO <statement>
36. <simple statement> ::= GOTO INTCONST |
 IDENT <assignment> | <empty>
37. <assignment> ::= (<expression list>) |
 <var tail> := <expression>
38. <var tail> ::= @ <dot part> |
 [<expression list>] <dot part> |
 <dot part>
39. <dot part> ::= .<variable> | <empty>

40. $\langle \text{variable} \rangle ::= \text{IDENT } \langle \text{var tail} \rangle$
41. $\langle \text{record var list} \rangle ::= \langle \text{variable} \rangle \{, \langle \text{variable} \rangle\}_o$
42. $\langle \text{expression list} \rangle ::= \langle \text{expression} \rangle \{, \langle \text{expression} \rangle\}_o$
43. $\langle \text{expression} \rangle ::= \langle \text{simple expression} \rangle \{ \text{RELOP } \langle \text{simple expression} \rangle \}_o$
44. $\langle \text{simple expression} \rangle ::= \text{ADDOP } \langle \text{term} \rangle \{ \text{ADDOP } \langle \text{term} \rangle \}_o \mid$
 $\langle \text{term} \rangle \{ \text{ADDOP } \langle \text{term} \rangle \}_o$
45. $\langle \text{term} \rangle ::= \langle \text{factor} \rangle \{ \text{MULOP } \langle \text{factor} \rangle \}_o$
46. $\langle \text{factor} \rangle ::= \text{INTCONST} \mid \text{REALCONST} \mid \text{STRINGCONST} \mid$
 $\text{CHARCONST} \mid [\langle \text{subset list} \rangle \mid$
 $\text{NOT } \langle \text{factor} \rangle \mid (\langle \text{expression} \rangle) \mid$
 $\text{IDENT } \langle \text{factor tail} \rangle$
47. $\langle \text{factor tail} \rangle ::= (\langle \text{expression list} \rangle) \mid \langle \text{var tail} \rangle$
48. $\langle \text{subset list} \rangle ::= [\mid \langle \text{element list} \rangle]$
49. $\langle \text{element list} \rangle ::= \langle \text{element} \rangle \{, \langle \text{element} \rangle\}_o$
50. $\langle \text{element} \rangle ::= \langle \text{expression} \rangle \langle \text{range part} \rangle$
51. $\langle \text{range part} \rangle ::= .. \langle \text{expression} \rangle \mid \langle \text{empty} \rangle$
52. $\langle \text{case list} \rangle ::= \langle \text{case list element} \rangle \{ ; \langle \text{case list element} \rangle \}$
53. $\langle \text{case list element} \rangle ::= \langle \text{constant list} \rangle : \langle \text{statement} \rangle$
54. $\langle \text{empty} \rangle ::= =$

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